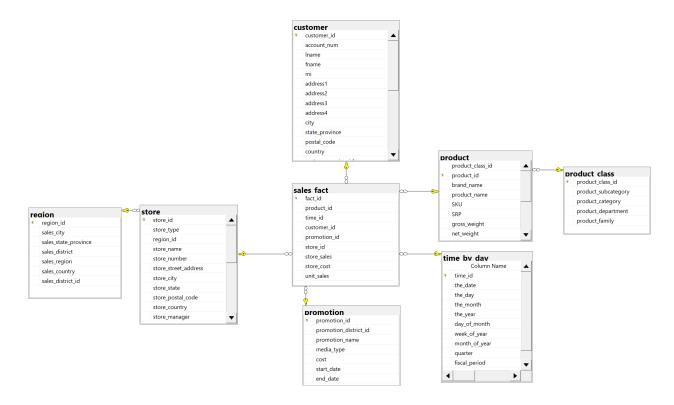
Second MidTerm of Decision Support Systems / Decision Support Databases

It is forbidden to consult any material during the test. Duration of written exam is 1.5h.

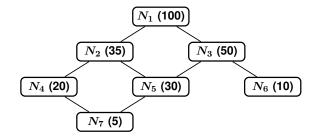
1. (4 points) Consider the FoodMart datawarehouse.



Write an **analytic** SQL query to solve the following problem: the list of store_city and product_id's such that the product_id is the most sold one (wrt store_sales) in at least one store_id of the city.

Solution:

2. (4 points) Let us consider the following lattice of possible candidate views to materialize. The numbers associated with the nodes represent the view size, measured in terms of the number of tuples in the view. Select 2 views to materialize, different from N_1 , with the greedy algorithm HRU.



Solution:

Benefits at the first step are:

$$N_2 (100 - 35) \cdot 4 = 260$$

$$N_3 \ (100 - 50) \cdot 4 = 200$$

$$N_4 (100 - 20) \cdot 2 = 160$$

$$N_5 (100 - 30) \cdot 2 = 140$$

$$N_6 (100 - 10) \cdot 4 = 90$$

$$N_7 (100 - 5) \cdot 1 = 95$$

Benefits at the second step are:

$$N_3 (100 - 50) \cdot 2 + 0 \cdot 2 = \underline{100}$$

$$N_4 (35-20) \cdot 2 = 30$$

$$N_5 (35 - 30) \cdot 2 = 10$$

$$N_6 (100 - 10) \cdot 4 = 90$$

$$N_7 (35-5) \cdot 1 = 30$$

Thus, HRU selects $\{N1, N2, N3\}$.

3. (4 points) Consider the data mart about bus tickets, without null values, and the query:

Travellers(<u>TPk</u>, Name, Address) TravelPlans(<u>PlanPk</u>, PlanName)

Trips(TFk, PlanFk, Day, Month, Year, Duration, Charge)

Q: SELECT PlanPk, PlanName, SUM(Charge) AS SC

FROM Trips, TravelPlans

WHERE PlanFk = PlanPk AND Year = 2025

GROUP BY PlanPk, PlanName;

Show how to rewrite the query Q using the view V, if possible.

V: SELECT PlanFk, Year, SUM(Charge) AS SC

FROM Trips

 $\begin{array}{ll} \text{WHERE} & \text{Year} >= 2020 \\ \text{GROUP BY} & \text{PlanFk, Year;} \end{array}$

Solution: Let us use the approach with a *compensation on the view* starting from the logical query plans of query and view.

PlanPK, PlanName $^{\gamma}$ SUM(Charge) **AS** SC $\sigma_{\text{Year} = 2025} \qquad \qquad \text{PlanFK, Year}^{\gamma}$ SUM(Charge) **AS** SC $\sigma_{\text{Year} = 2025} \qquad \qquad \sigma_{\text{Year} >= 2025}$ $\sigma_{\text{Year} >= 2025}$ $\sigma_{\text{Trips}} \qquad \qquad \sigma_{\text{Trips}} >= 2025$ $\sigma_{\text{Year} >= 2025}$ $\sigma_{\text{Year} >= 2025}$

Figure 1: Query and view logical query plans

(join) There is a partial match, which can be compensated by joining the view with TravelPlans. The compensation can float since PlanFk is in the output of the view.

(selection) There is a partial match, which can be compensated by restricting to Year=2025. The compensation can float since Year is in the output of the view.

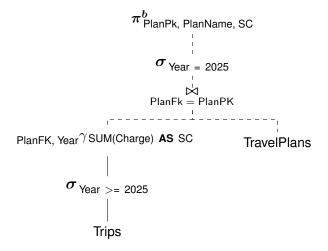
(grouping) Let us check the functional dependencies for $g(Q) = \{ \text{PlanPk, PlanName} \}$ and $g(V) = \{ \text{PlanFk, Year } \}$:

 $g(Q) \to q(V)$ is valid, in fact {PlanPk, PlanName} $^+ = \{$ PlanPk, PlanName, PlanFk, Year, ...} since PlanPk=PlanFk and Year=2025.

 $g(V) \to q(Q)$ is valid, in fact $\{\text{PlanFk, Year}\}^+ = \{\text{PlanFk, Year, PlanPk, PlanName, } \dots \}$ since PlanPk=PlanFk and PlanPk is a key of TravelPlans.

Therefore, the compensation consists only of a projection.

In summary, we have the following compensation:



and the query Q can be rewritten as:

SELECT PlanPk, PlanName, SC
FROM V, TravelPlans
WHERE PlanPk AND Year = 2025;

- 4. (4 points) Answer the following questions with reference to the FoodMart datawarehouse:
 - (a) assuming a BMFCJ index on customer.city write the physical query plan for the following query:

V: SELECT product_id, SUM(store_sales) AS TotalSales FROM sales_fact AS S, customer AS C

WHERE S.customer_id = C.customer_id AND C.city='Los Angeles'

GROUP BY product_id;

- (b1) (for Italian speaking students) discuss the benefits of the optimizations implemented in the SadasDB DBMS over the FoodMart databases;
- **(b2)** (for non-Italian speaking students) write two SQL queries on FoodMart in order to discuss when column-based storage is better than traditional (heap-based) storage.

Solution: (a) Let us call the index as Idx. We have:

