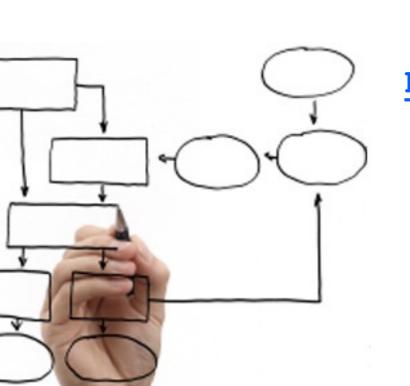
## Business Processes Modelling MPB (6 cfu, 295AA)

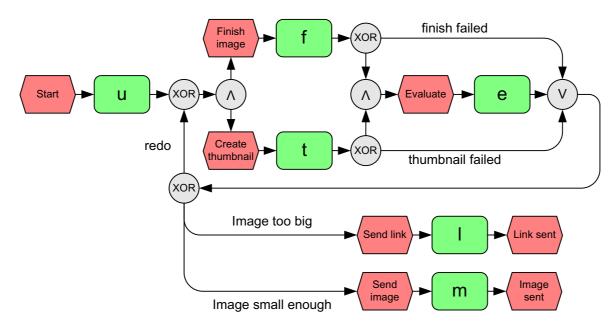


#### Roberto Bruni

http://www.di.unipi.it/~bruni

16a - EPC analysis

### Object



We overview the main challenges that arise when analysing EPC diagrams with Petri nets

Ch. 6 of Business Process Management: Concepts, Languages, Architectures

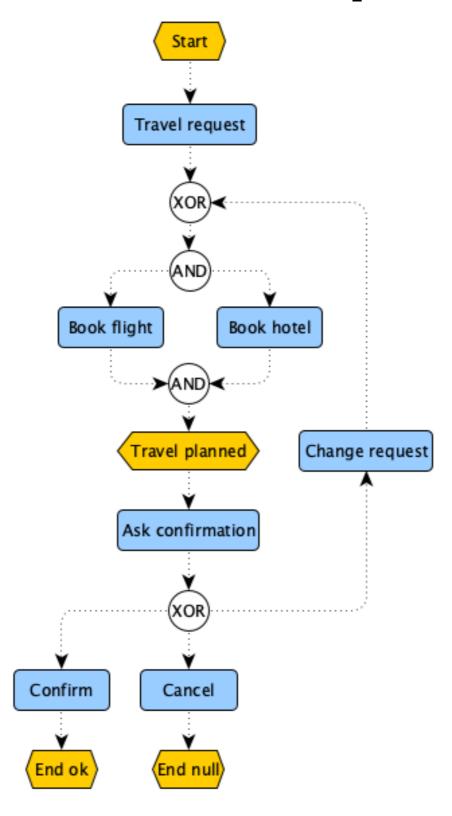
## EPC Diagrams

# EPC ingredients at a glance

**Event Function** Connectors XOR **Control Flow** 

M. Weske: Business Process Management, Springer-Verlag Berlin Heidelberg 2007

### EPC: Example



#### EPC Semantics

### Sound EPC diagrams

We exploit the formal semantics of nets to give unambiguous semantics to EPC diagrams

We transform EPC diagrams to Workflow nets: the EPC diagram is sound if its net is so

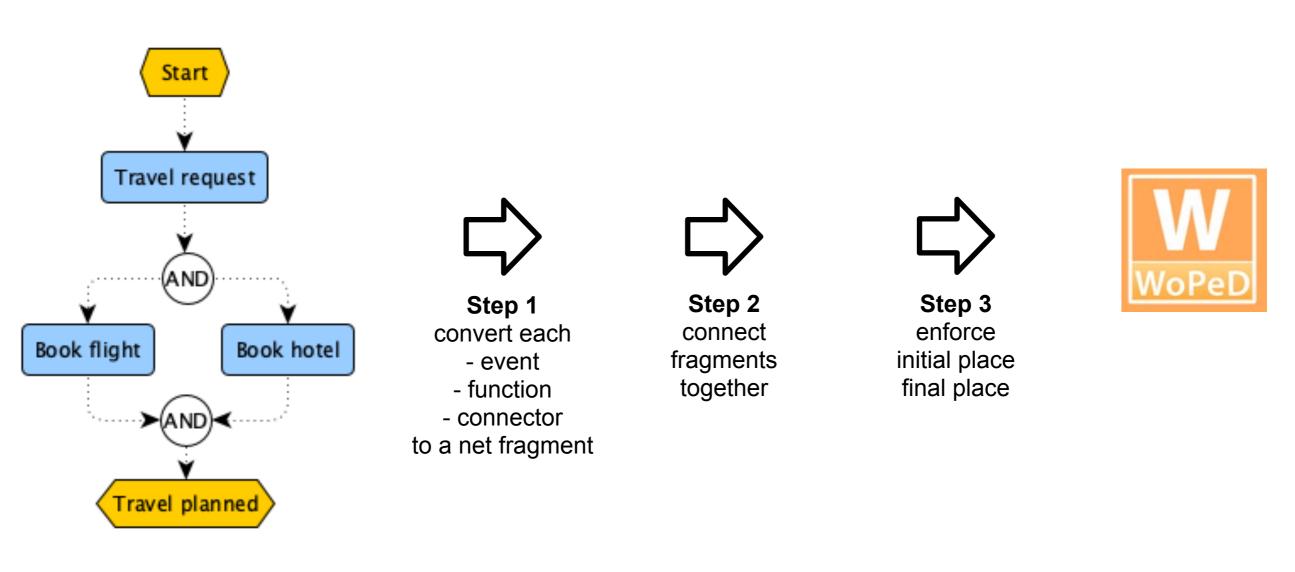
We can reuse the verification tools to check if the net is sound

Is there a unique way to proceed? Not necessarily!

# Translation of EPC to Petri nets

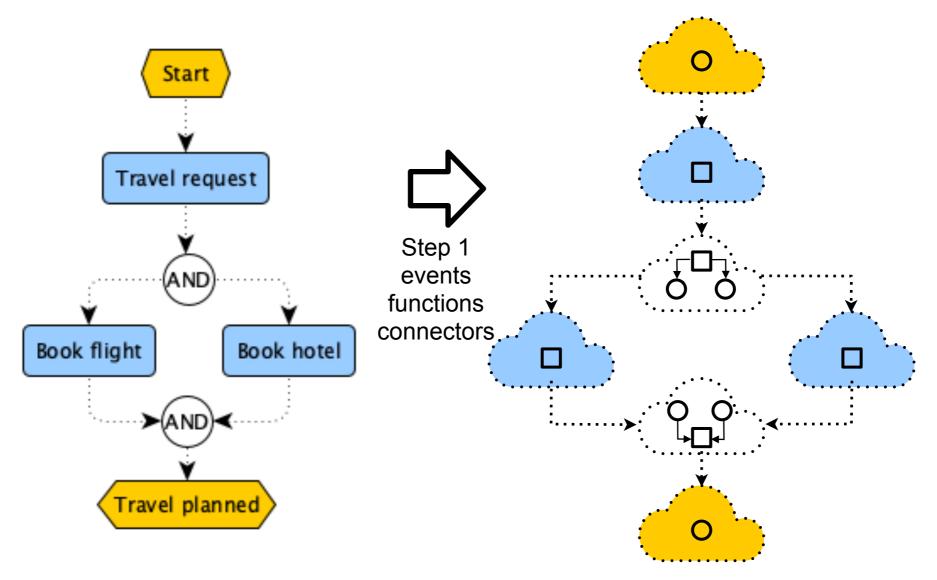
#### The idea

#### From EPC to wf nets in three steps



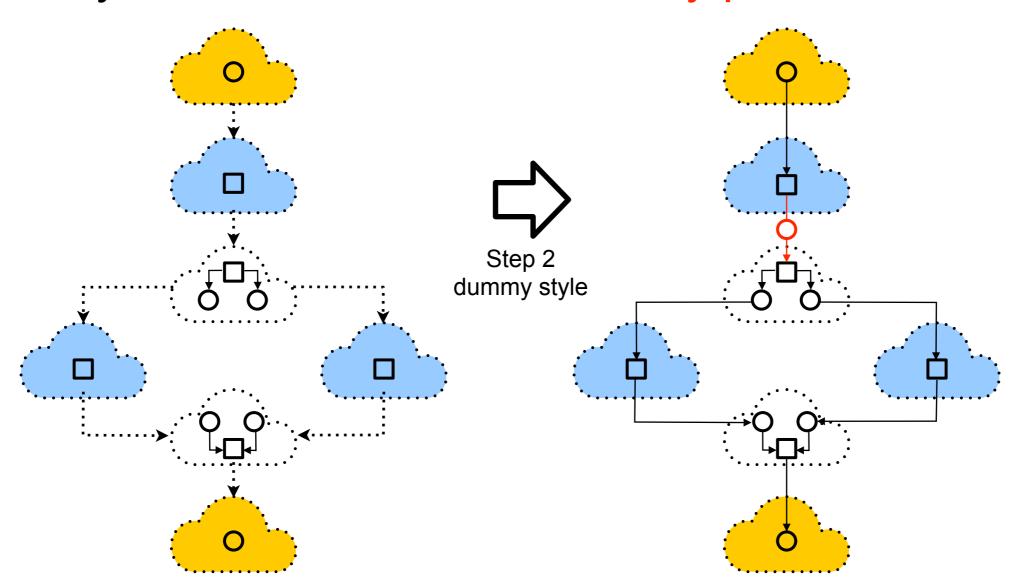
#### Step 1

We replace each event, function and connector separately with small net fragments



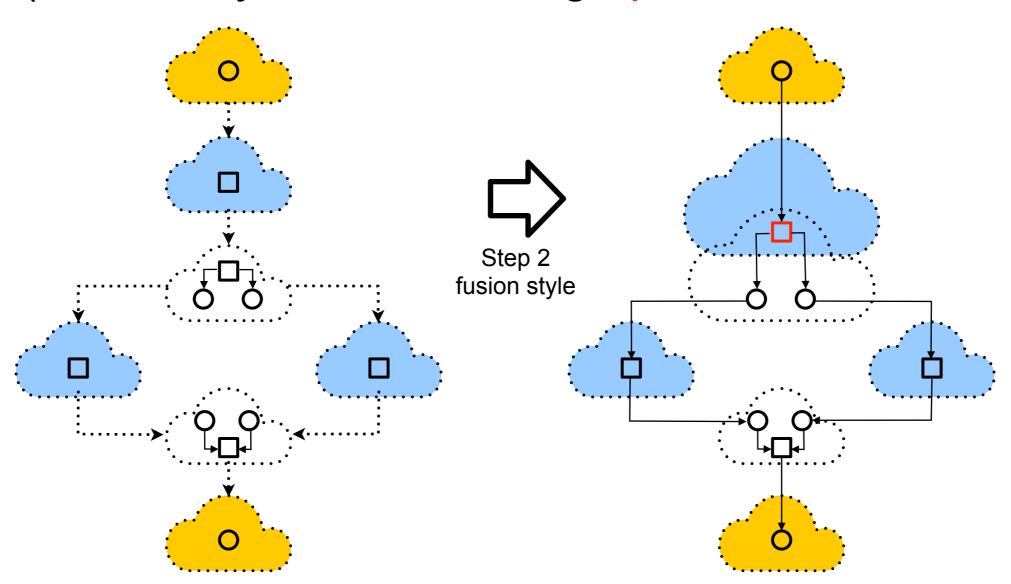
## Step 2: dummy style

Then we connect the fragments together (we may decide to introduce dummy places / transitions)



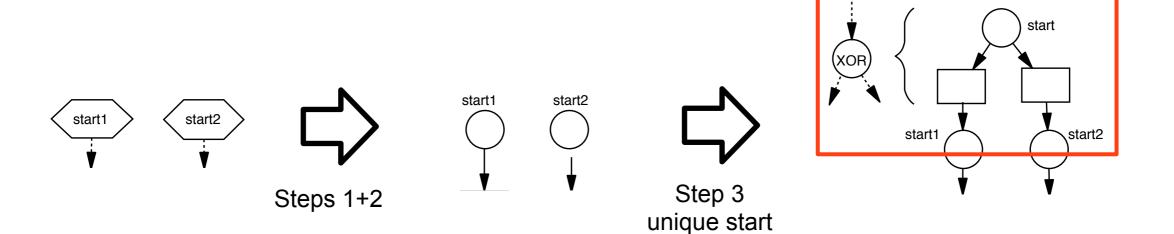
### Step 2: fusion style

Then we connect the fragments together (or we may decide to merge places / transitions)

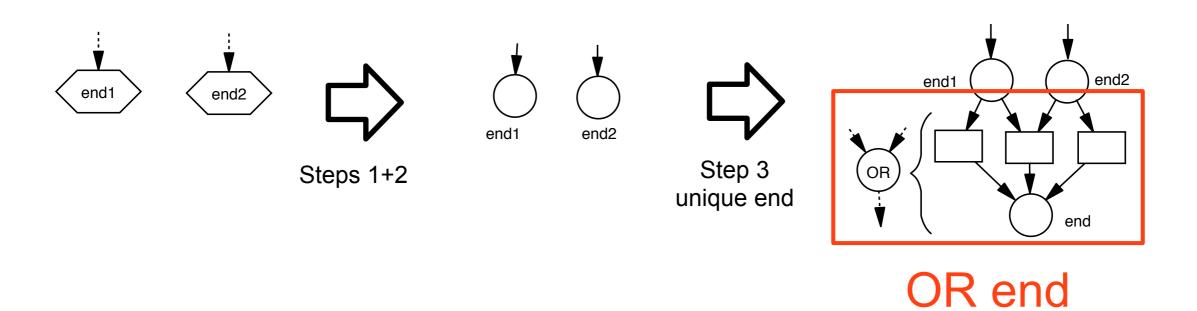


#### Step 3: unique start

#### **XOR** start



## Step 3: unique end



(sometimes XOR/AND can be preferred)

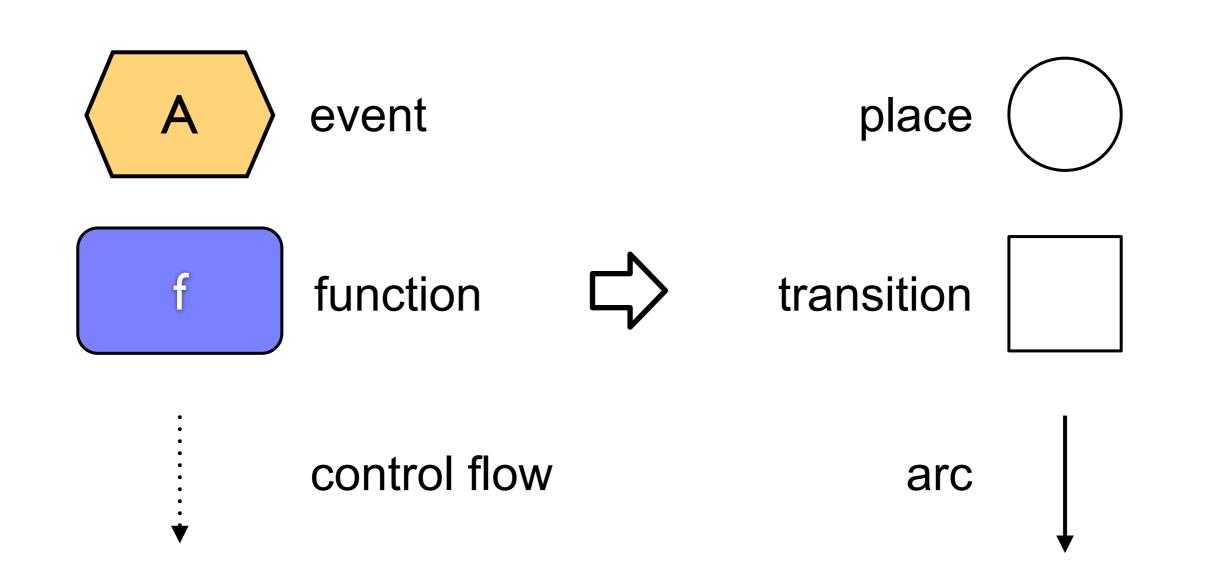
#### Three approaches

#### We overview three different translations

n.	ingenuity	style	applicability	outcome
1st	easy	fusion	any EPC	likely unsound, (relaxed soundness)
2nd	medium, context dependent	(dummy)	simplified EPC: event function alternation, no OR connectors	free-choice net
3rd	hard, context dependent	dummy	decorated EPC: join-split correspondence, OR policies	accurate analysis

#### Commonalities

net fragment



# First attempt (straight translation)

#### Relaxed Soundness of Business Processes

Juliane Dehnert<sup>1,\*</sup> and Peter Rittgen<sup>2</sup>

<sup>1</sup> Institute of Computer Information Systems, Technical University Berlin, Germany dehnert@cs.tu-berlin.de

<sup>2</sup> Institute of Business Informatics, University Koblenz-Landau, Germany rittgen@uni-koblenz.de

K.R. Dittrich, A. Geppert, M.C. Norrie (Eds.): CAiSE 2001, LNCS 2068, pp. 157–170, 2001. © Springer-Verlag Berlin Heidelberg 2001

#### Rationale

EPC success is due to its simplicity

EPC diagrams lack a consistent semantics: ambiguous and flawed process descriptions can arise in the design phase

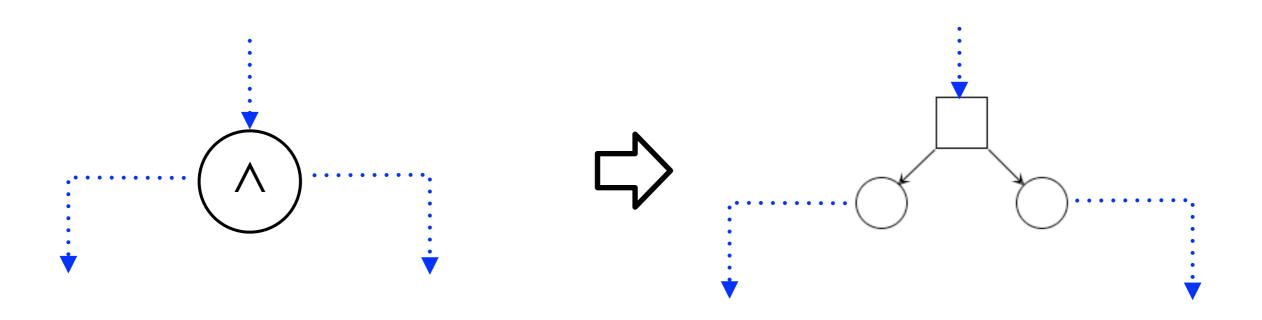
it is important to find out flaws as soon as possible

therefore

we need to fix a **formal representation** that **preserves all ambiguities** 

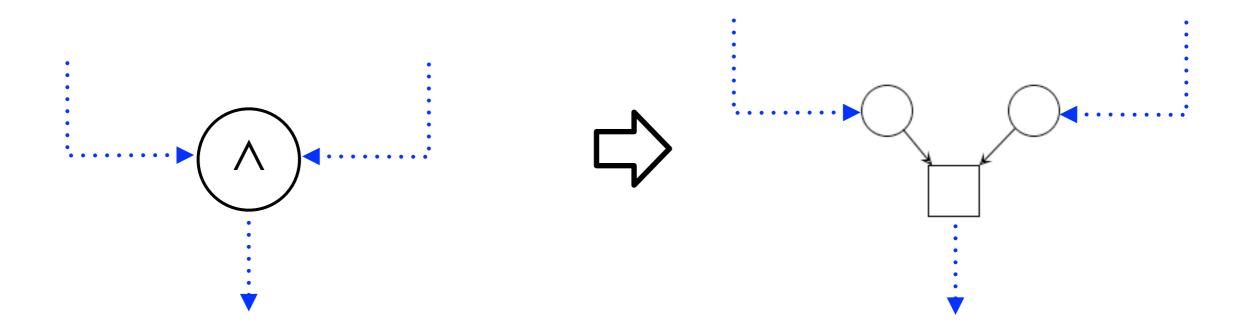
## Step 1: AND split

net fragment



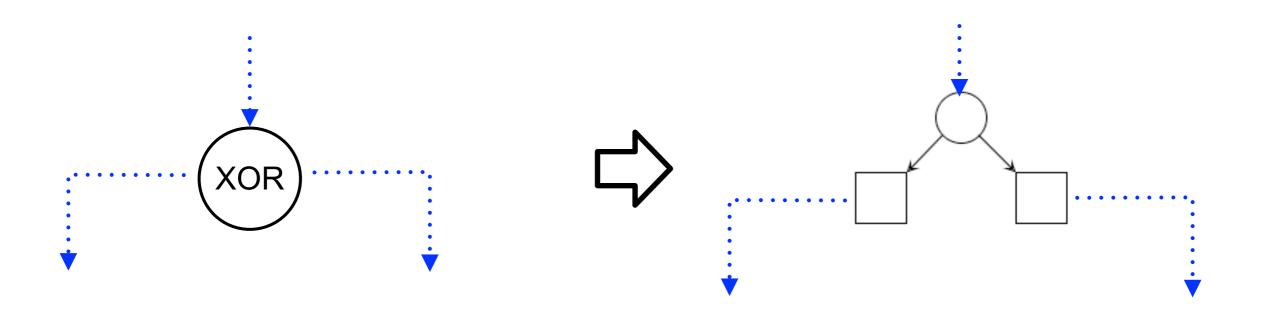
### Step 1: AND join

net fragment



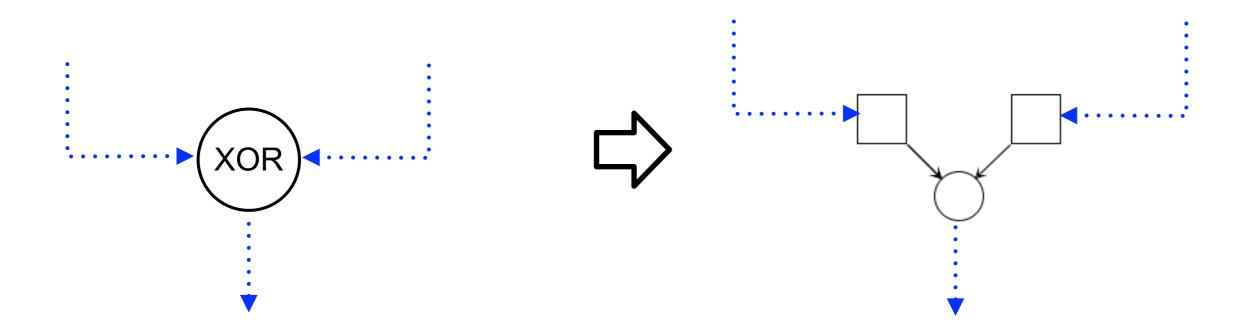
## Step 1: XOR split

net fragment



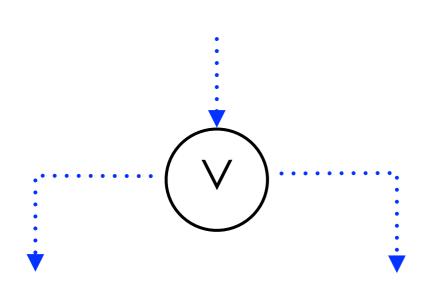
### Step 1: XOR join

net fragment

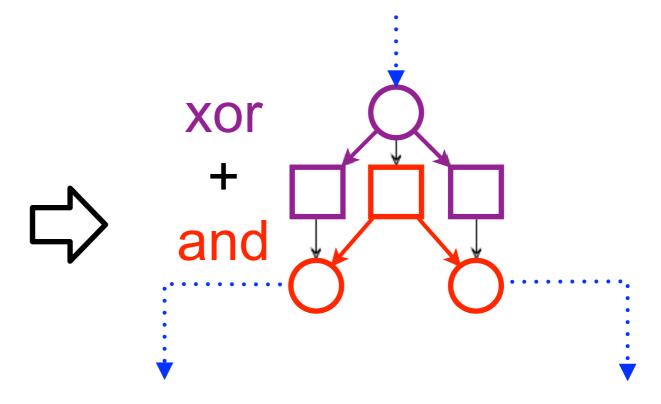


## Step 1: OR split

#### **EPC** element



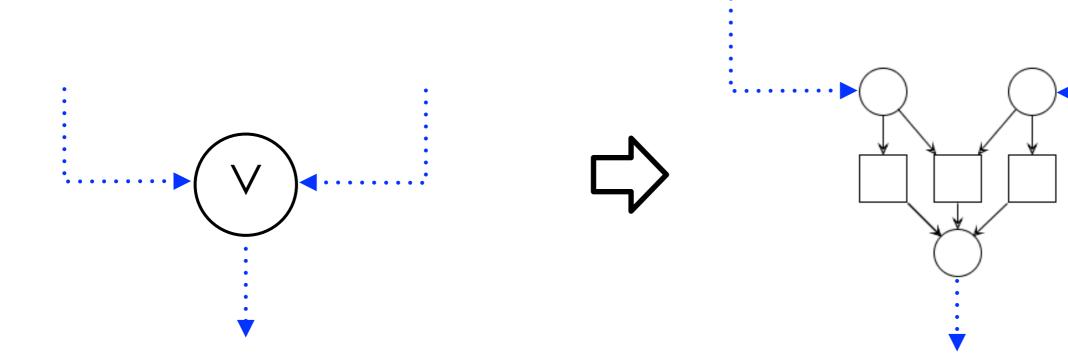
#### net fragment



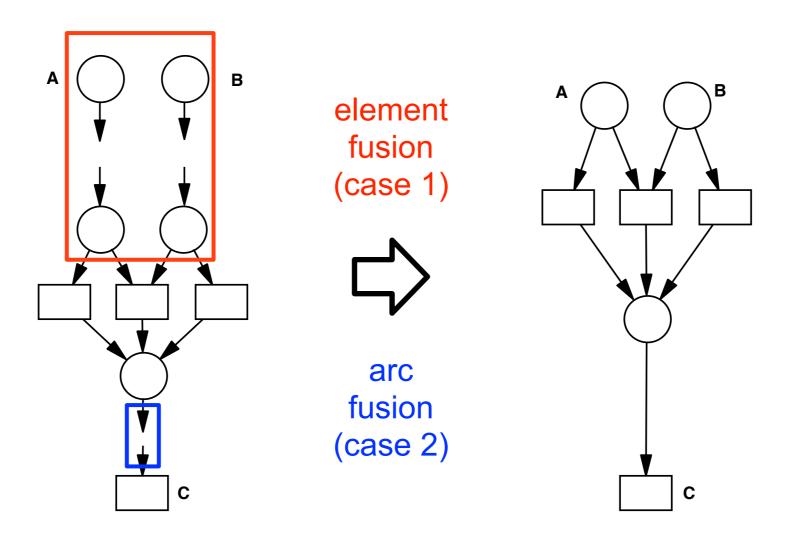
## Step 1: OR join

#### **EPC** element

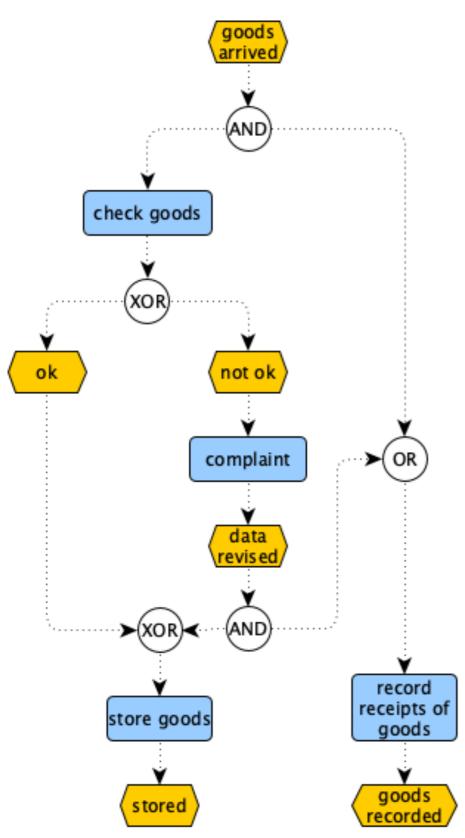
#### net fragment



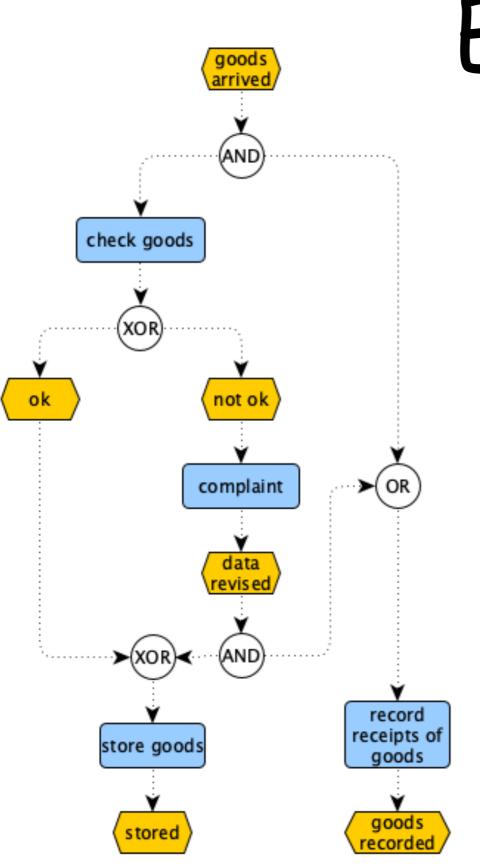
## Step 2: fusion style



## Example

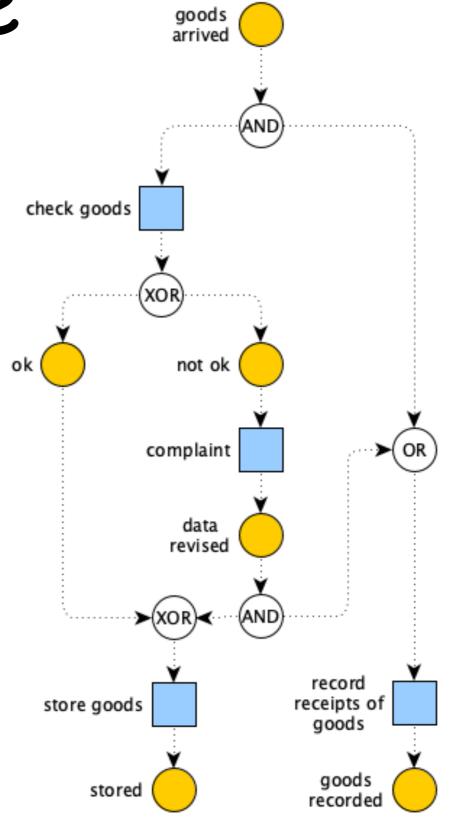


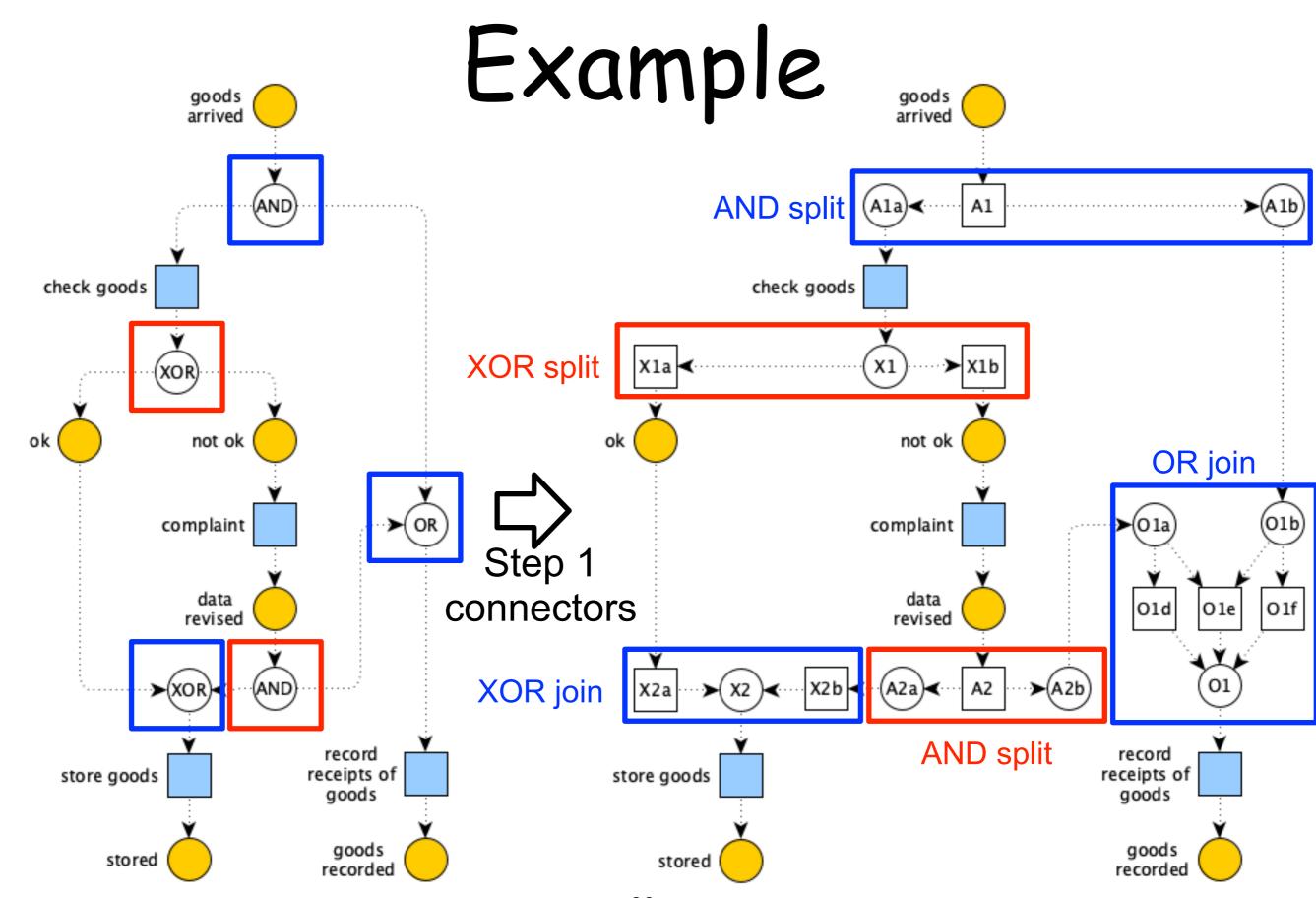
Sound?



Example

Step 1 events and functions

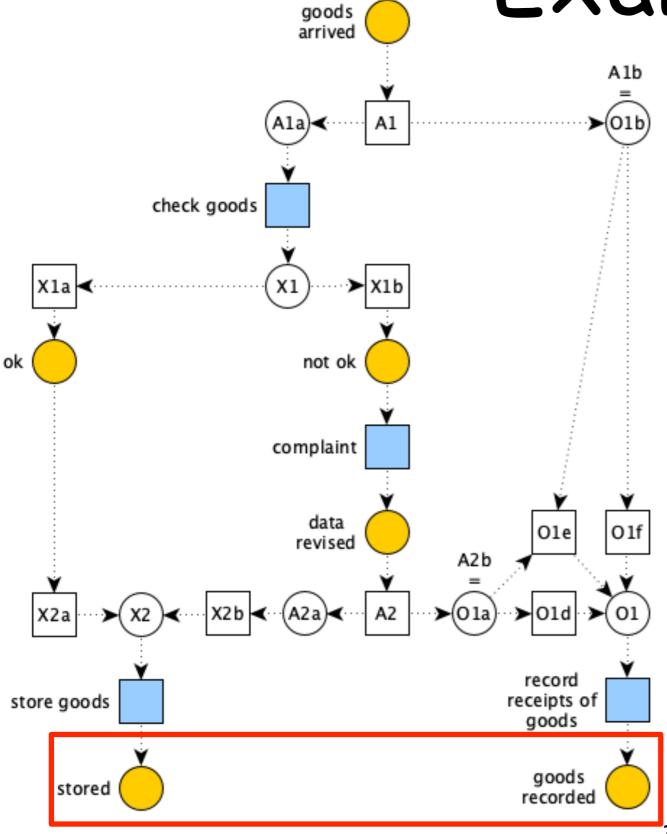




#### Example goods arrived Α1 check goods ➤ X1b not ok (O1b) complaint **≻**(O1a Step 2 data O1f O1d O1e fusion revised 01 record receipts of store goods goods goods stored recorded

#### Example goods arrived A1b =**≻**(01b check goods ➤ X1b not ok complaint Step 2 data O1e O1f fusion revised A2b > 01d record receipts of store goods goods goods stored recorded

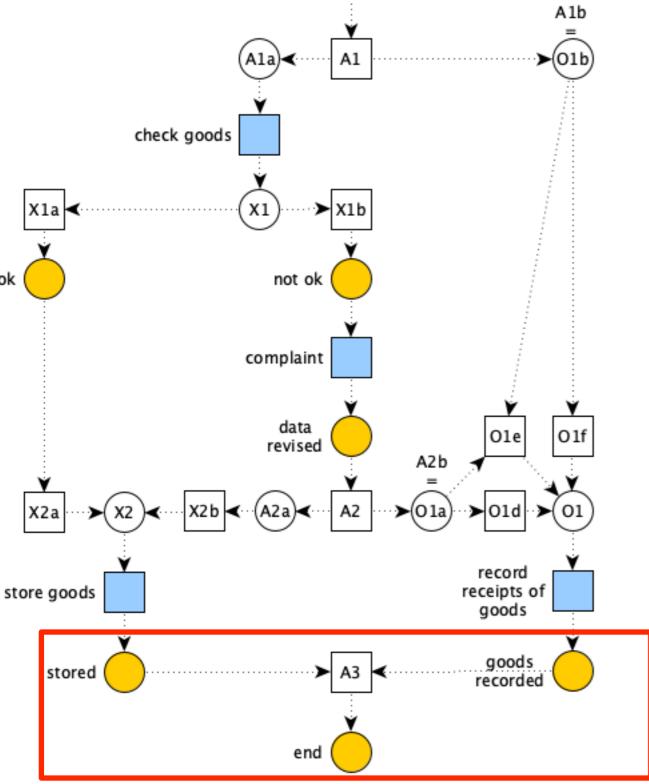
## Example





implicit AND join (because of A2)

## Example Ala Alb Olb



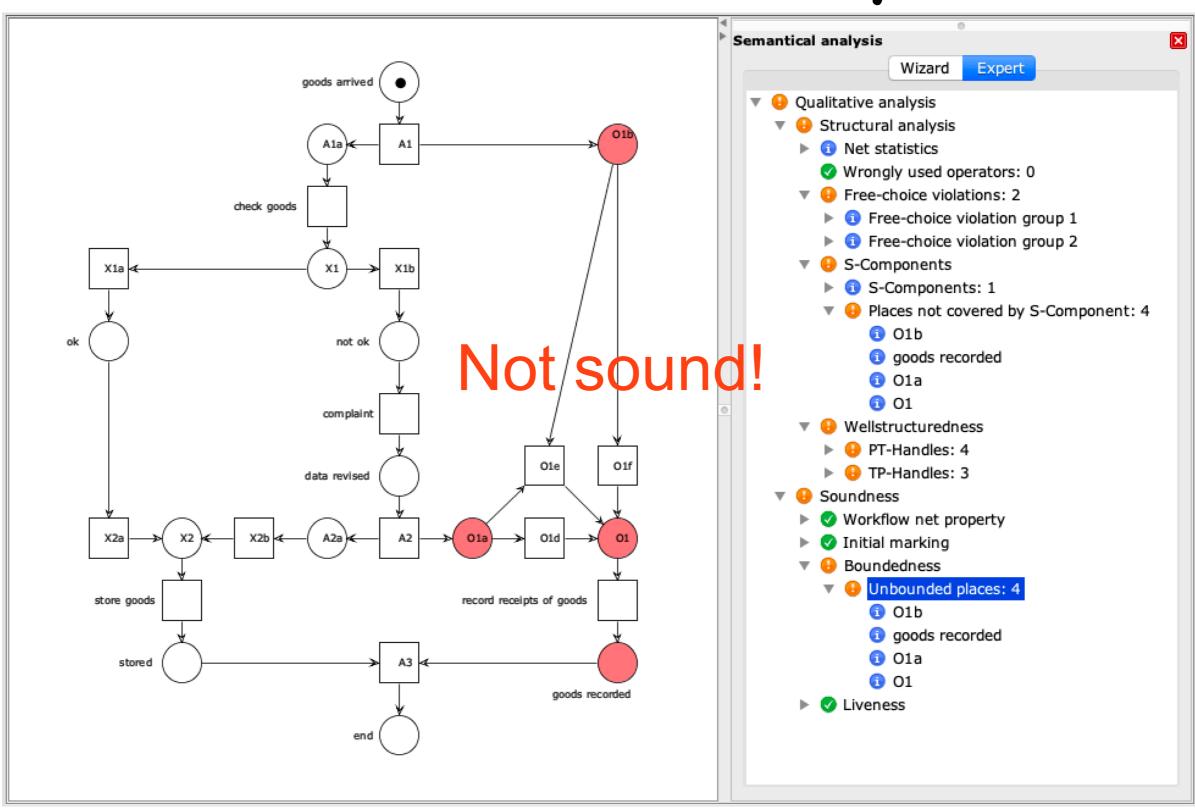


unique end

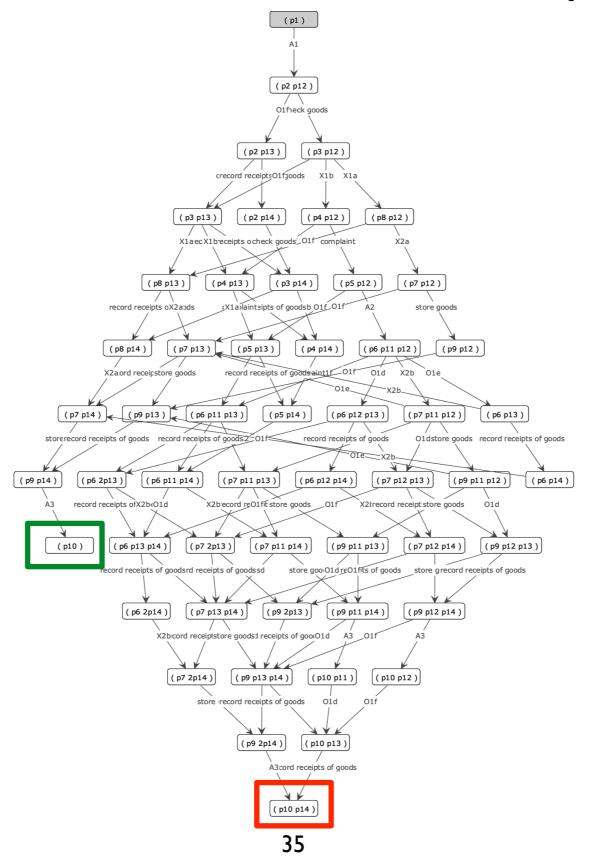
implicit AND join (because of A2)

#### **EPC** wf net Example goods goods arrived arrived A1b check goods check goods X1a ◀ Sound? ... not ok not ok complaint OR complaint 01f O1e **Steps** revised A2b data 1+2+3 record receipts of store goods goods record receipts of store goods goods goods recorded goods stored

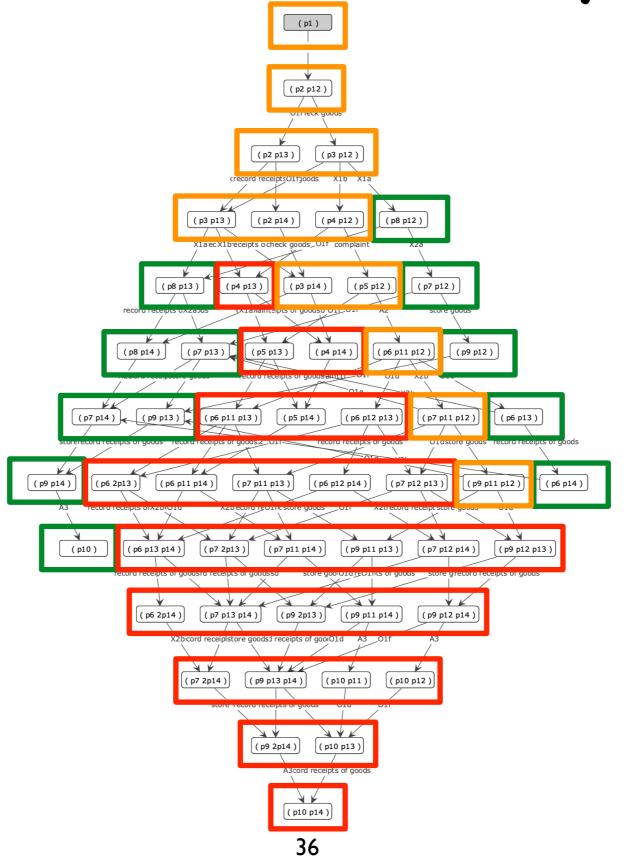
## Soundness analysis

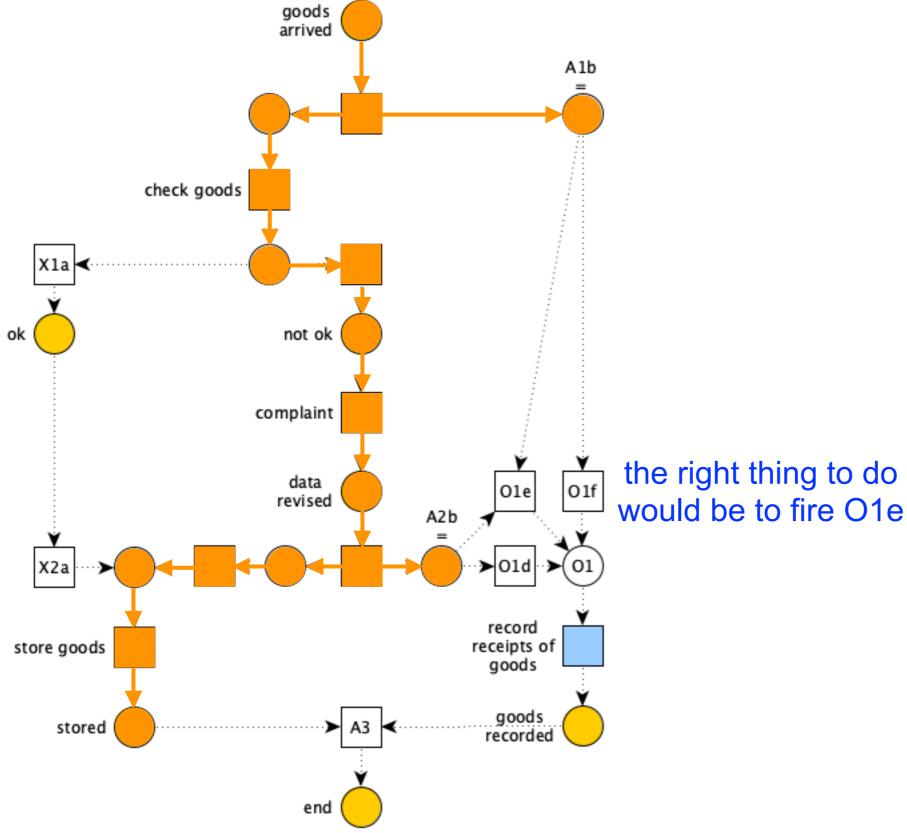


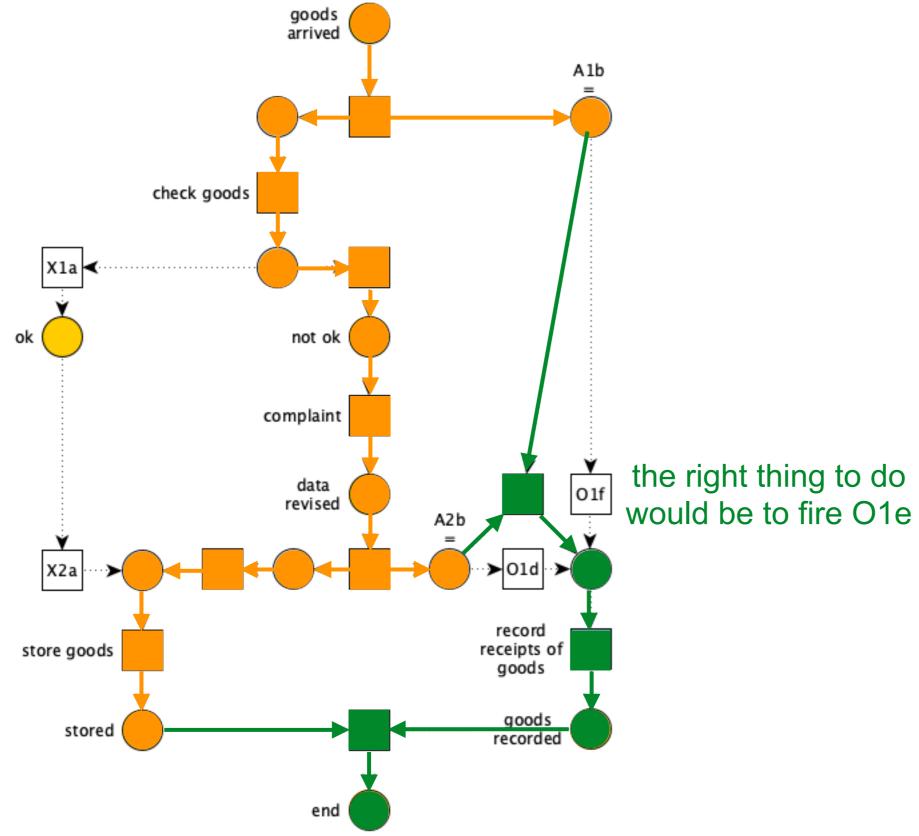
### Soundness analysis

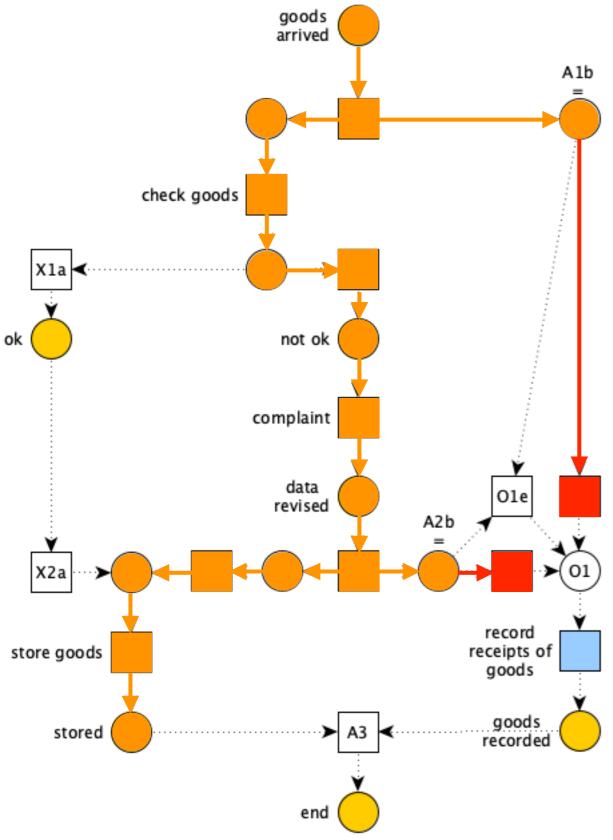


#### Soundness analysis

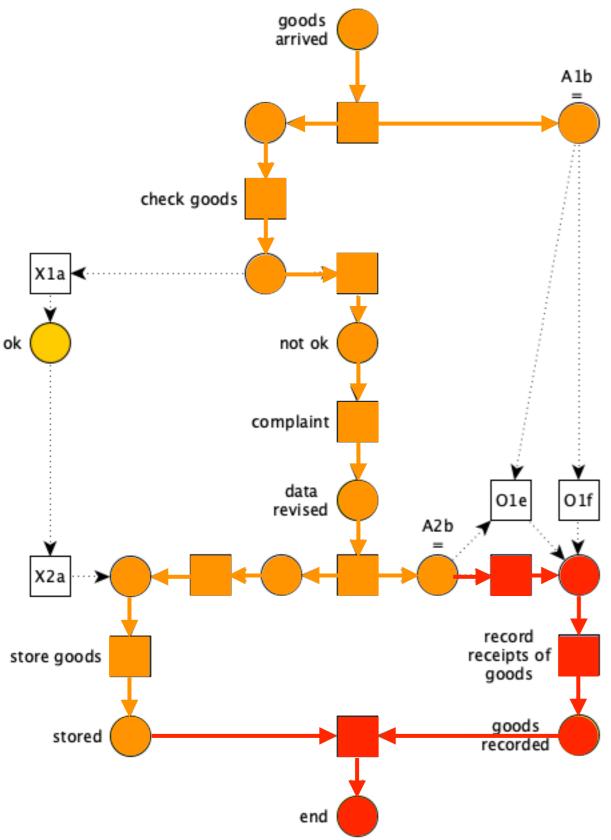




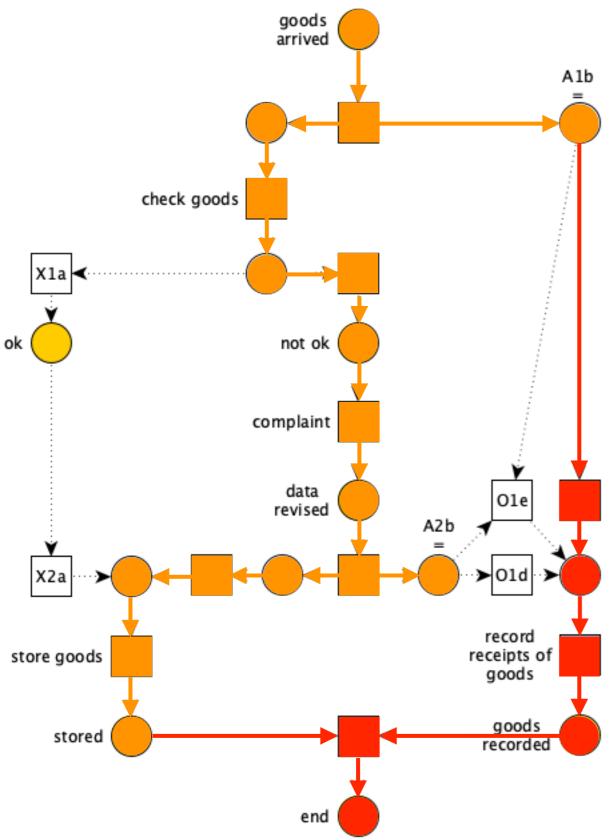




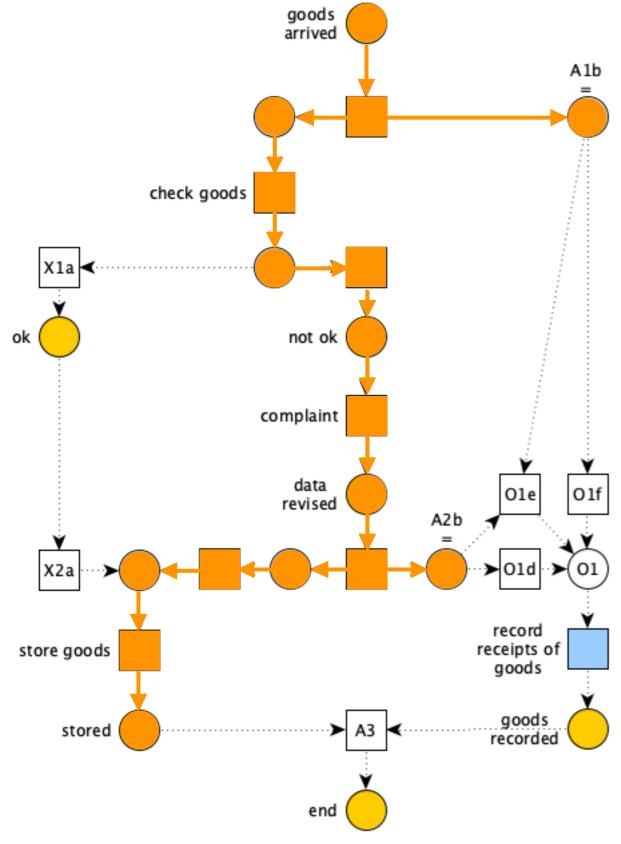
but O1f and O1d are enabled as well (OR semantics!)



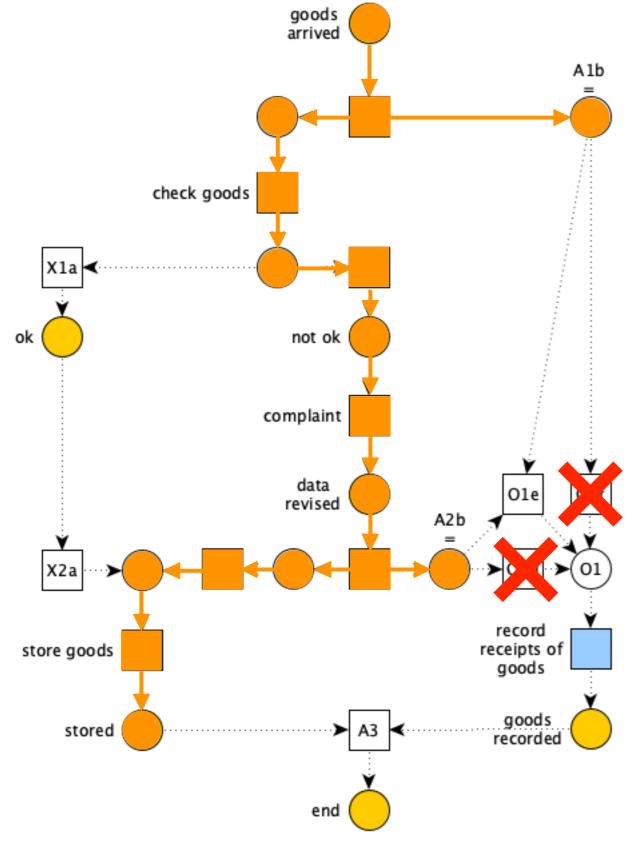
proper completion is not guaranteed (N\* unbounded)

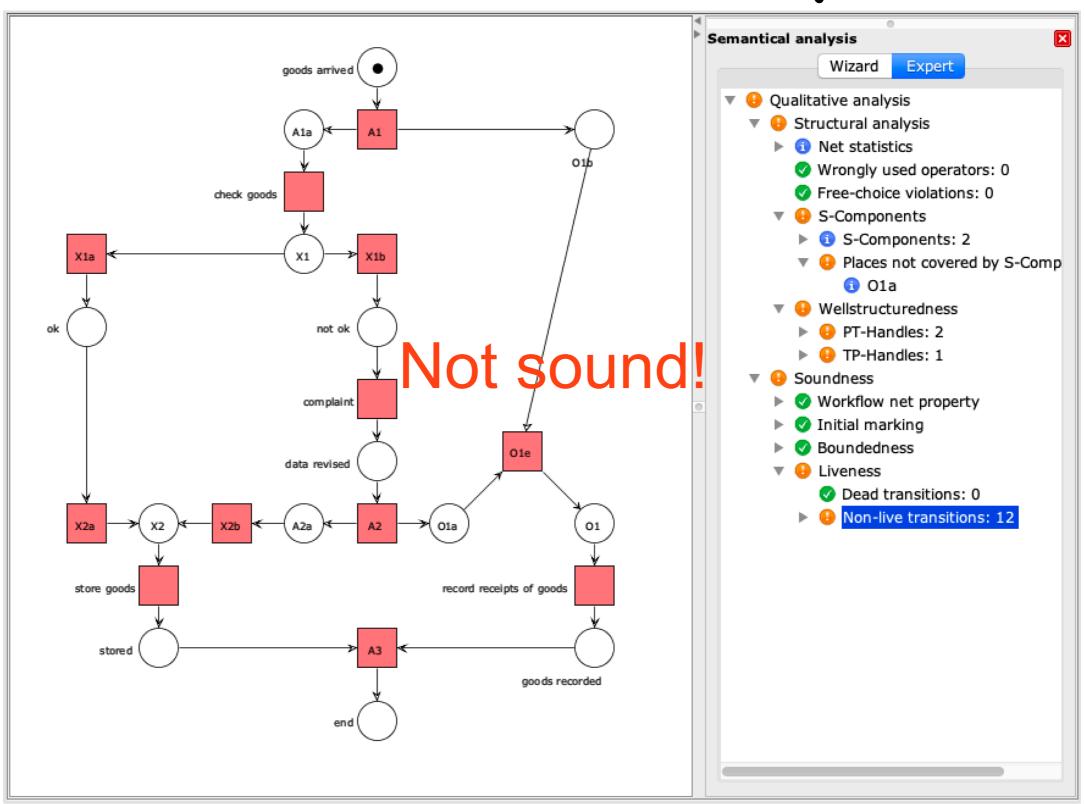


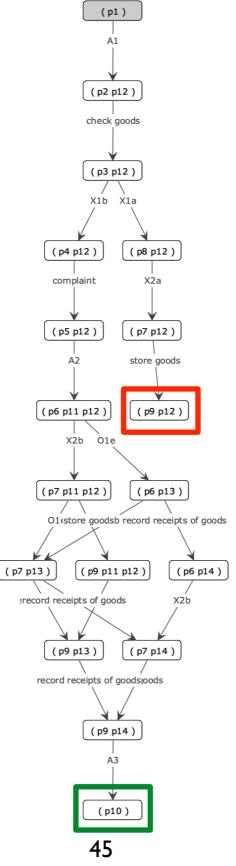
proper completion is not guaranteed (N\* unbounded)

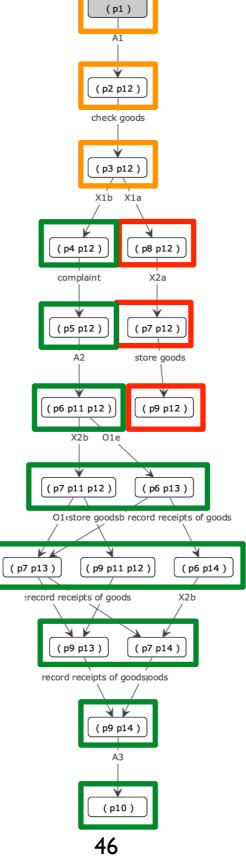


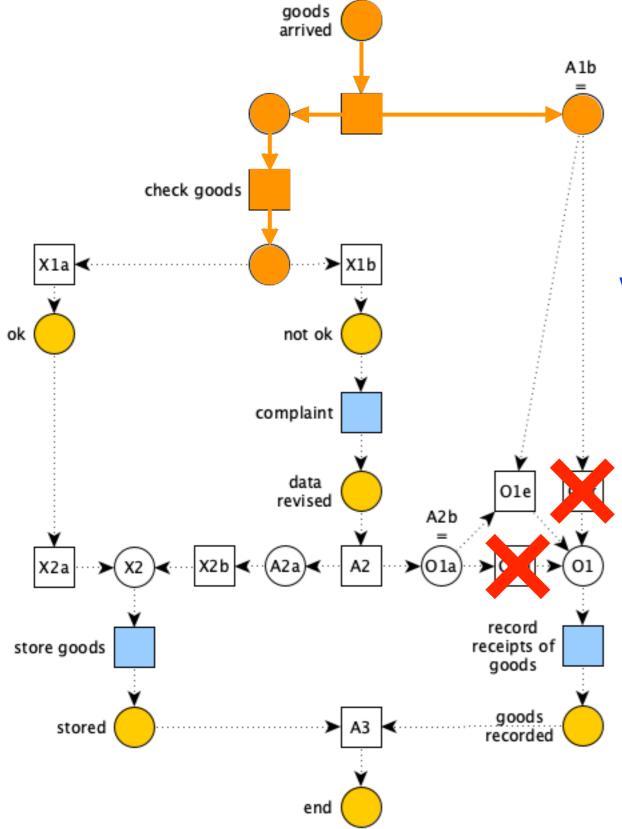
### Can we repair the model?



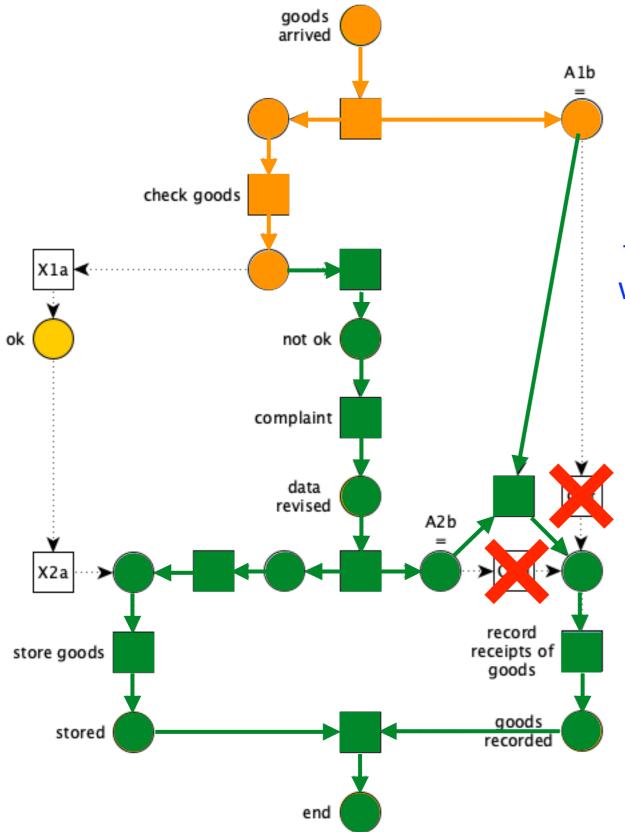




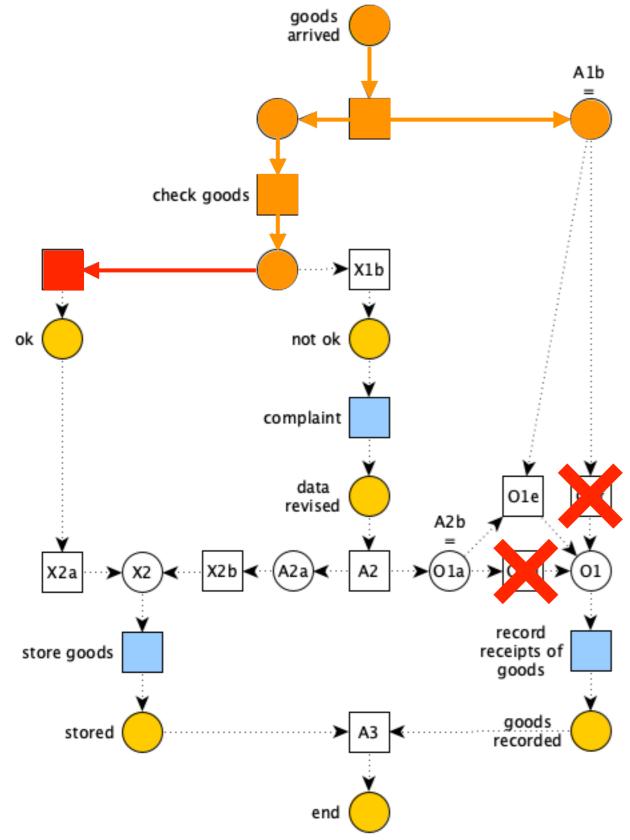




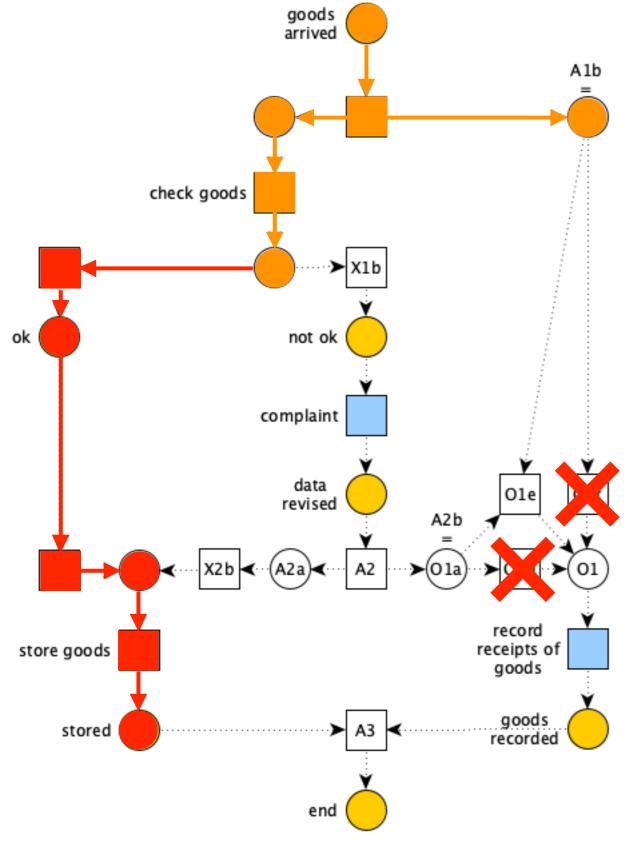
the right thing to do would be to fire X1b



the right thing to do would be to fire X1b

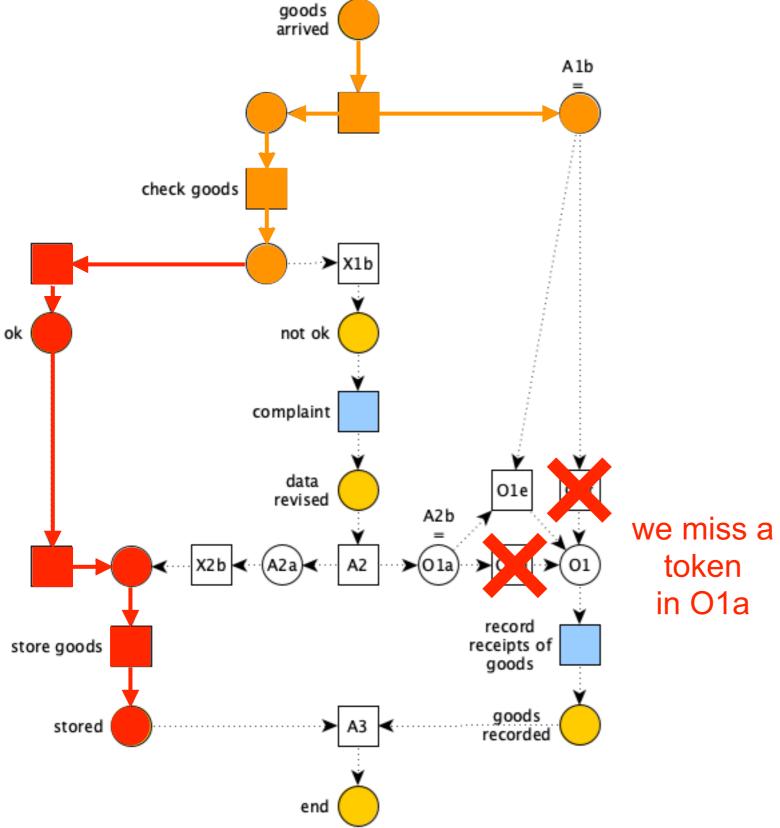


but X1a is enabled as well



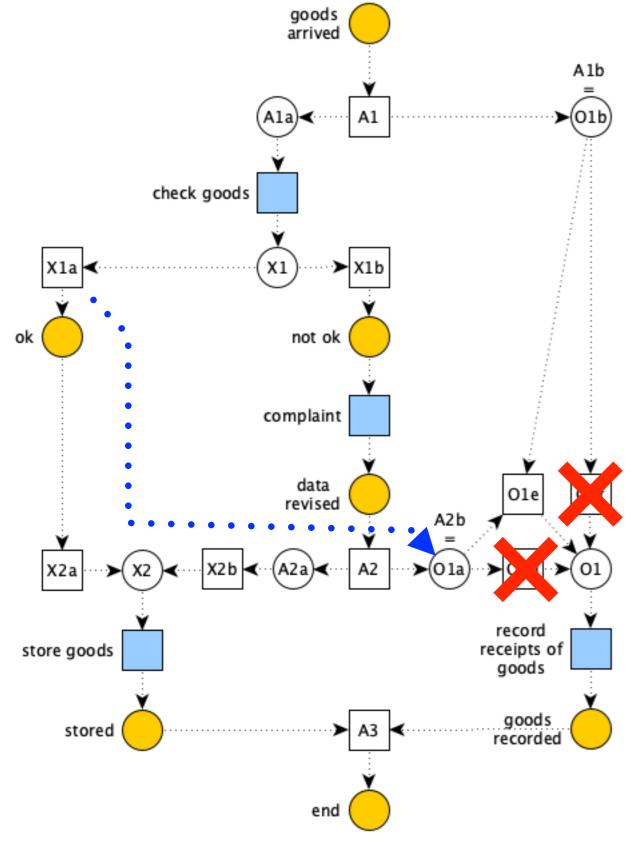
## AND join instead of OR join?

possible deadlock! option to complete is not guaranteed (N\* non-live)

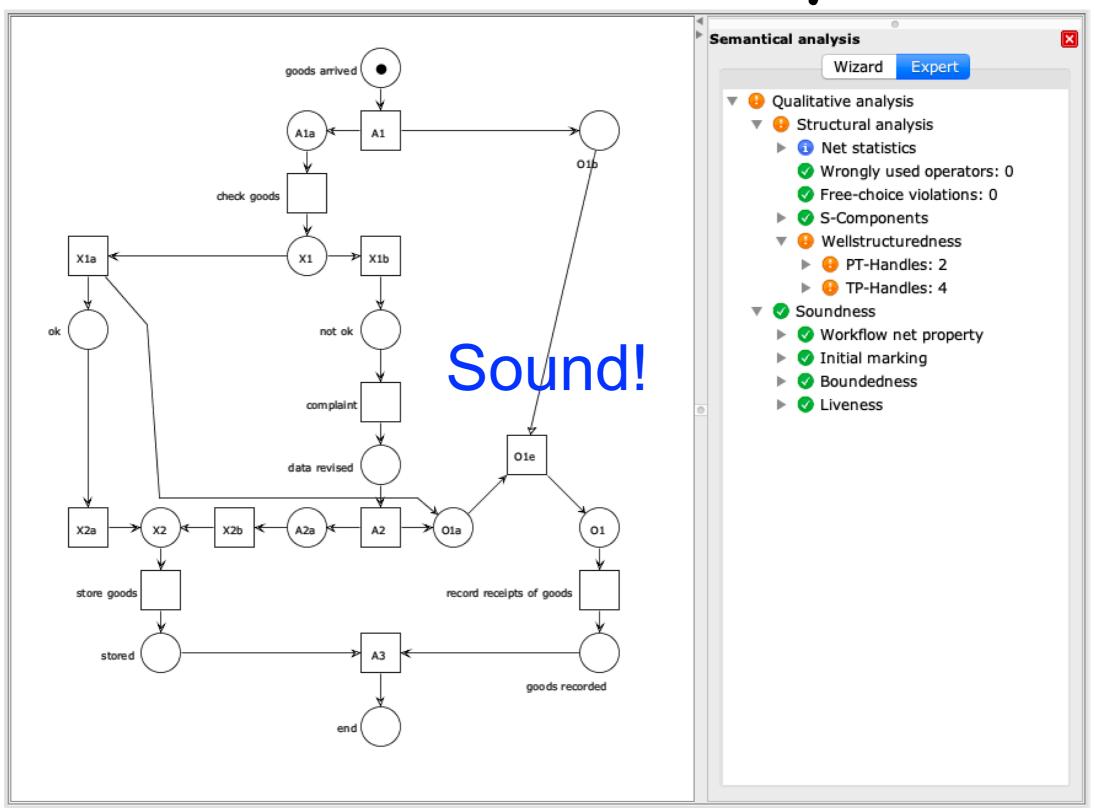


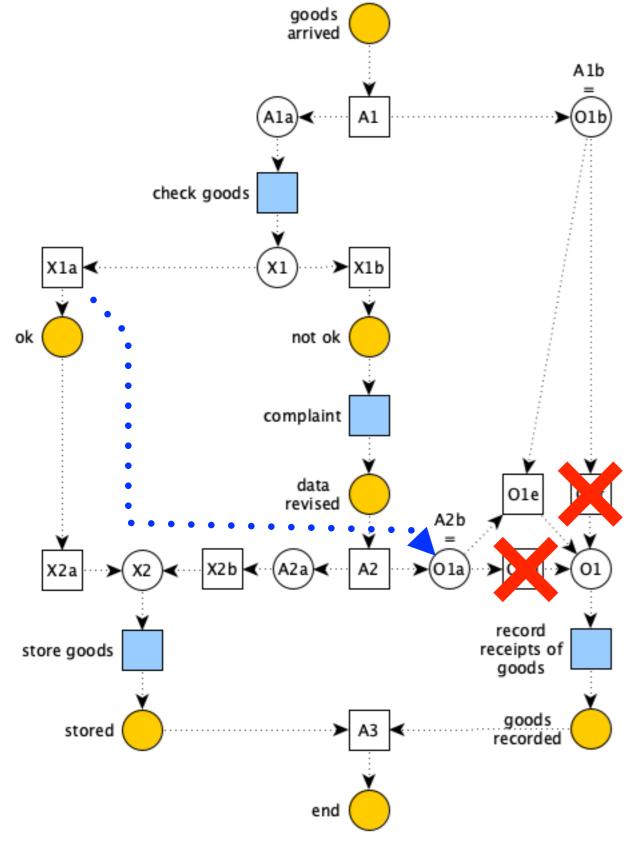
AND join instead of OR join

+ ad hoc flow?

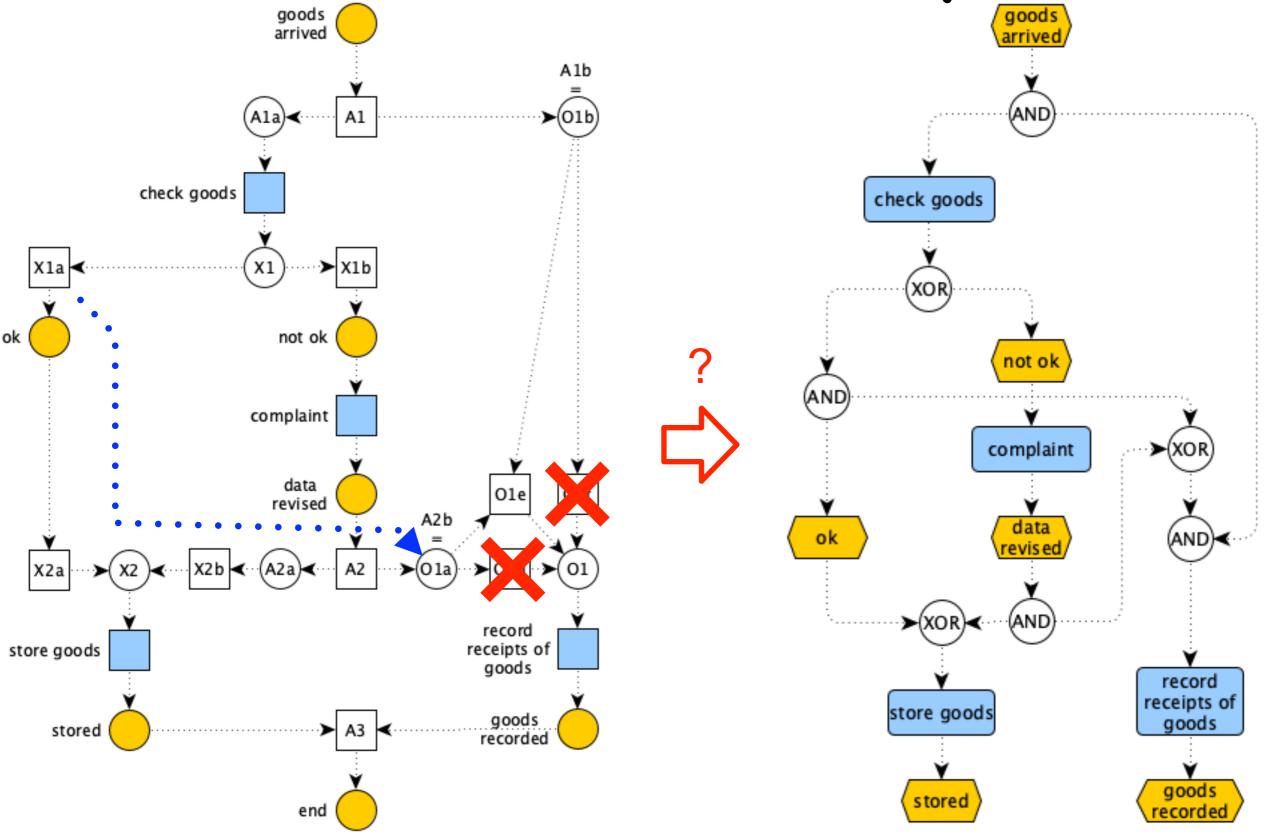


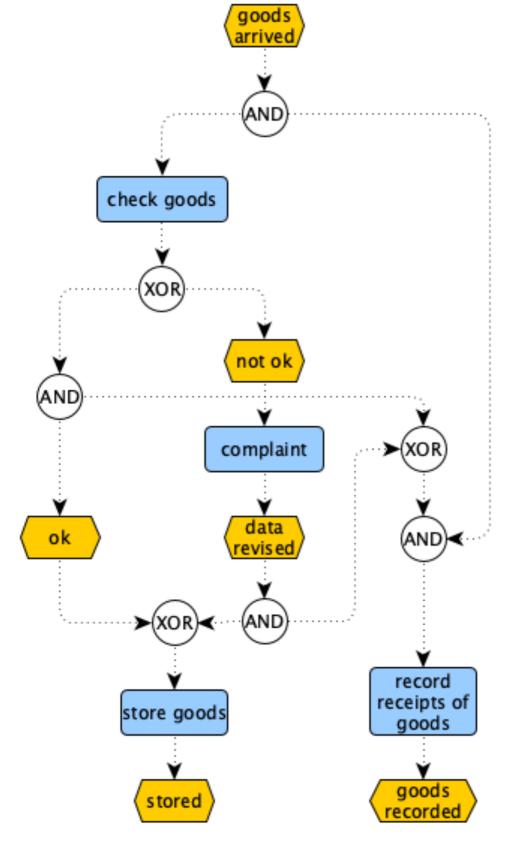
AND join instead of OR join + ad hoc flow?





Sound, but...
we have repaired the wf net,
not the original EPC diagram!





The diagram is now more complex and less readable than the original one!

Are we sure that its translation is the same sound wf net that we have designed ad hoc?

Are we sure it is sound?

Need to restart the analysis!!

# Relaxed Soundness (optional reading)

#### Problem

EPC is widely adopted also at early stages of design

WF nets offer a useful tool

but

Soundness can be too demanding at early stages

#### (Un)sound behaviours

#### A **sound** behaviour:

we move from a start event to an end event so that nothing blocks or remains undone

The language of the net collects all and only its sound behaviours

$$L(N) = \{ \sigma \mid i \xrightarrow{\sigma} o \}$$

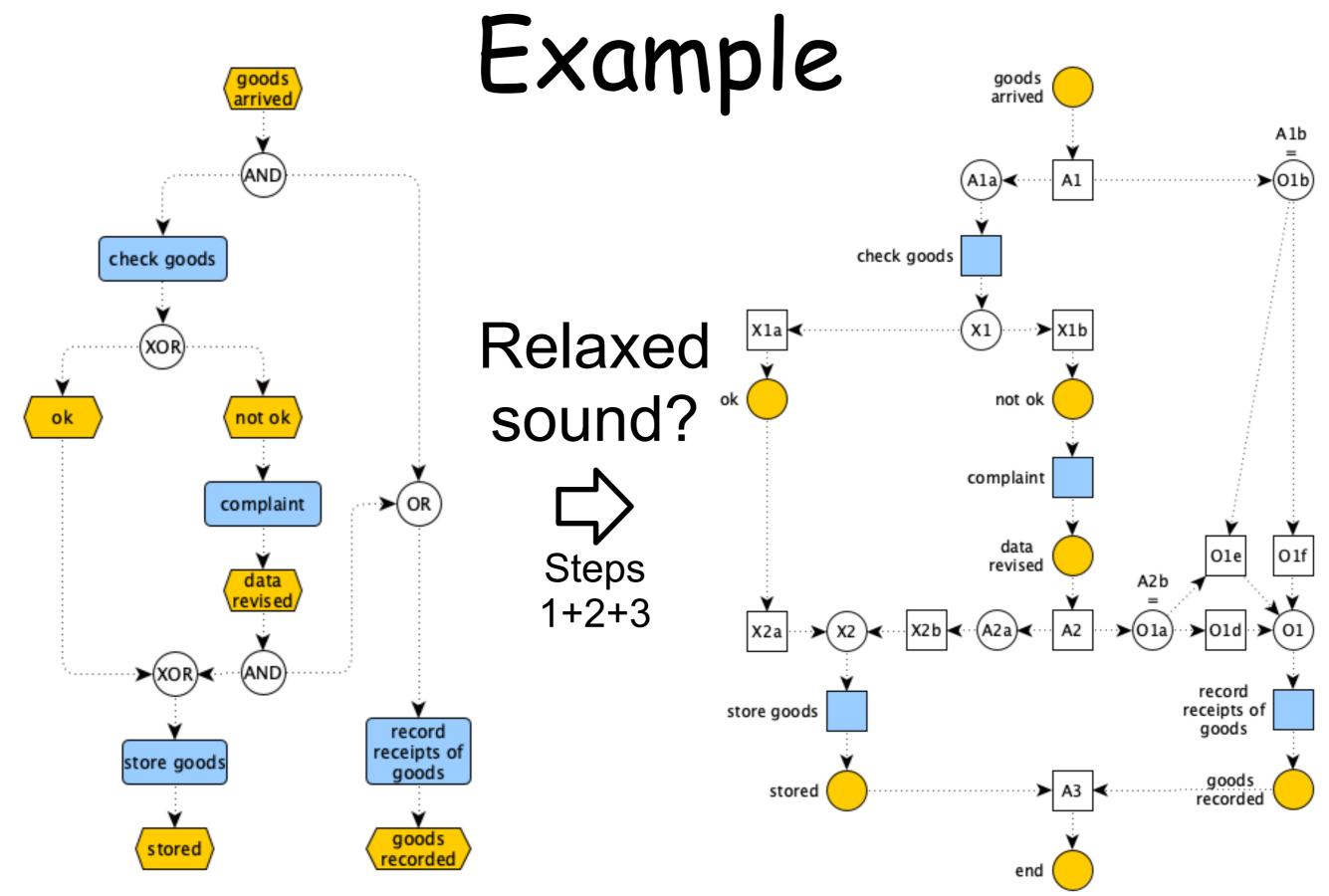
Execution paths leading to **unsound** behaviours can be used to infer potential mistakes

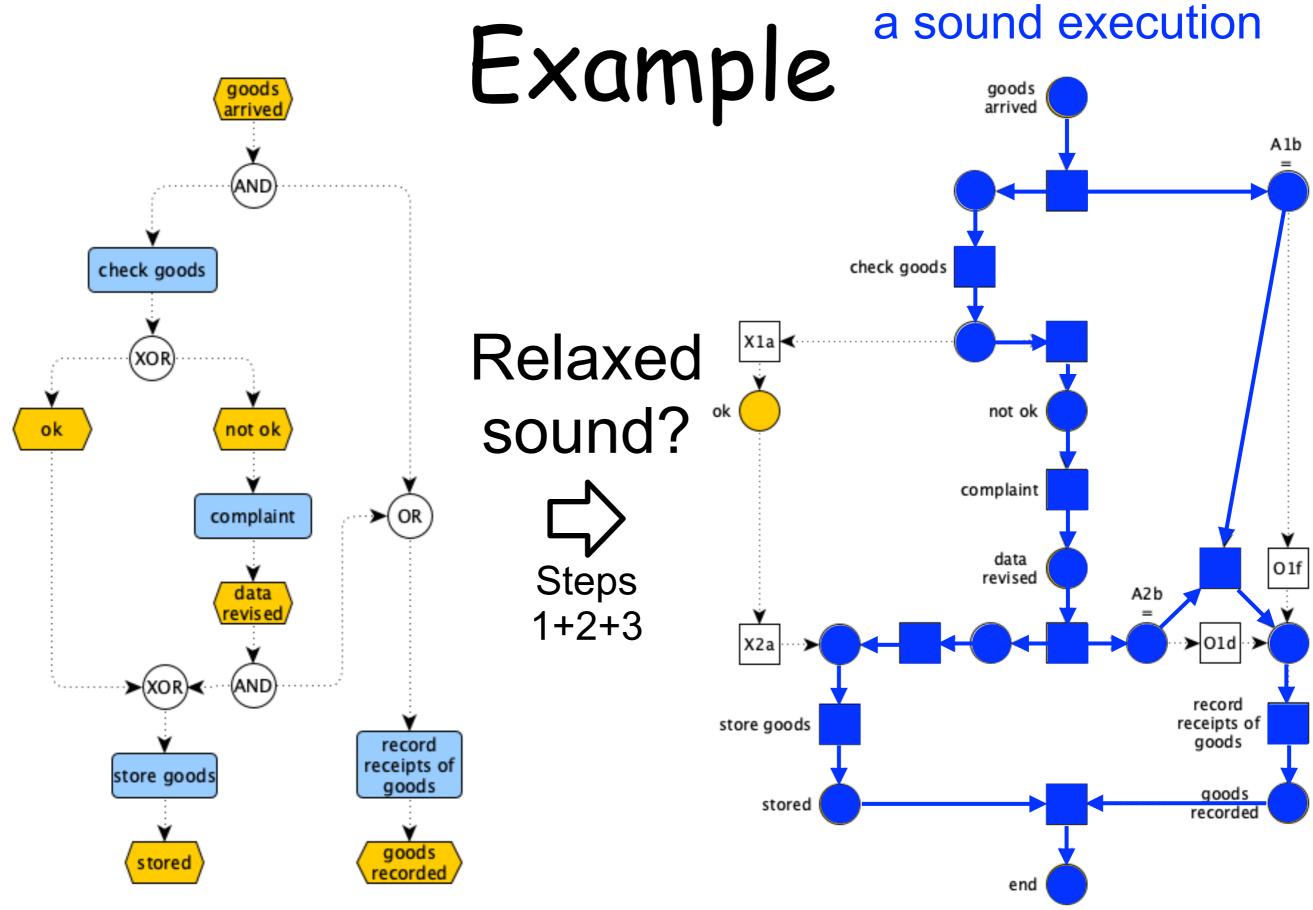
#### Relaxed soundness

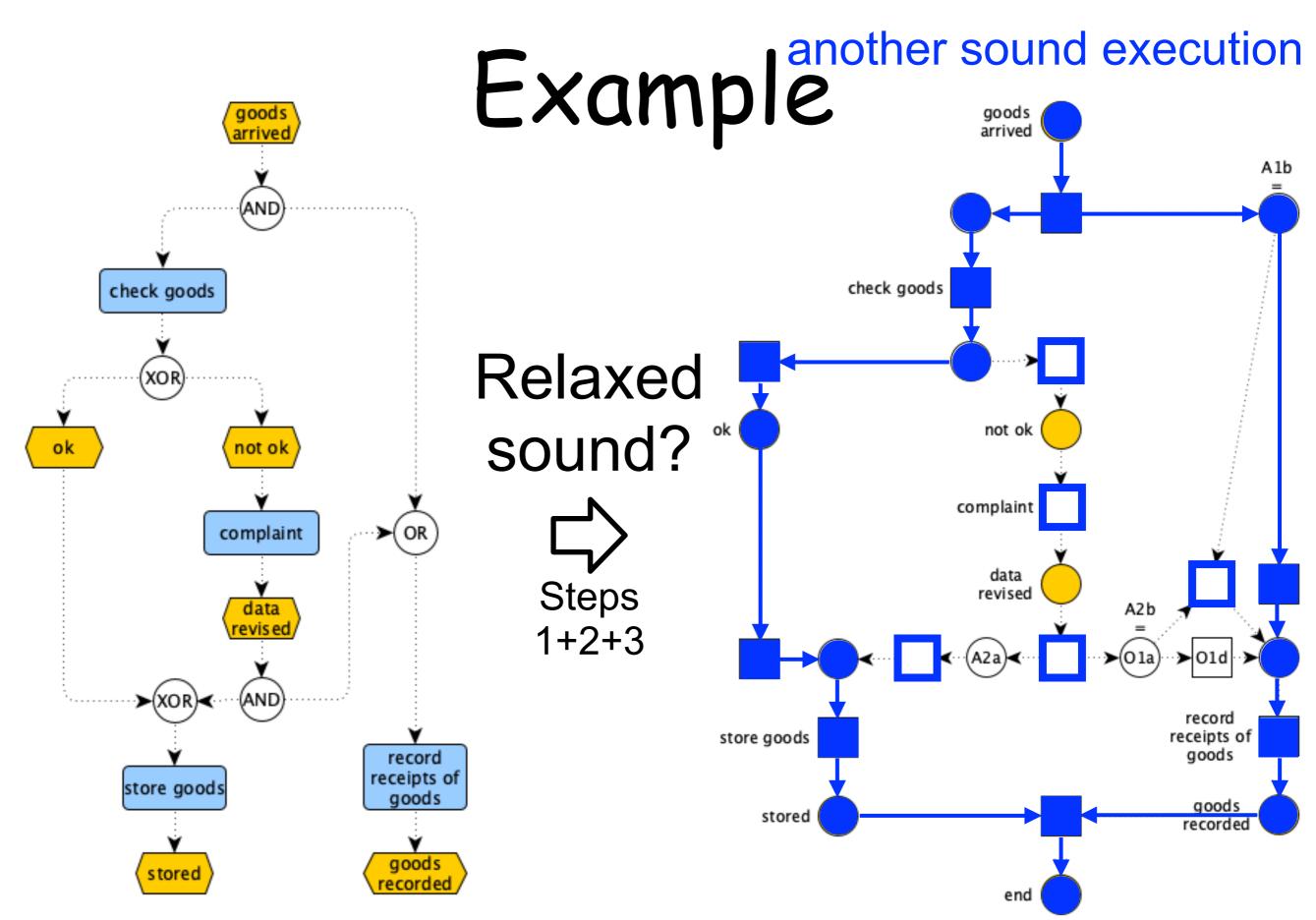
If some unsound behaviour is possible but any transition can take part to one sound execution, then the process is called **relaxed sound** 

**Definition**: A WF net is **relaxed sound** if every transition belongs to a firing sequence that starts in state i and ends in state o (i.e. it appears in the language of the net)

$$\forall t \in T. \ \exists \sigma \in L(N). \ \vec{\sigma}(t) > 0$$

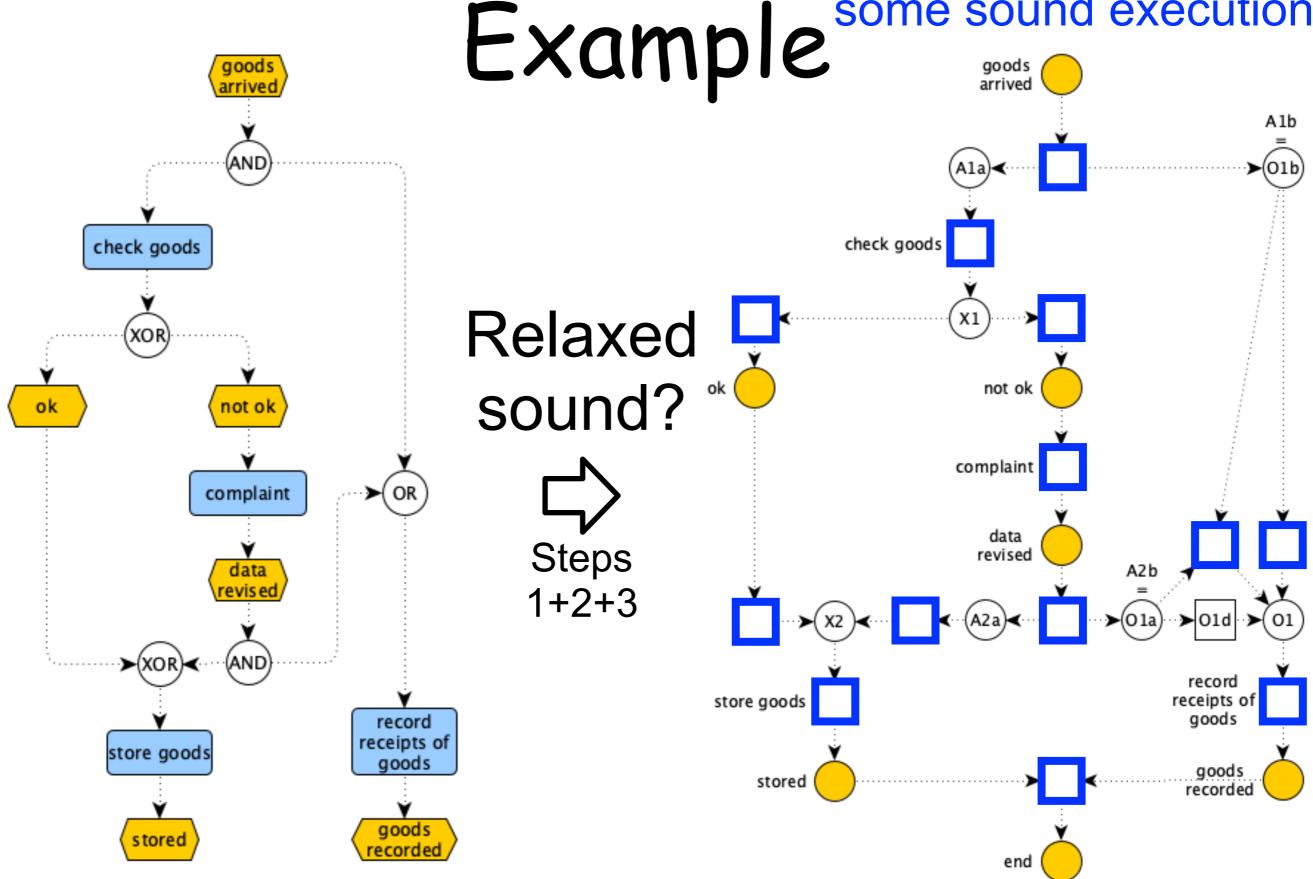




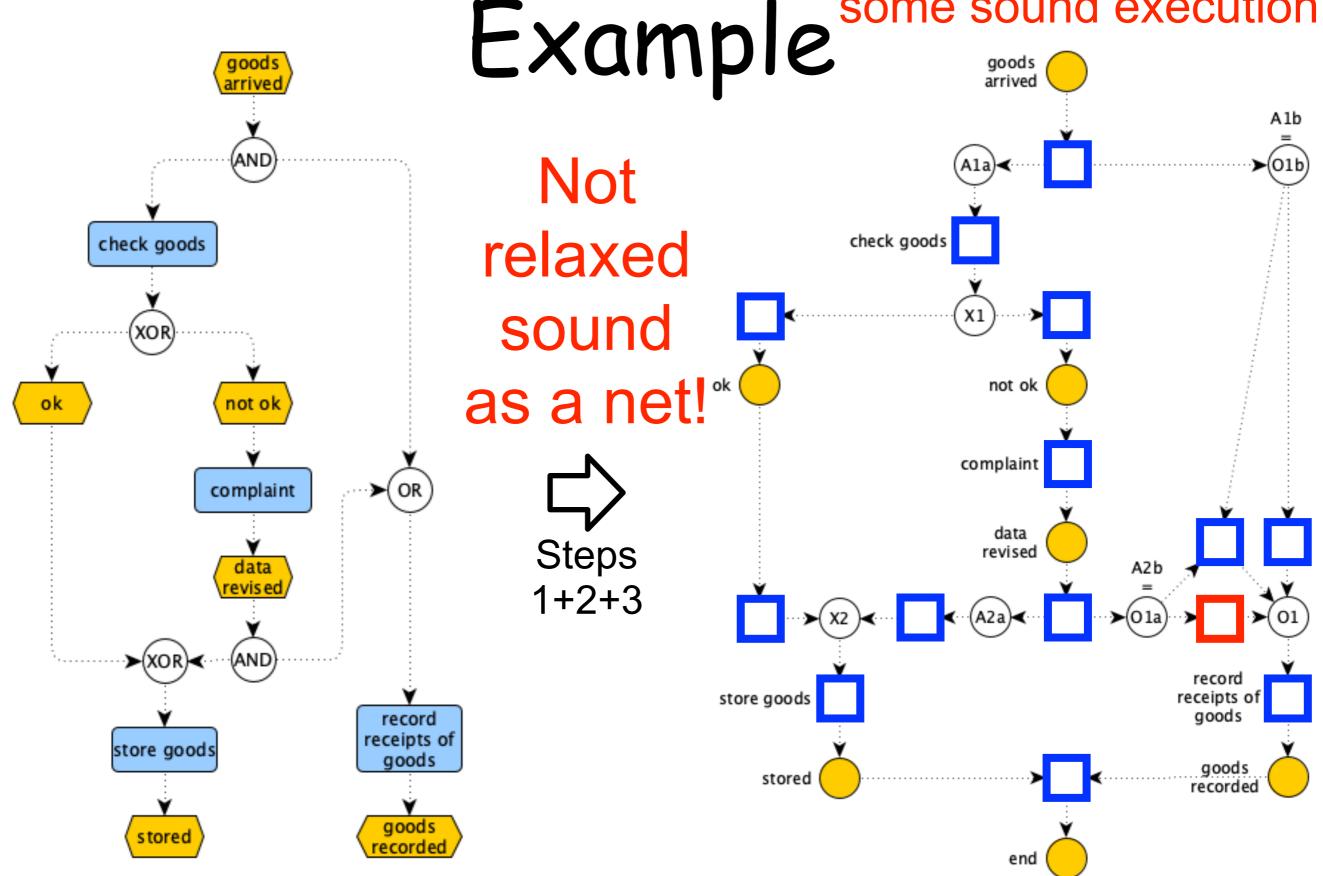


#### tasks involved in

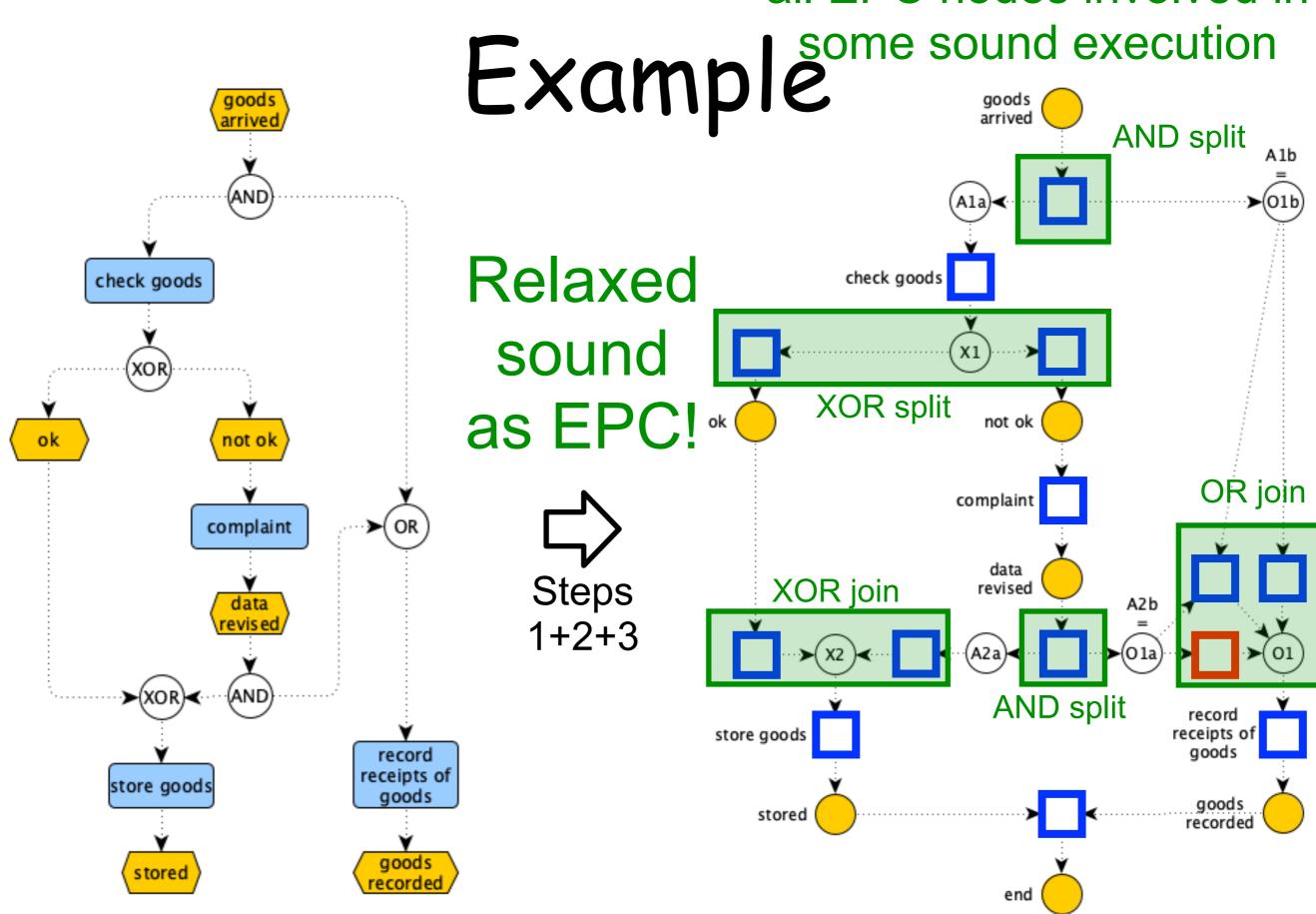
some sound execution



one task not involved in some sound execution



all EPC nodes involved in



#### Relaxed soundness?

If the WF net is **not relaxed sound** there are transitions that are not involved in sound executions (not included in a firing sequence of L(N))

Their EPC counterparts may need improvements

Relaxed soundness can be proven only by enumeration (of enough firing sequences of L(N))

#### Open problem

No equivalent characterization is known that is more convenient to check

# Second attempt (no OR connectors)

## Formalization and Verification of Event-driven Process Chains

W.M.P. van der Aalst

Department of Mathematics and Computing Science, Eindhoven University of Technology, P.O. Box 513, NL-5600 MB, Eindhoven, The Netherlands, telephone: -31 40 2474295, e-mail: wsinwa@win.tue.nl

#### Simplified EPC

We restrict the analysis to a sub-class of EPC diagrams

We require:

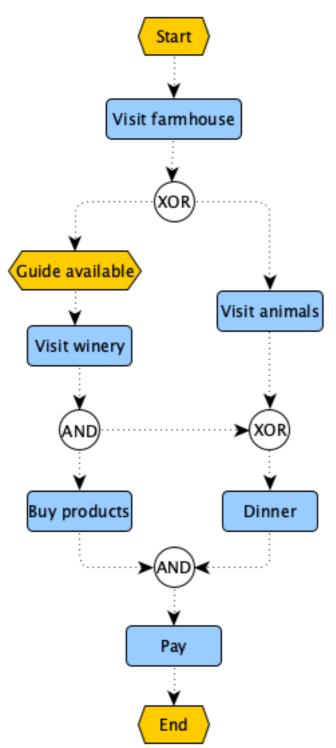
event / function alternation

(also along paths between two connectors) (fusion not needed, dummy places/transitions not needed)

OR-connectors are not present

(avoid intrinsic problems with OR join)

## OR-connectors are not present alternation is not satisfied

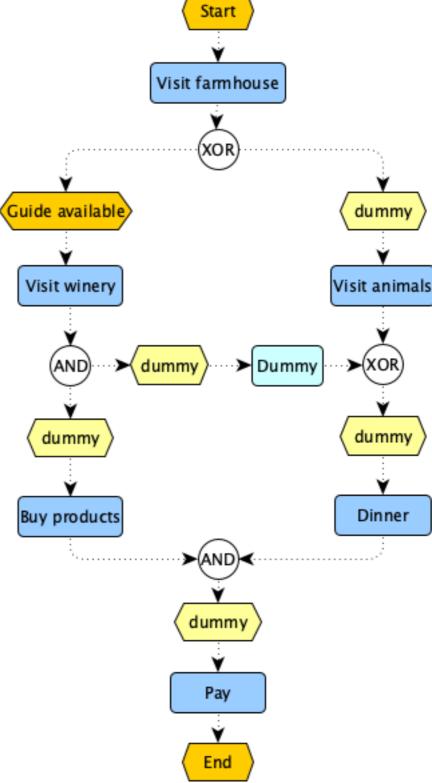


Example

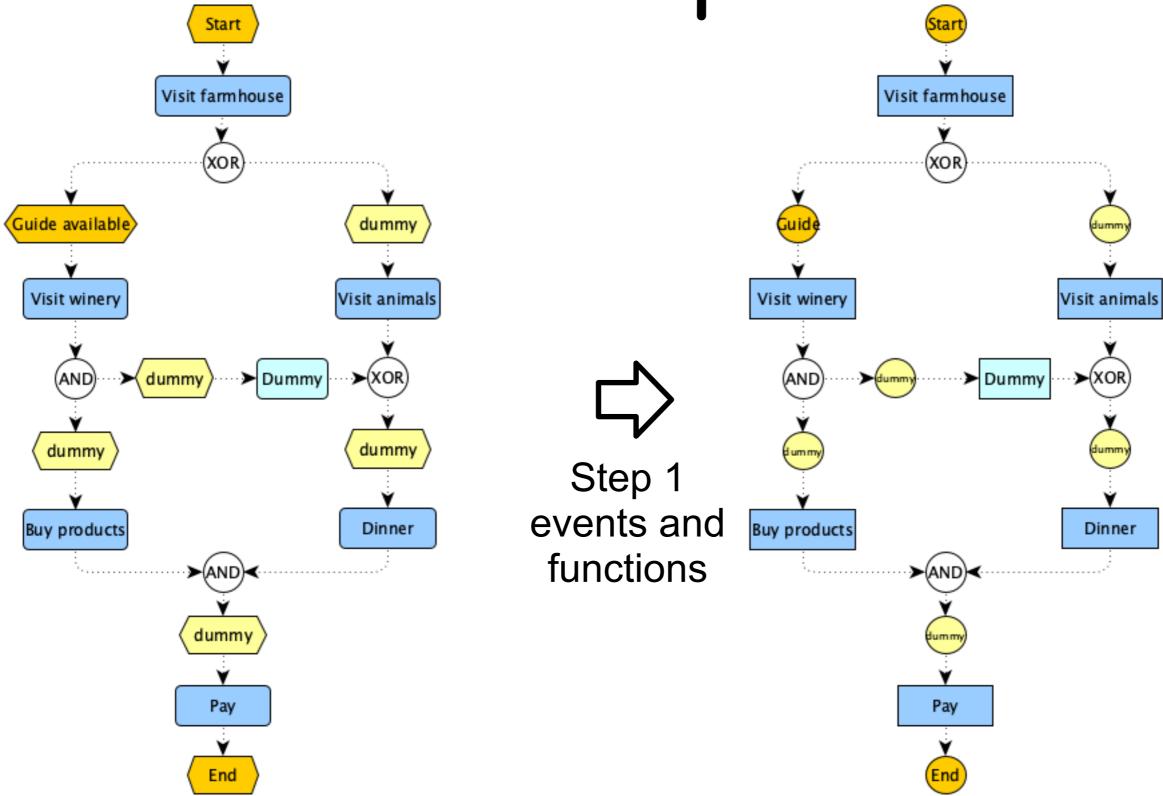
Add dummy events and functions to force alternation



Step 0



Example



# Step 1: split/join connectors

The translation of logical connectors depends on the context:

if a connector connects **functions to events** we apply a certain translation

if it connects **events to functions** we apply a different translation

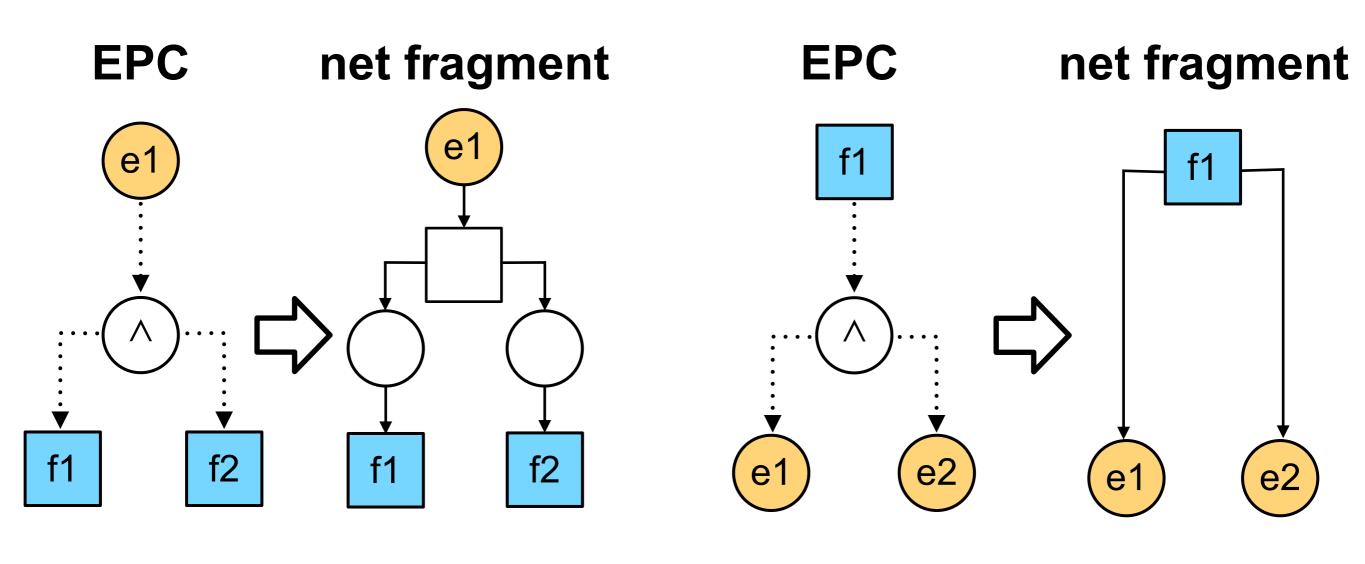
# Step 1: split/join connectors

The translation of logical connectors depends on the context:

if a connector connects transitions to places we apply a certain translation

if it connects places to transitions we apply a different translation

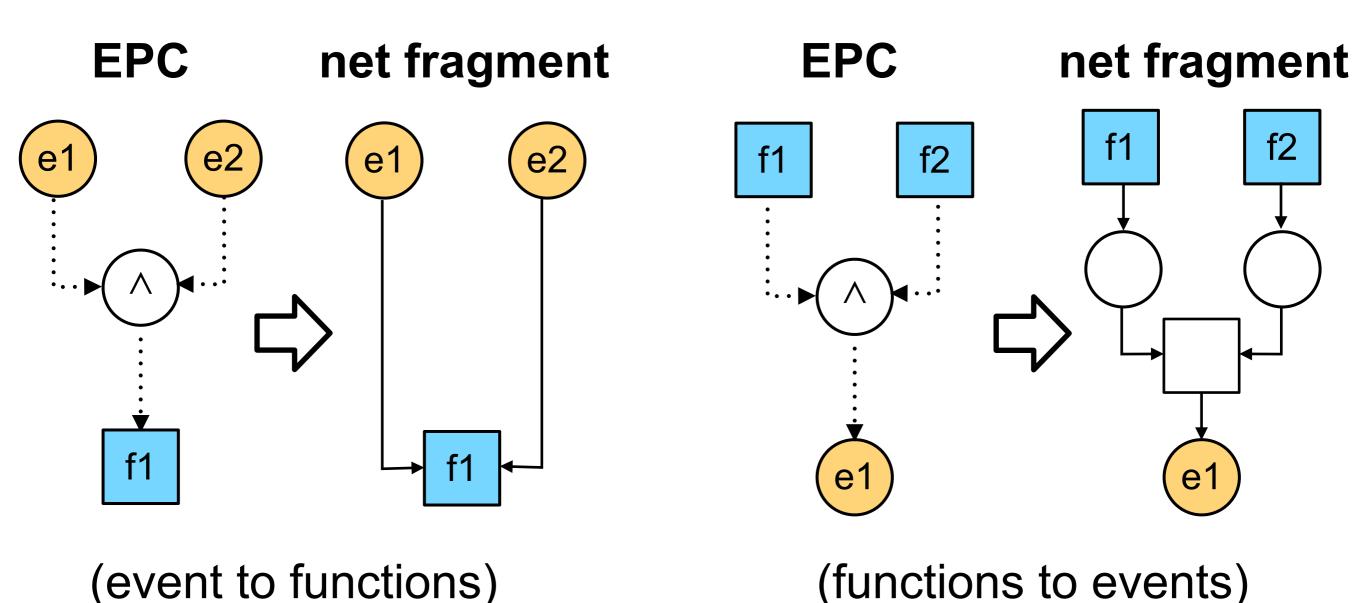
#### Step 1: AND split

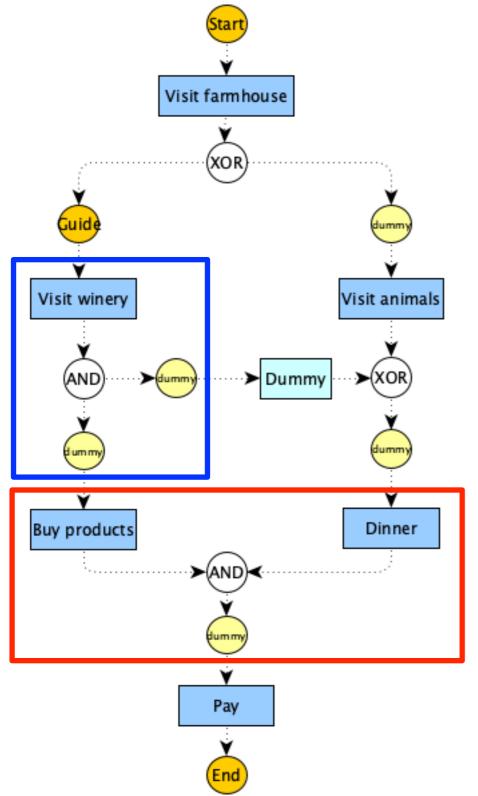


(event to functions)

(functions to events)

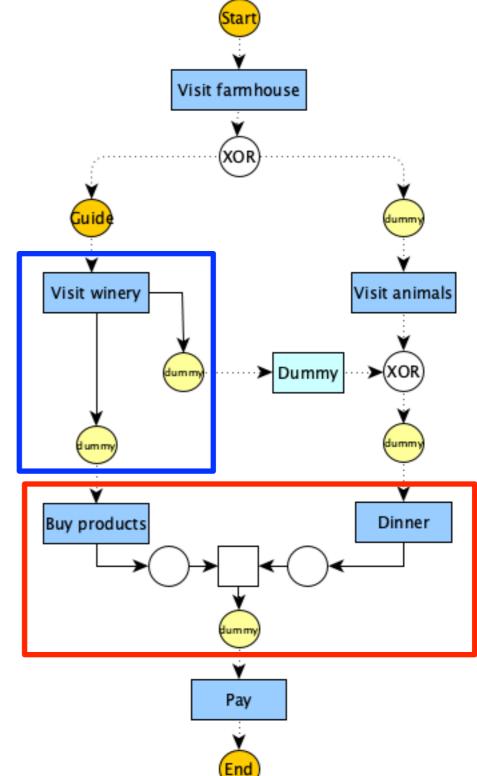
#### Step 1: AND join



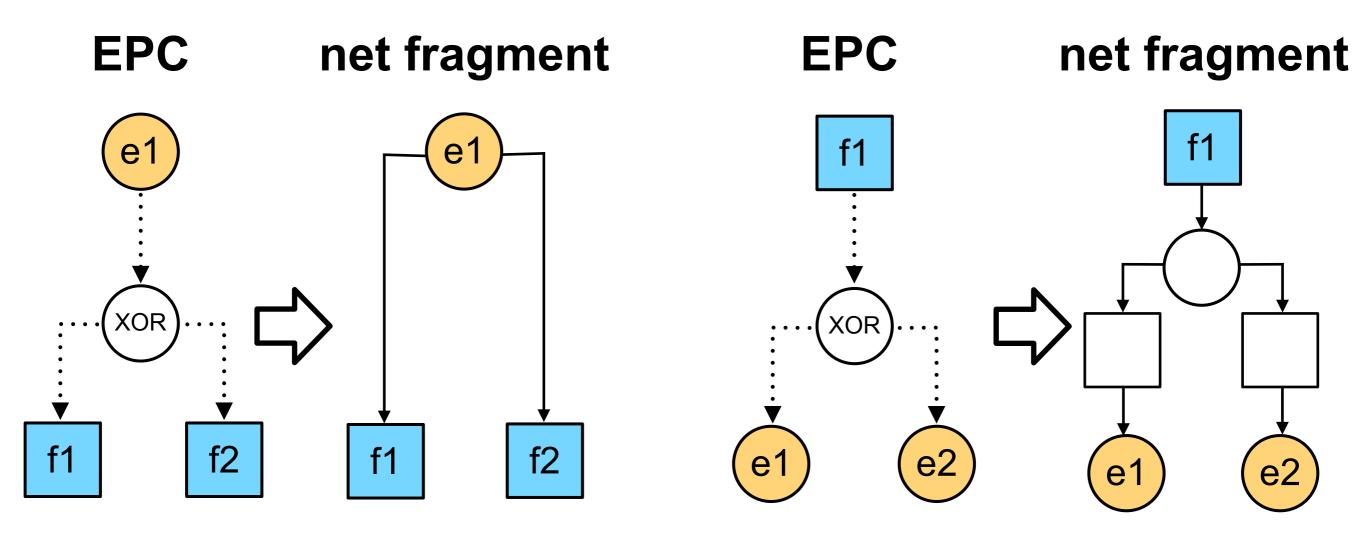




AND connectors



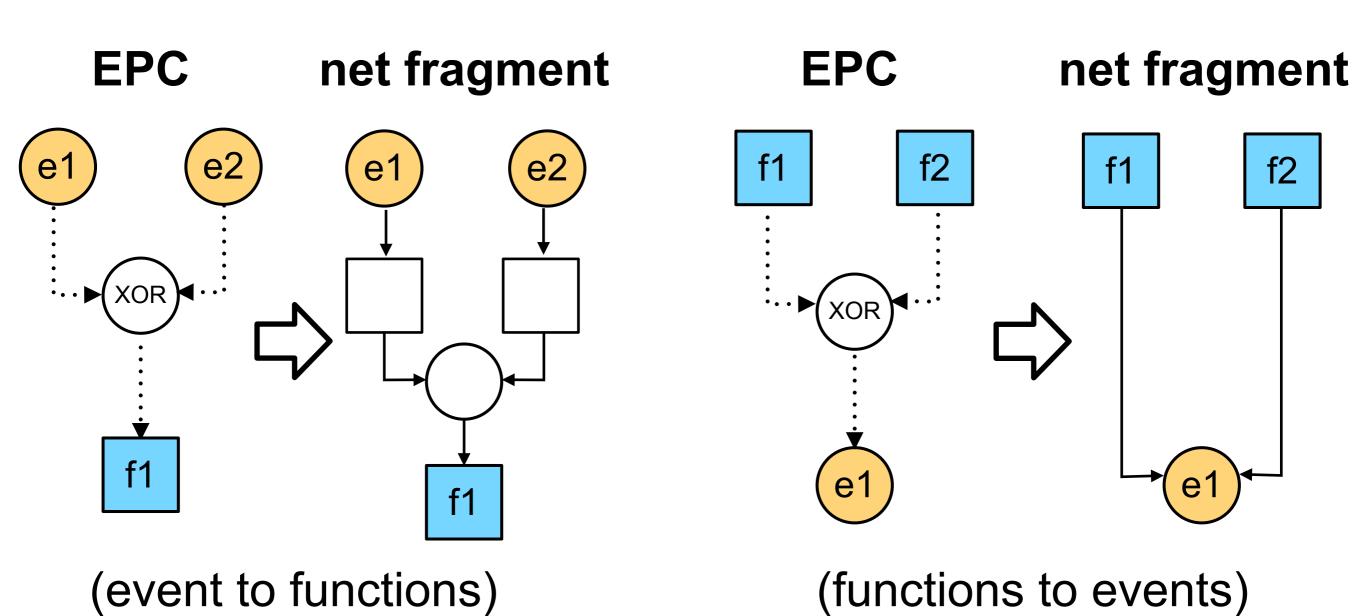
#### Step 1: XOR split

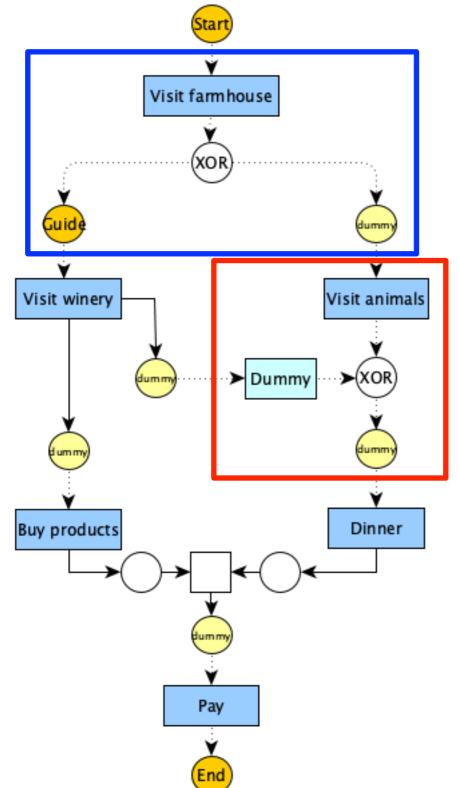


(event to functions)

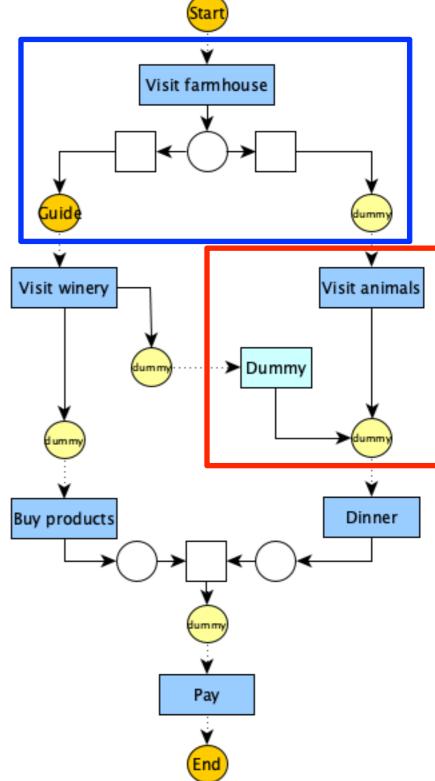
(functions to events)

#### Step 1: XOR join

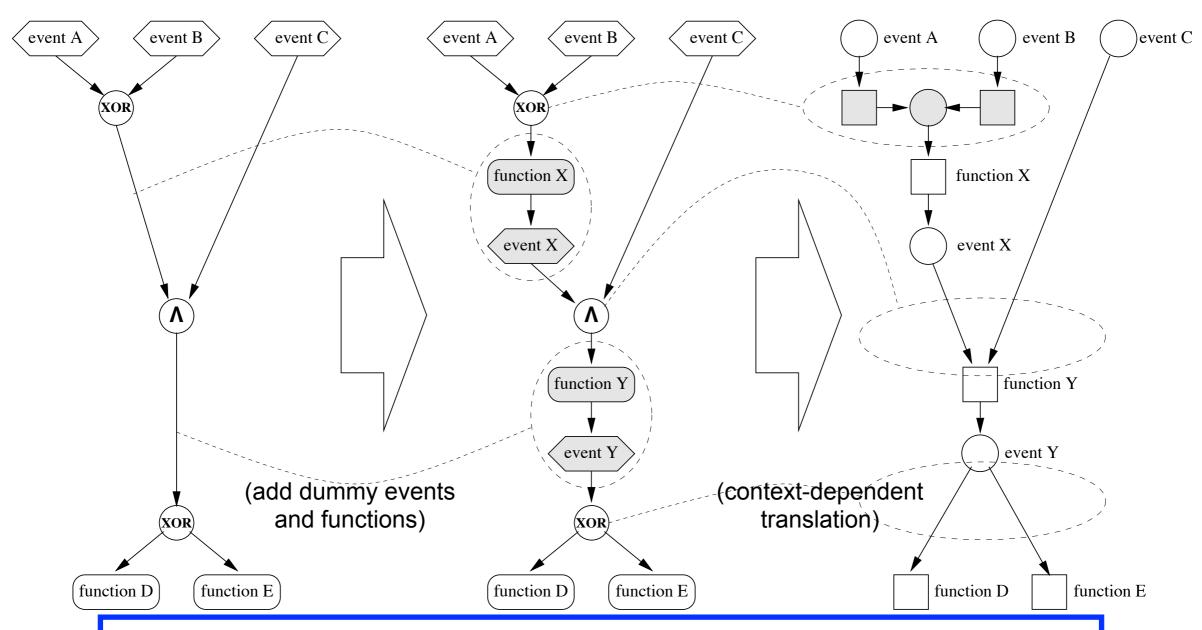




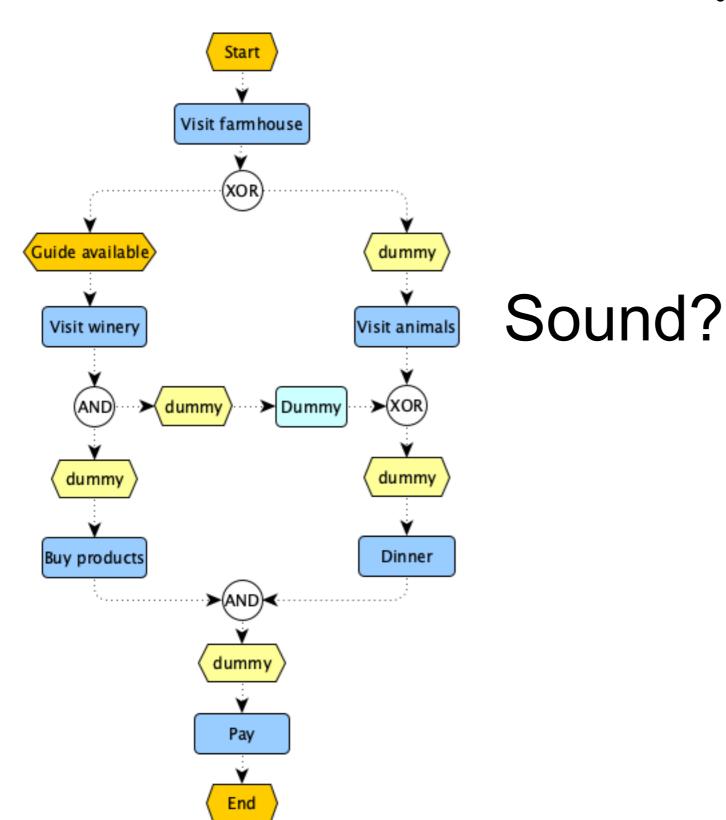


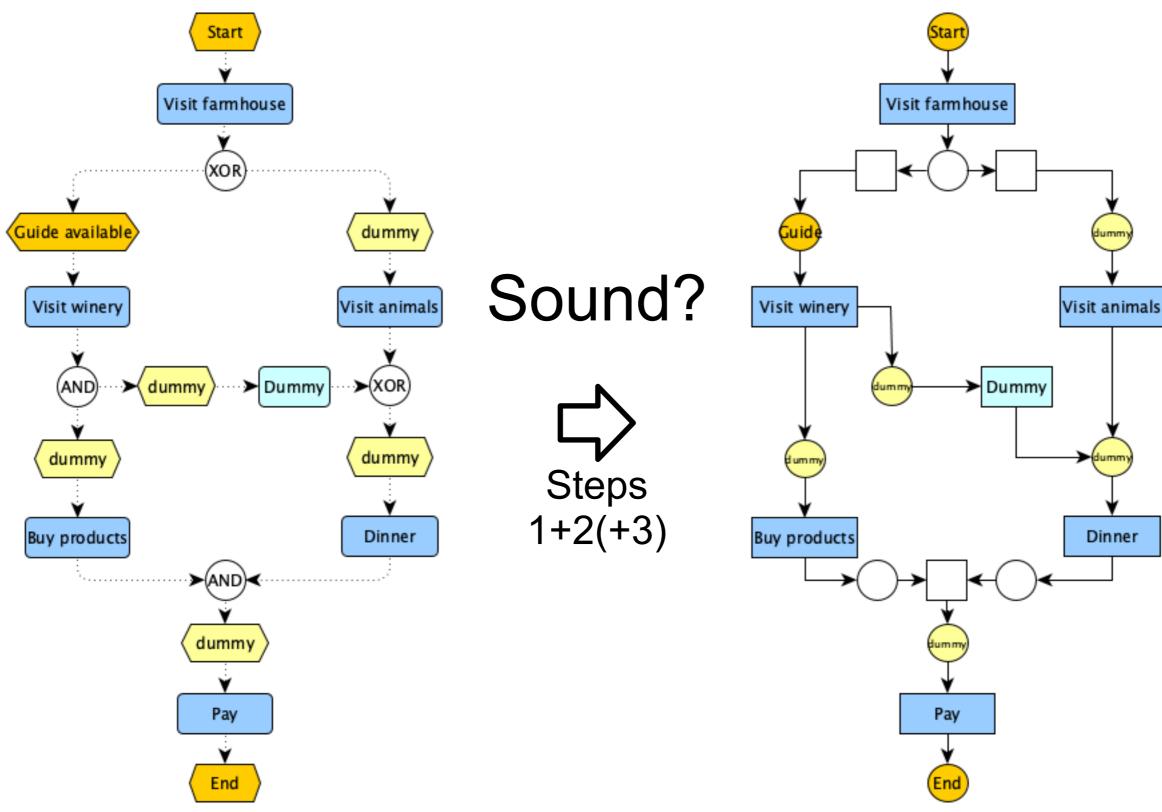


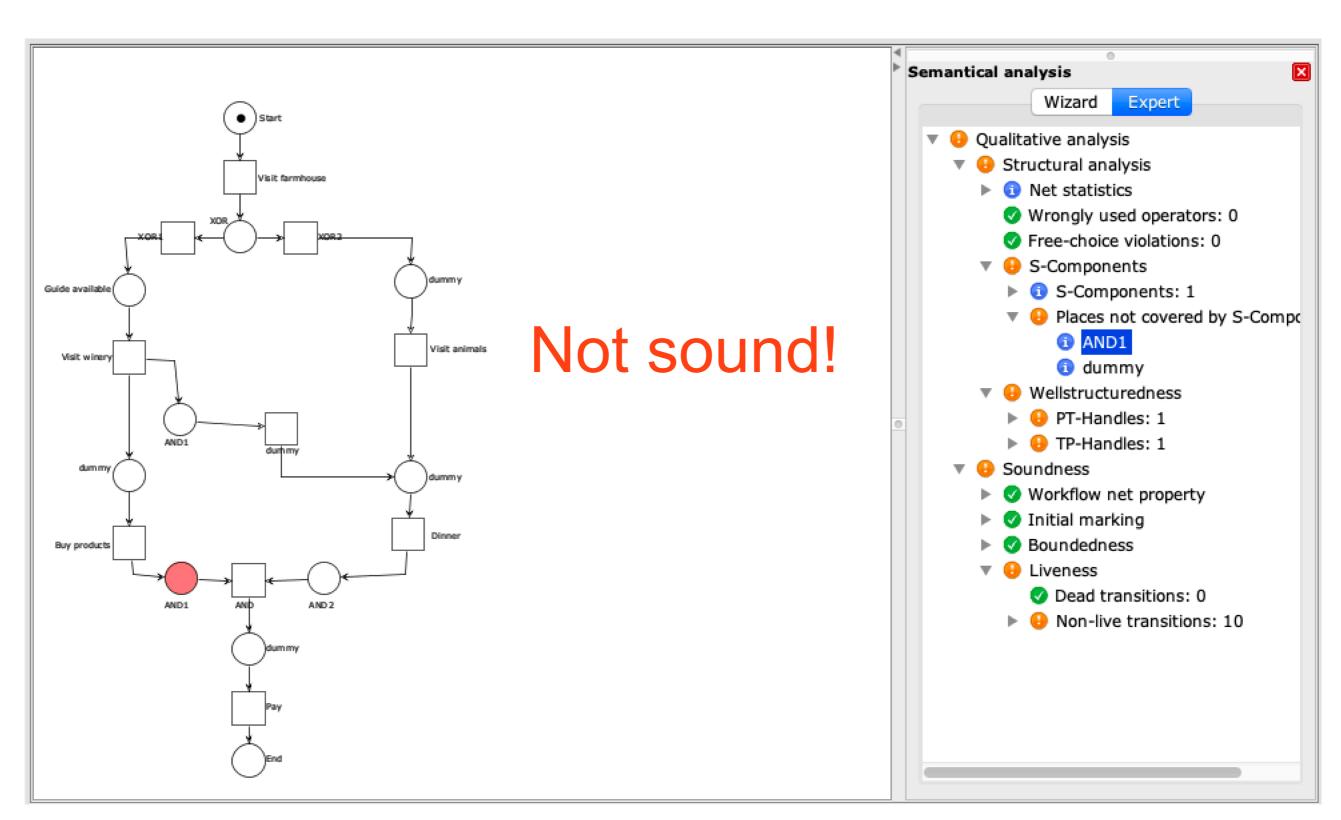
#### Overall strategy



From any EPC we derive a free-choice net







# Third attempt (decorated EPC)





PETER RITTGEN

MODIFIED EPCS AND THEIR FORMAL SEMANTICS

#### Decorated EPC

Applicable to any EPC diagram, provided that its designer add some information

We require:

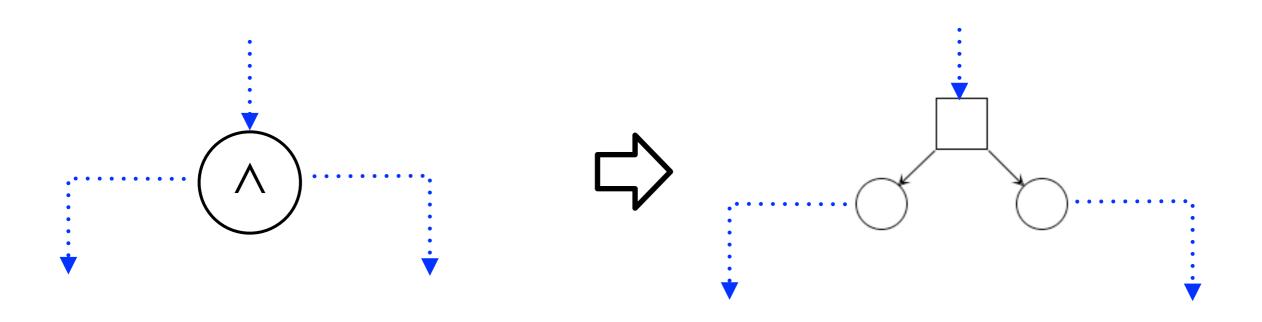
every (X)OR join is paired with a corresponding split (possibly of the same type)

OR-joins are decorated with a policy (avoid OR join ambiguous behaviour)

### Step 1: AND split

#### **EPC** element

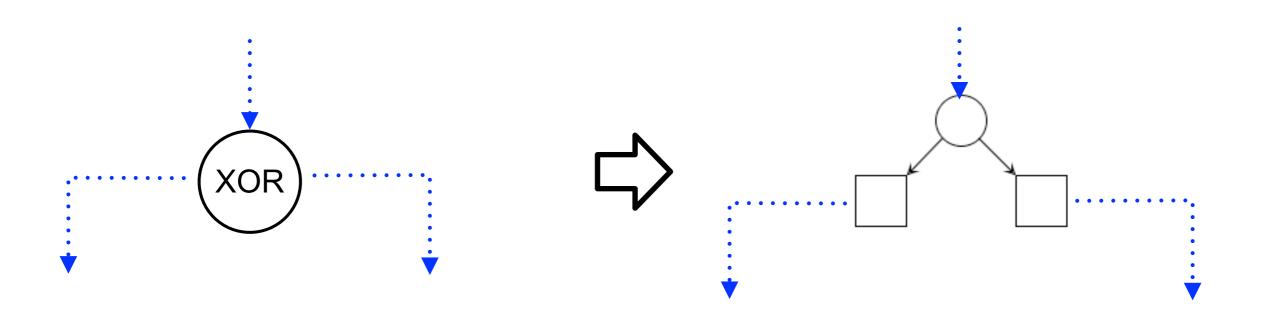
net fragment



#### Step 1: XOR split

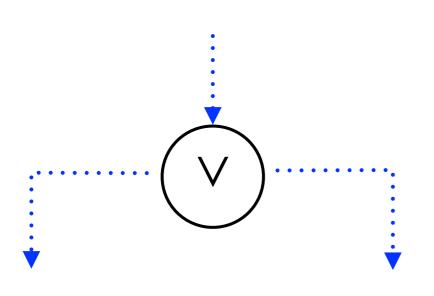
#### **EPC** element

net fragment

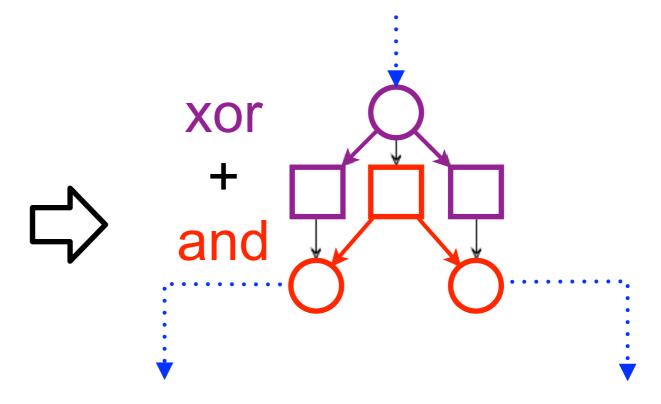


#### Step 1: OR split

#### **EPC** element



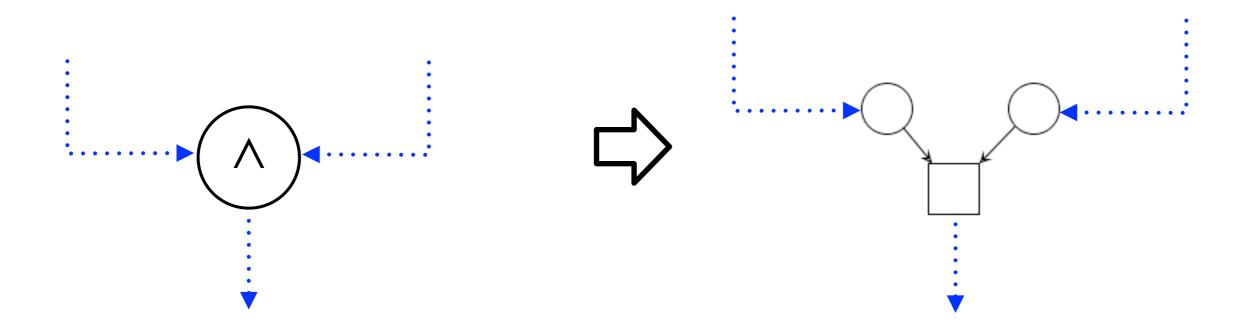
#### net fragment



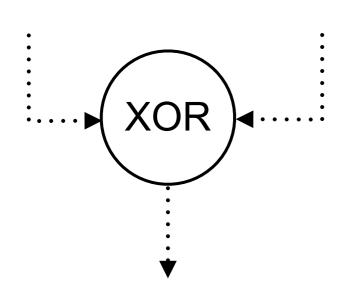
#### Step 1: AND join

**EPC** element

net fragment



### XOR join: intended meaning

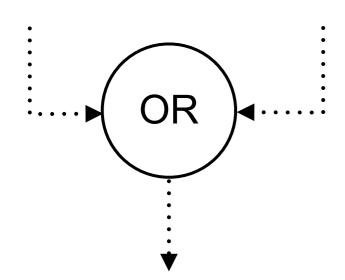


if both inputs arrive, it should block the flow

if one input arrives,
it cannot proceed unless
it is informed that
the other input will never arrive

### OR join: intended meaning

if only one input arrives, it should release the flow



if both inputs arrive, it should release only one output

if one input arrives,

it must wait until the other arrives or it is guaranteed that the other will never arrive

#### OR join: assumption

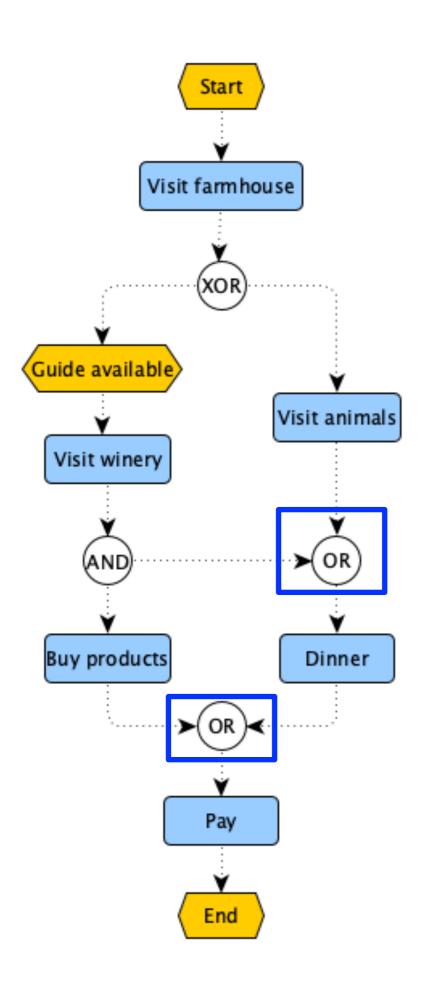
If an OR join has a **matching split**, its semantics is **wait-for-all**: wait for the completion of all *activated* paths

Otherwise, also other policies can be chosen:

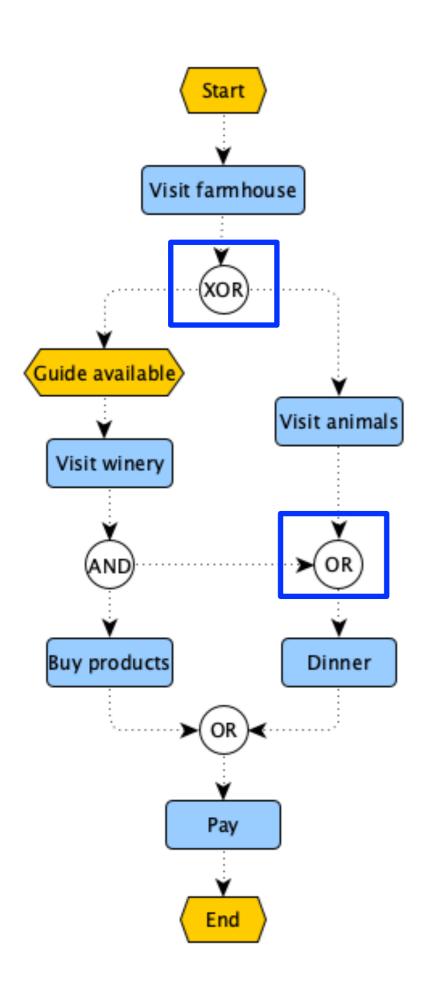
every-time: trigger the outgoing path on each input

first-come: wait for the first input and ignore the second

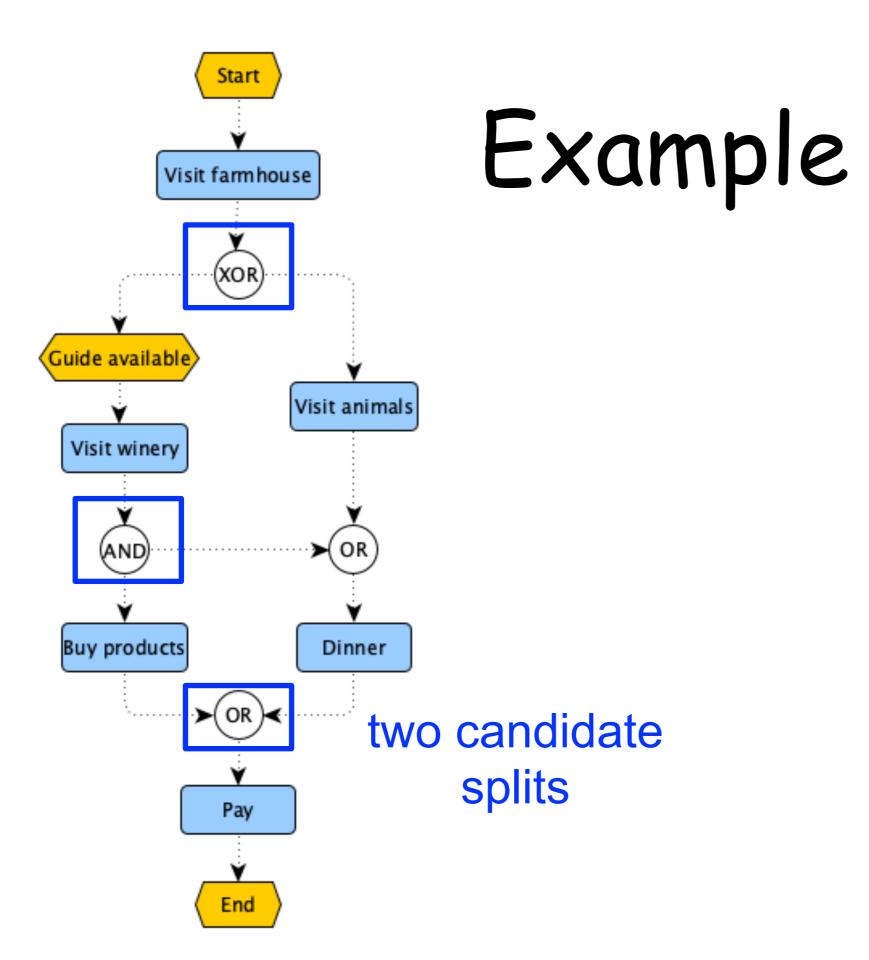
**Assumption**: every OR join is tagged with a policy (some suggested to have different trapezoid symbols)

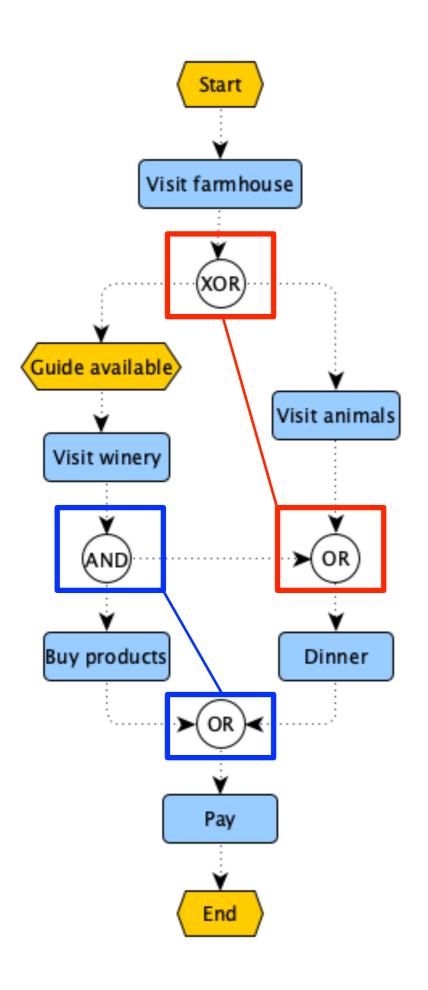


two OR joins but no OR split

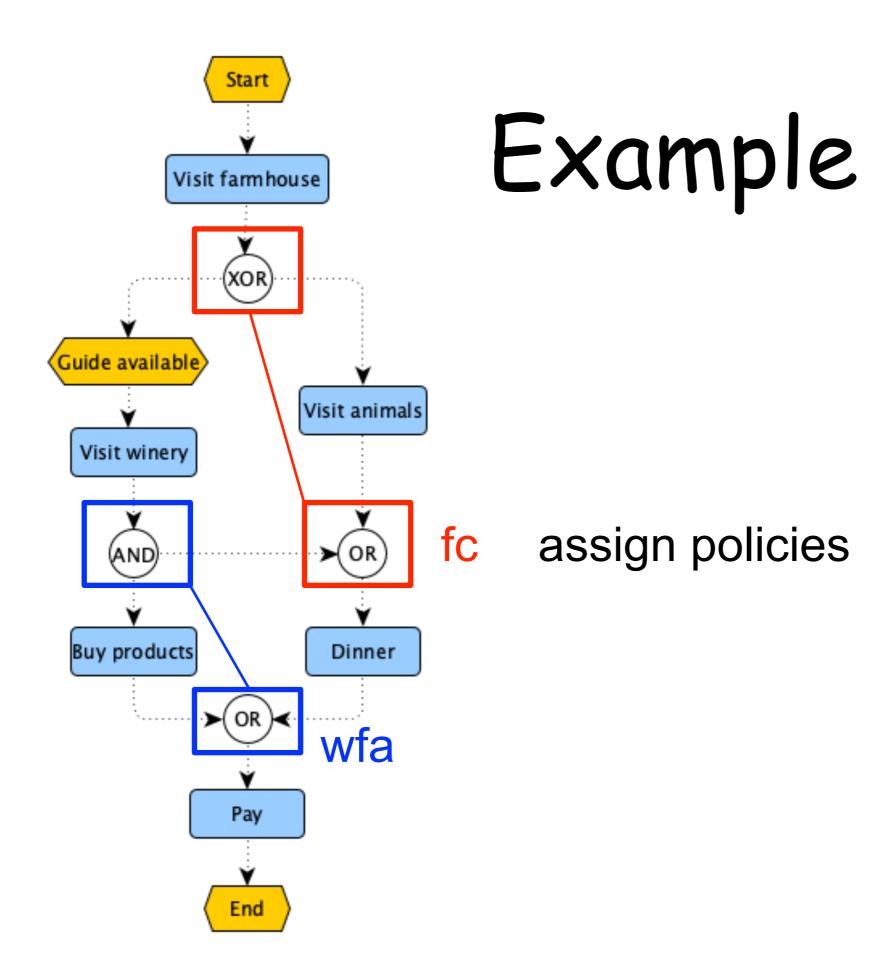


only one candidate split





assign corresponding splits



#### Assumption

. . .

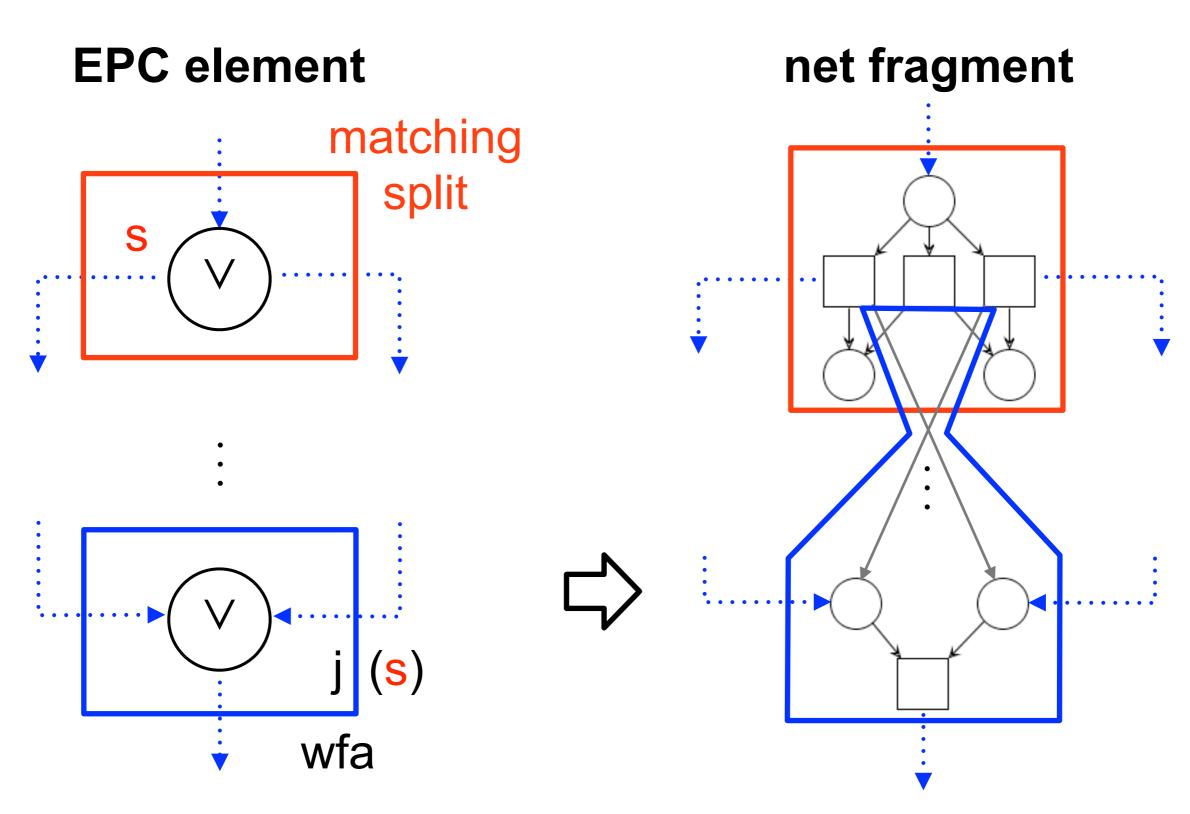
An OR join with matching split uses wfa

If an OR join has non-matching corresponding split it is decorated with a policy (wfa, fc, et)

wfa: wait-for-all works well with any corresponding split

. . .

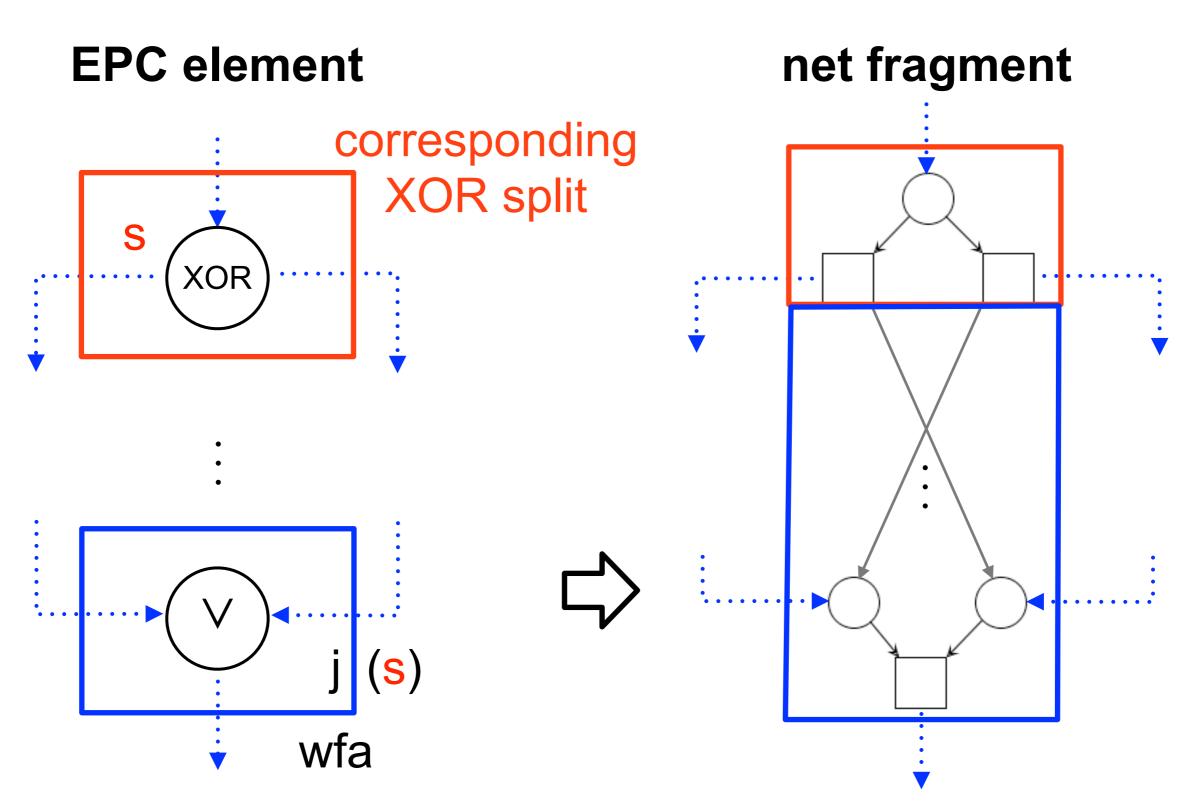
#### Step 1: OR join (wfa)



#### Step 1: OR join (wfa)

## **EPC** element net fragment corresponding **AND** split

#### Step 1: OR join (wfa)



#### Assumption

. . .

If an OR join has non-matching corresponding split it is decorated with a policy (wfa, fc, et)

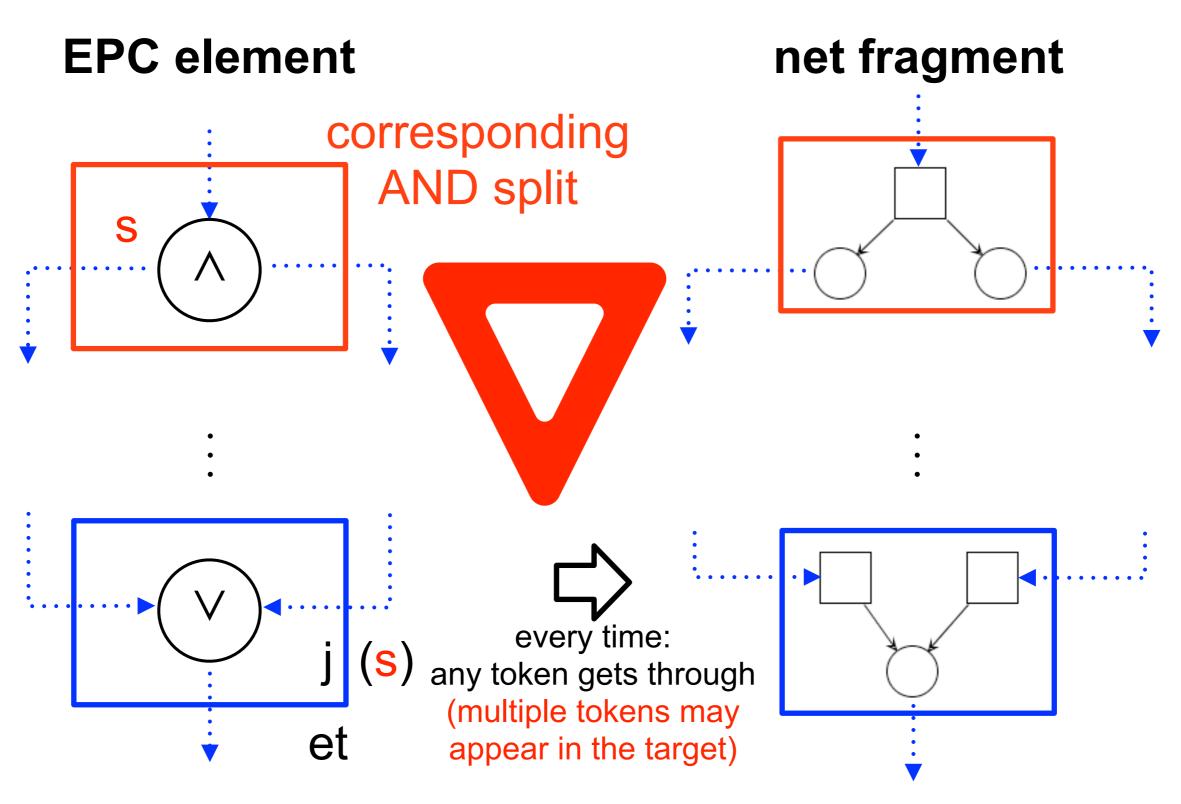
et: every-time works well with corresponding XOR split

. . .

### Step 1: OR join (et)

### **EPC** element net fragment corresponding XOR split XOR et

### Step 1: OR join (et)



#### Assumption

. . .

If an OR join has non-matching corresponding split it is decorated with a policy (wfa, fc, et)

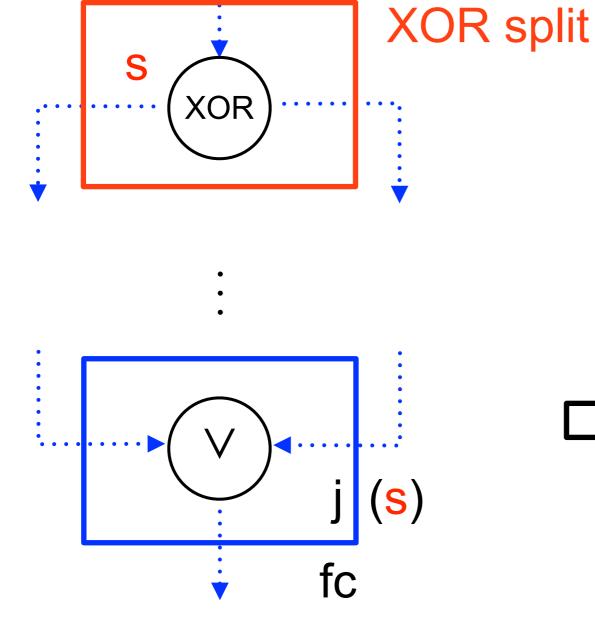
fc: first-come works well with corresponding XOR split

. . .

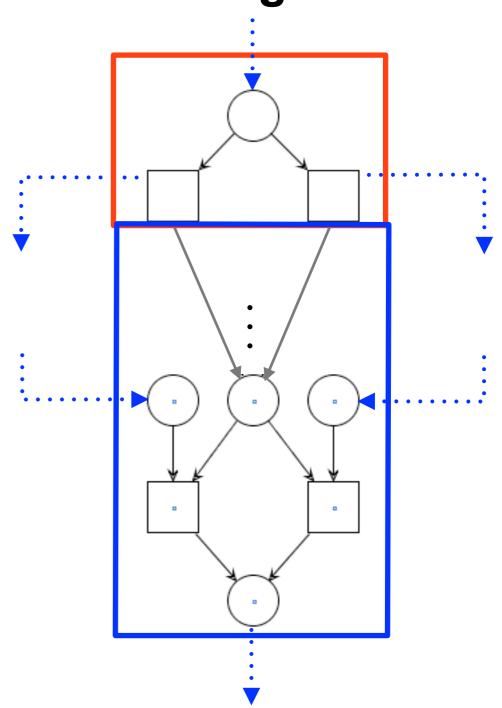
### Step 1: OR join (fc)

corresponding

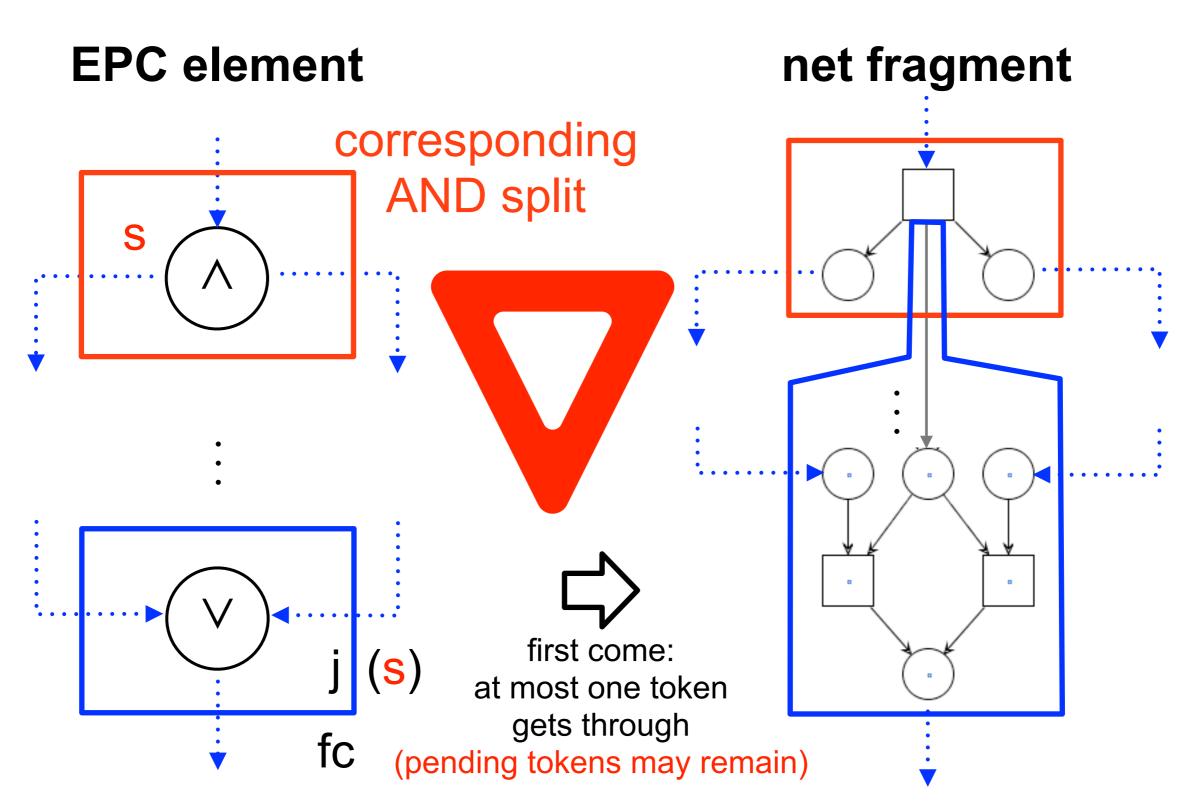
#### **EPC** element



#### net fragment



### Step 1: OR join (fc)



### XOR join: assumption

If a XOR join has a **matching split**, the semantics is: "it blocks if both paths are activated and it is triggered by a unique activated path"

Any policy (wait-for-all, first-come, every-time)

contradicts the exclusivity of XOR

(a token from one path can be accepted only if we make sure that no second token will arrive via the other path)

**Assumption**: every XOR join has a matching split (the implicit start split is allowed as a valid match)

#### Assumption

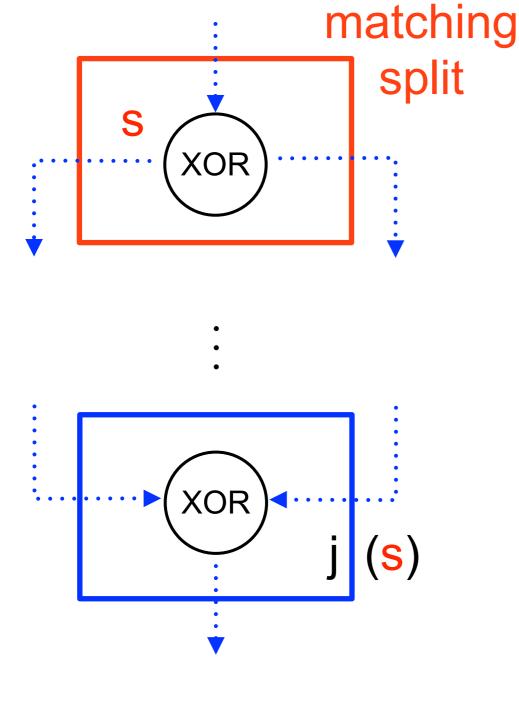
. .

Any XOR join has a corresponding matching split

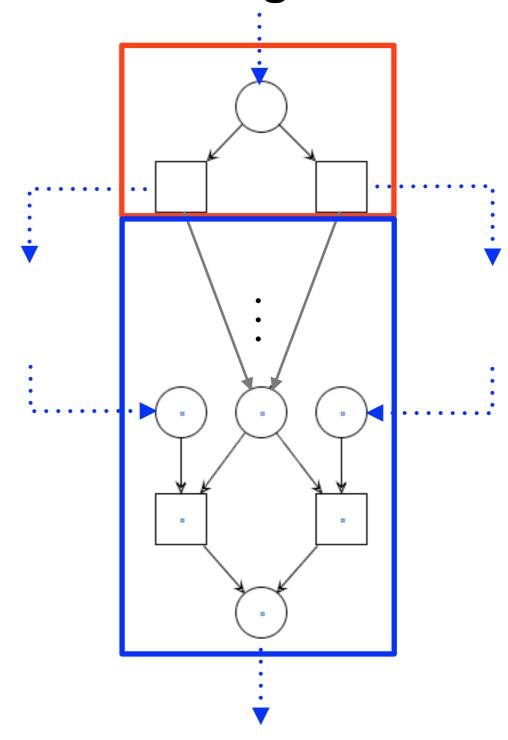
. . .

### Step 1: XOR join

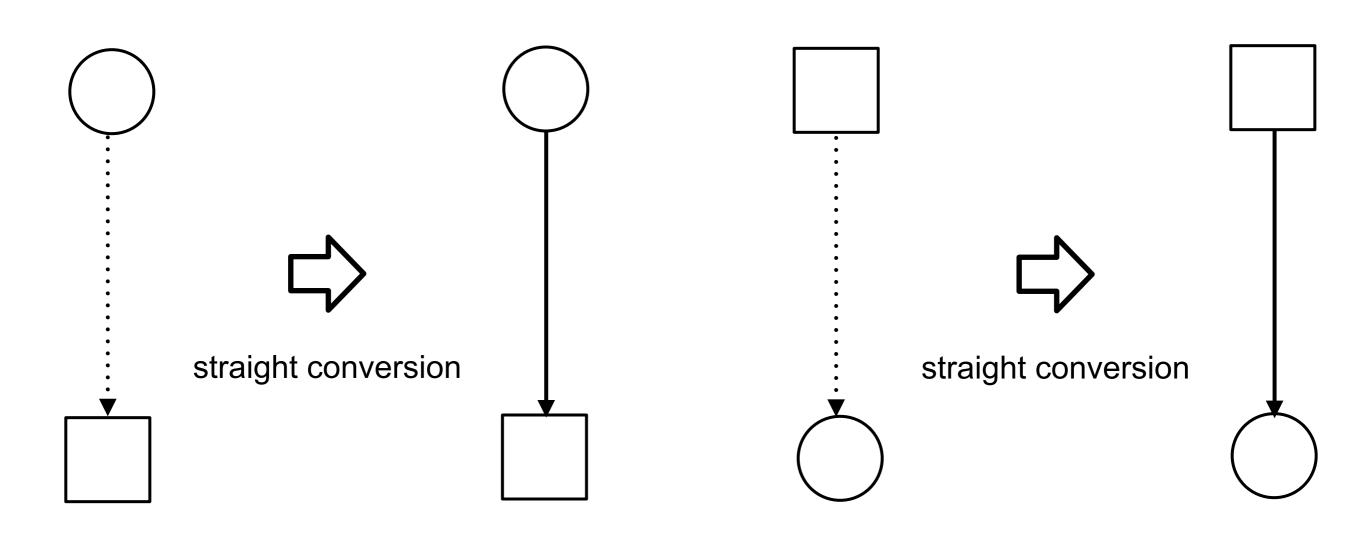
#### **EPC** element



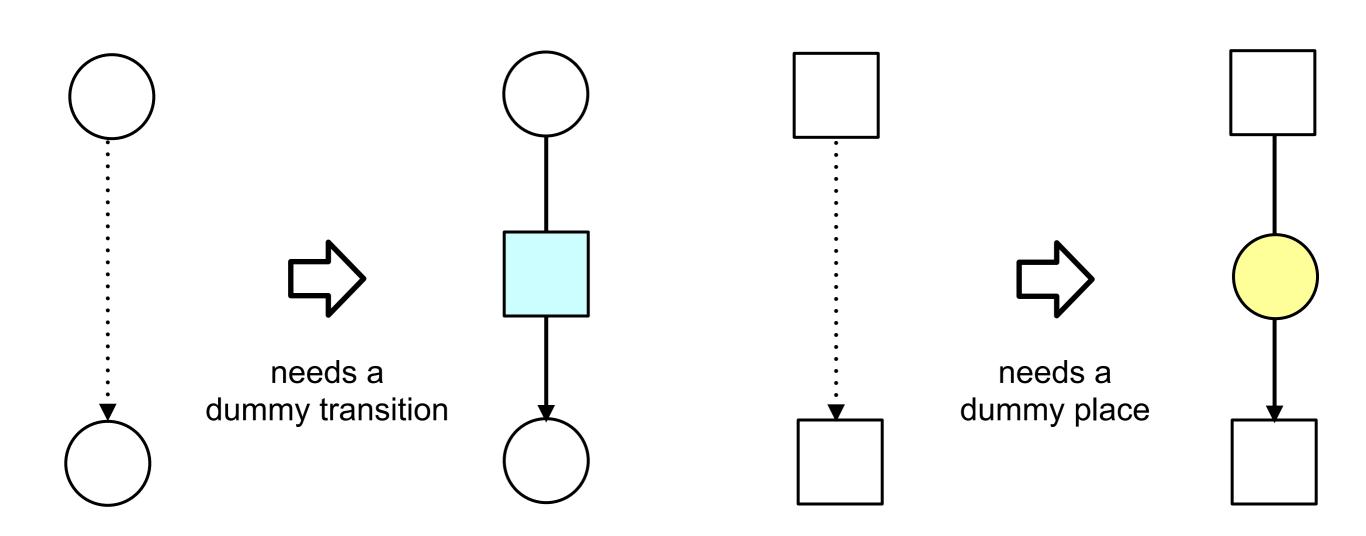
#### net fragment

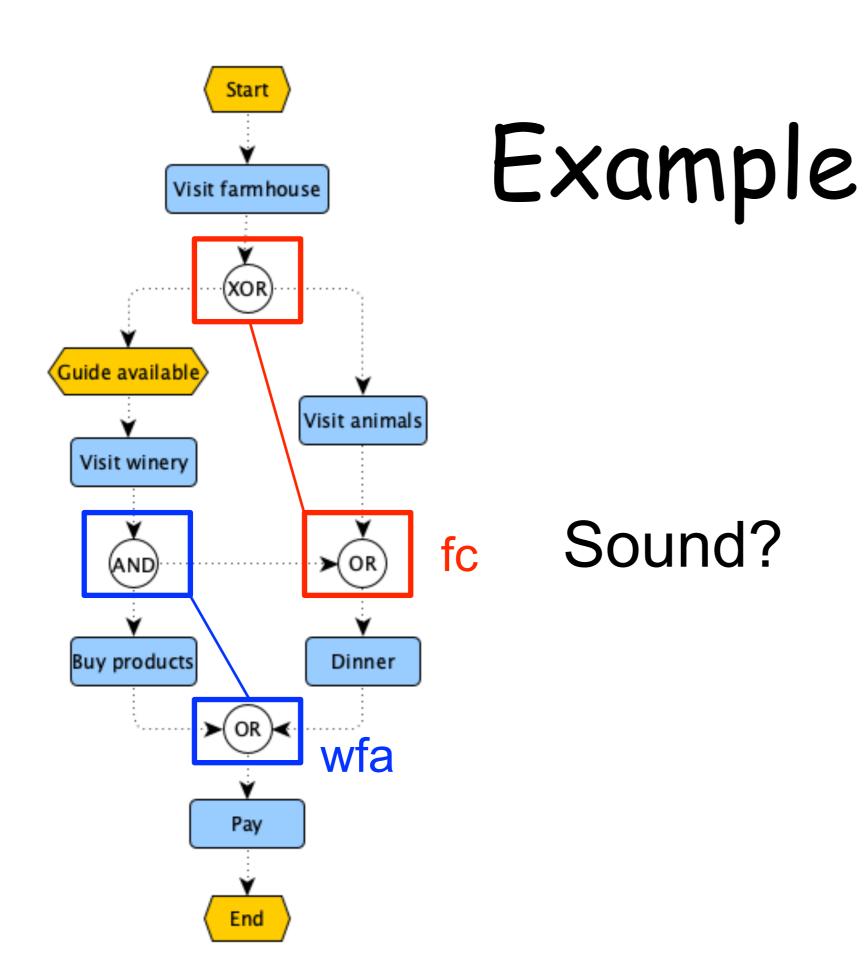


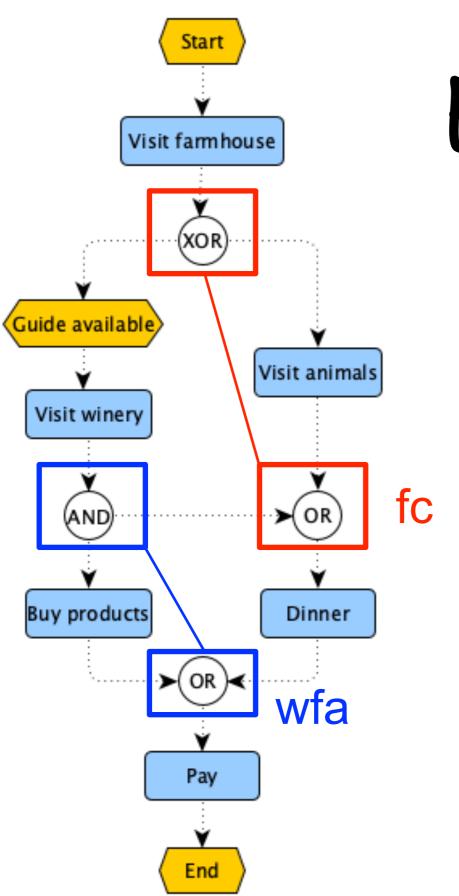
### Step 2: dummy style



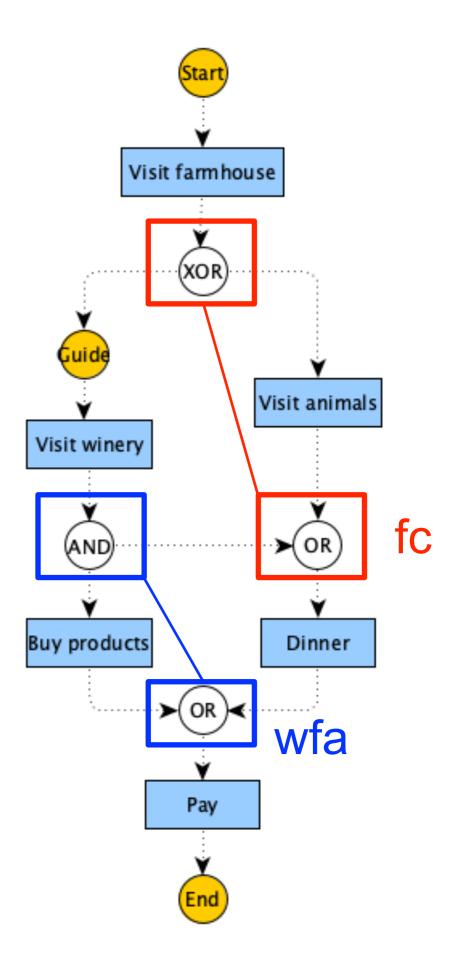
## Step 2: dummy style

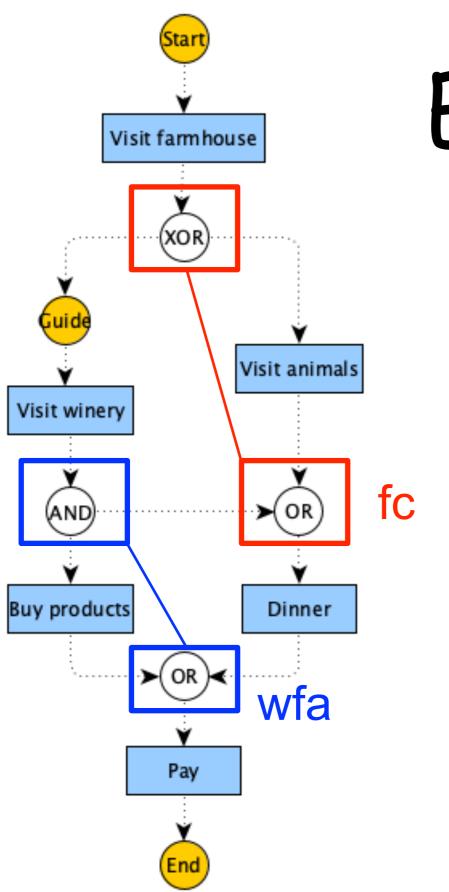




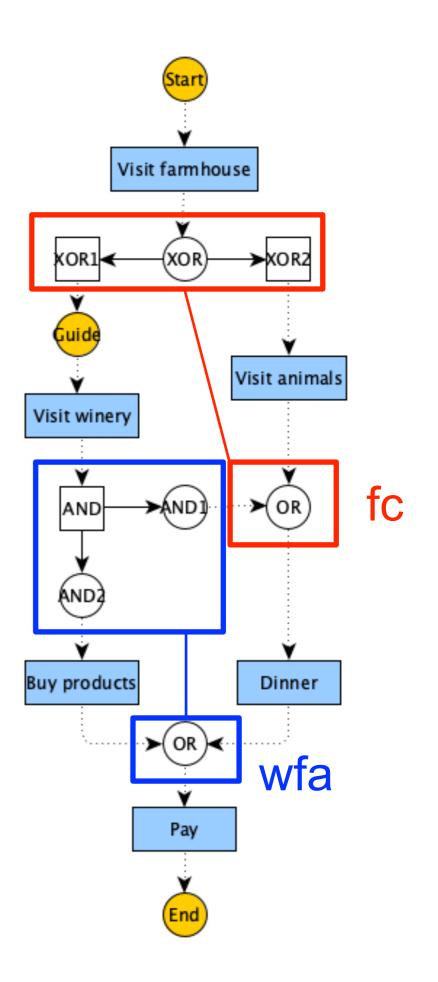


Step 1 events and functions





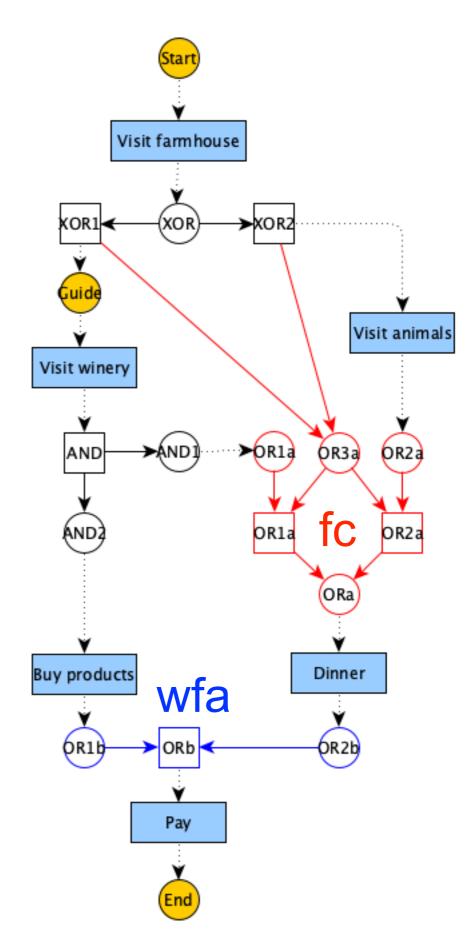


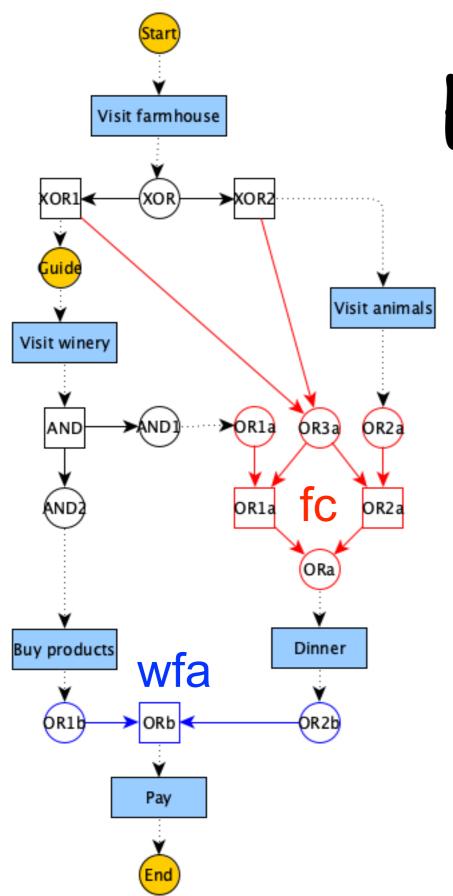


# Visit farmhouse Visit animals Visit winery fc Buy products Dinner wfa Pay

### Example

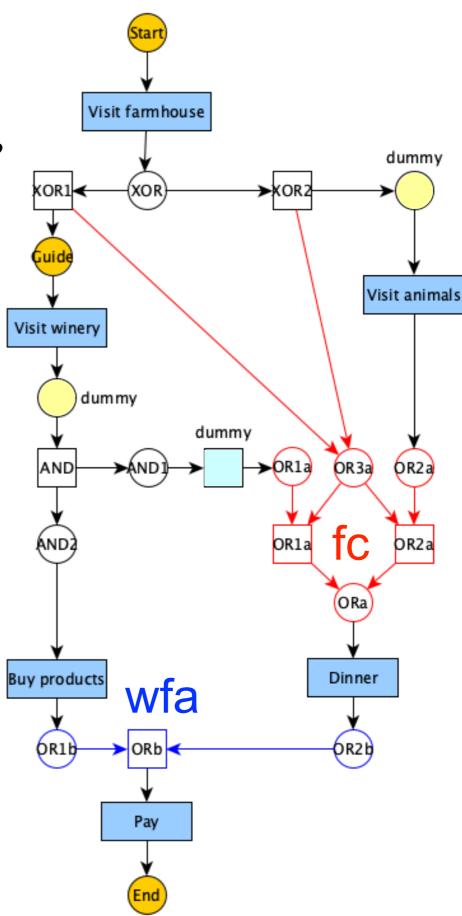
Step 1 splits and joins

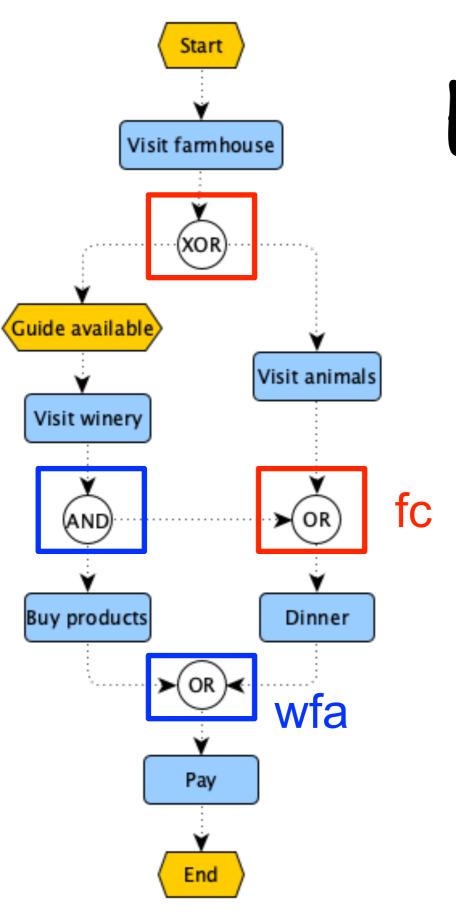






Step 2(+3) dummy style

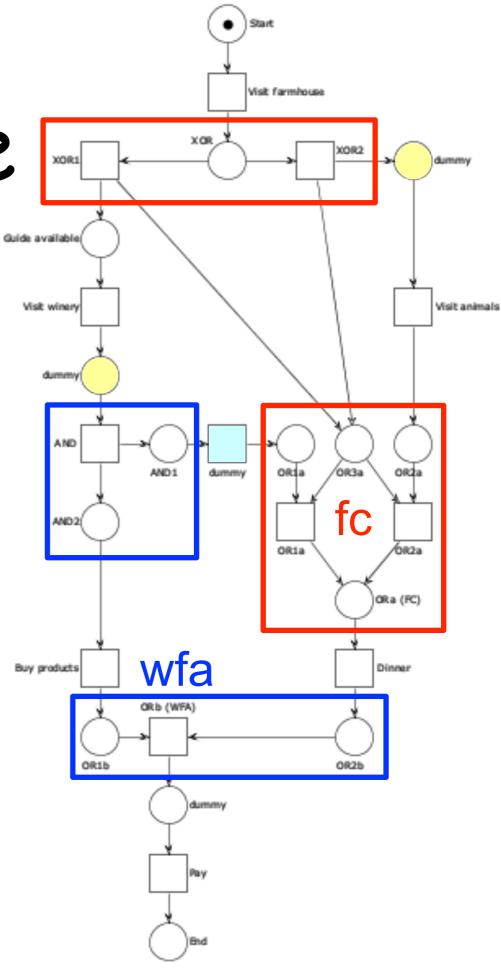


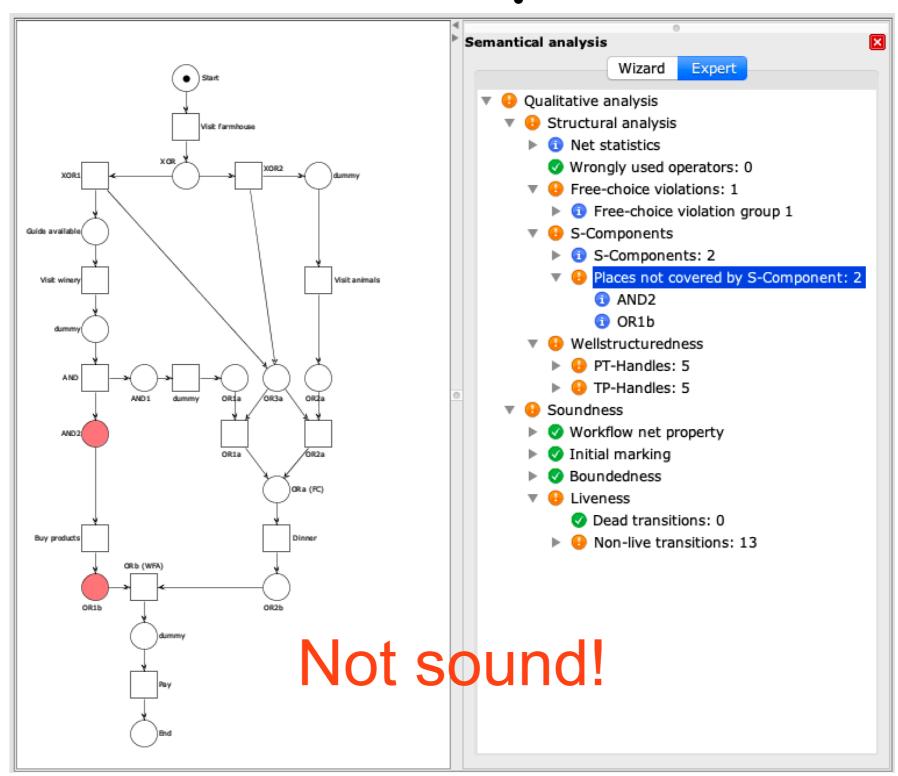


Sound?



Steps 1+2(+3)



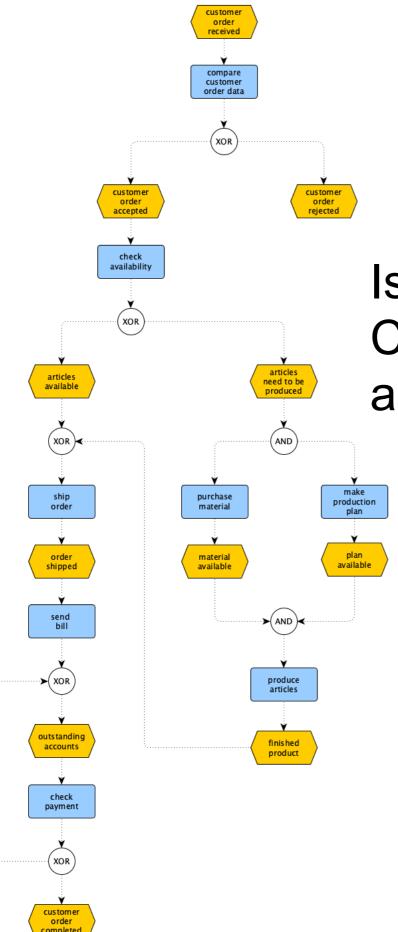


#### EPC pros and cons

You may leave complete freedom, but most diagrams will not be sound

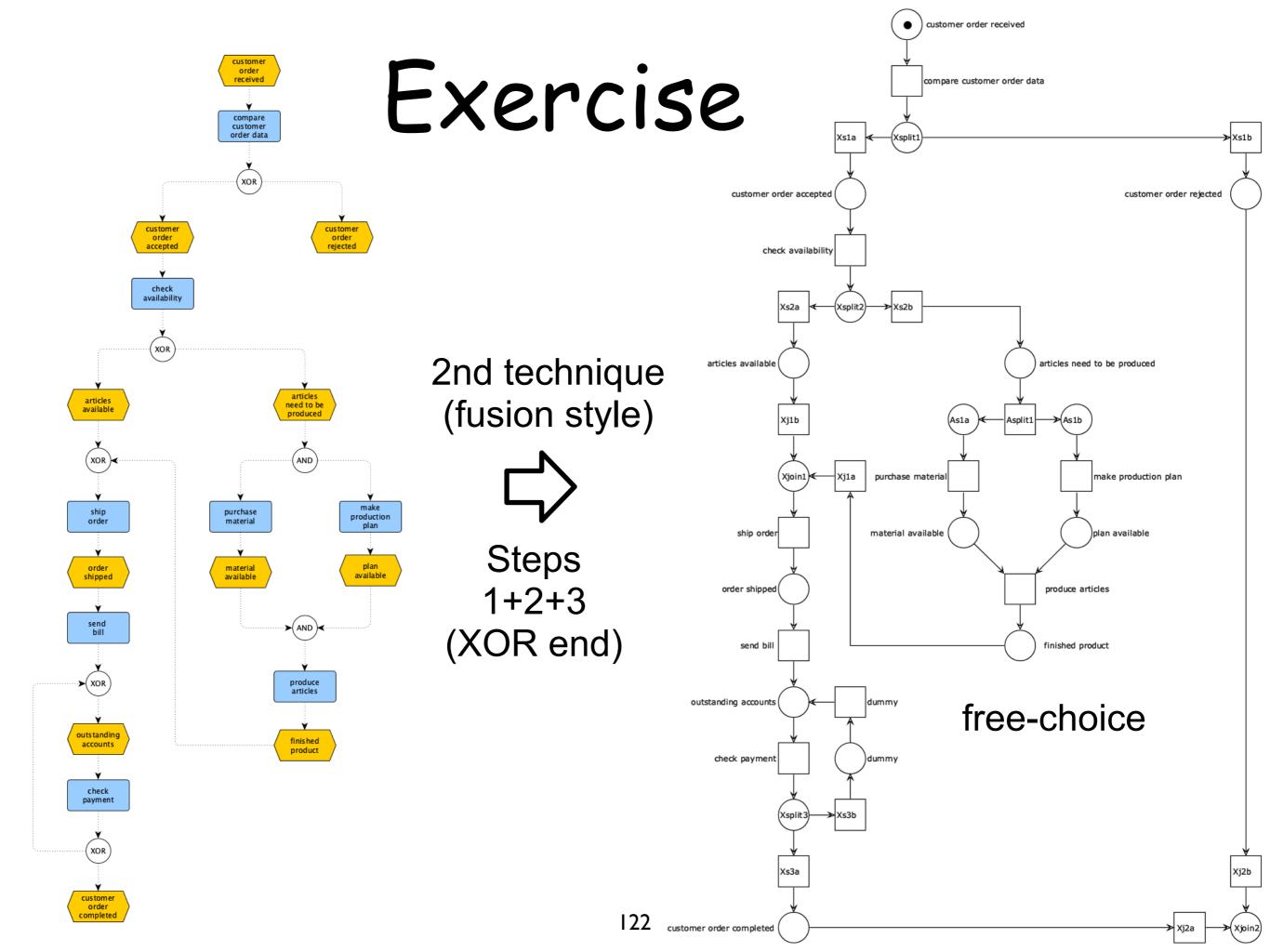
You may constrain diagrams, but people like flexible syntax and ignore guidelines

You may require to add decorations, but people will be lazy or misinterpret policies



#### Exercise

Is this EPC diagram sound? Choose one of the three techniques seen and apply it to answer the above question



#### Exercise

