



Preprocessing Mobility Data



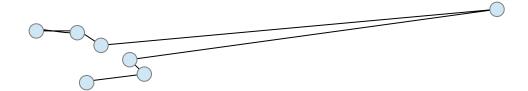
Content of this lesson

- Preprocessing trajectories
 - trajectory filtering
 - point map matching
 - route reconstruction
 - trajectory compression
 - Semantic enrichment
 - stop detection / trajectory segmentation
 - home location detection (GPS & MobPhones)
 - activity recognition (POI-based)

Trajectory filtering

- Data points are sometimes affected by errors
- Errors can have huge effects on results

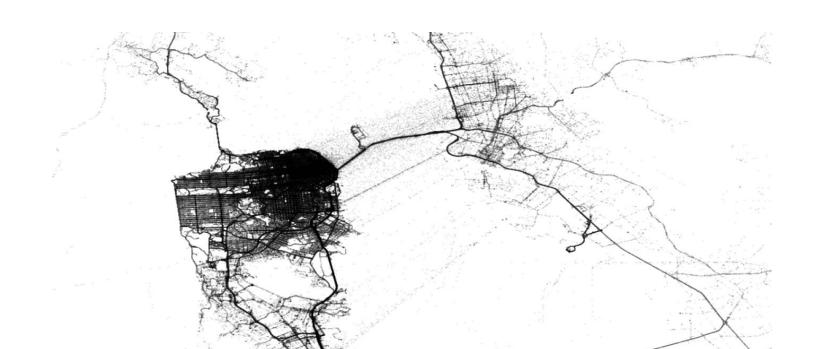
What is the real length of this trip?



- Two families of approaches:
 - Context-based filtering
 - Movement-based filtering

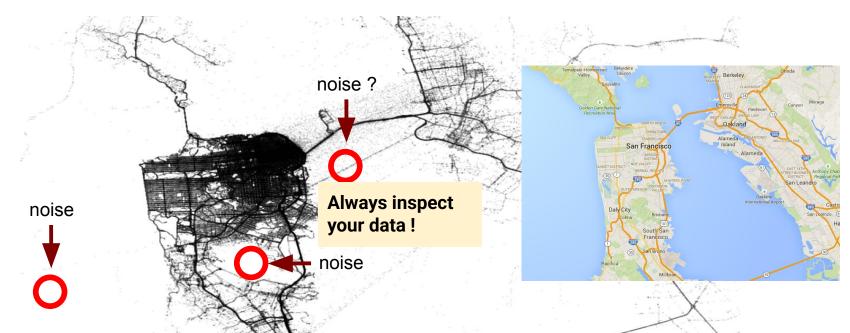
Context-based filtering

• Single points might contain errors of various kinds



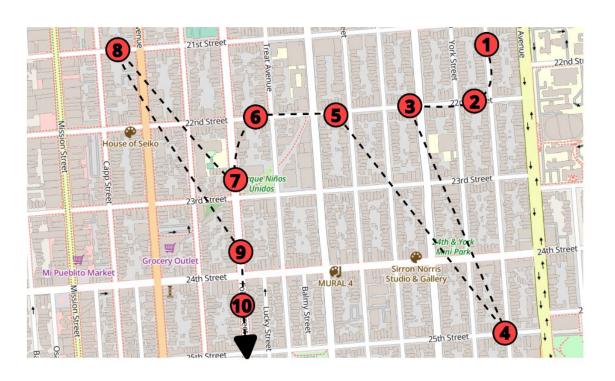
Context-based filtering

- Single points might contain errors of various kinds
- Map-based detection: cars on the water or out of roads are noise
 - Caution: do you trust 100% your map?



Context-based filtering

Why looking at the context might be not enough



Movement-based filtering

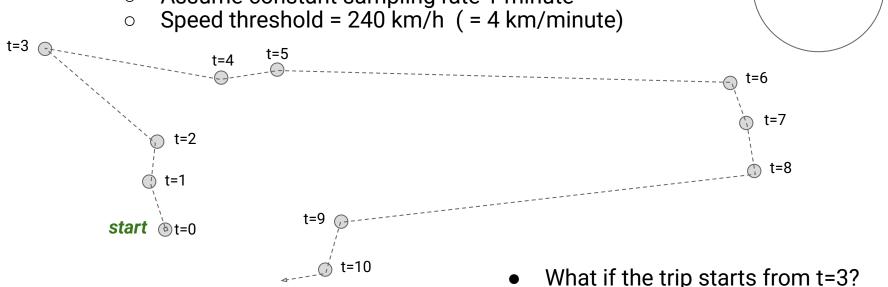
- No context is used, just the geometry / dynamics of movement
- Speed-based noise filtering approach:
 - The first point of the trajectory is set as valid
 - Scan all remaining points "p" of the trajectory (time order)
 - Compute "v" = average straight-line speed between point "p" and the previous valid one
 - If "v" is huge (e.g. larger than 400 km/h)
 => remove "p" from trajectory ("p" will not be used next to estimate speeds...)
 else
 => set "p" as valid



Movement-based filtering

Exercise

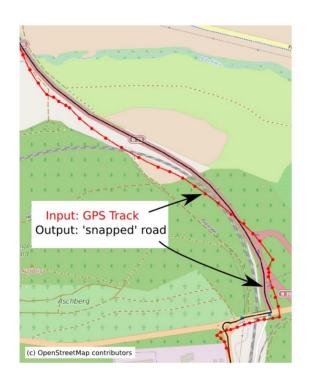
- What happens in this situation? (Multiple noisy points)
 - Assume constant sampling rate 1 minute



What if the trip starts from t=3?
 Which points of the trip survive?

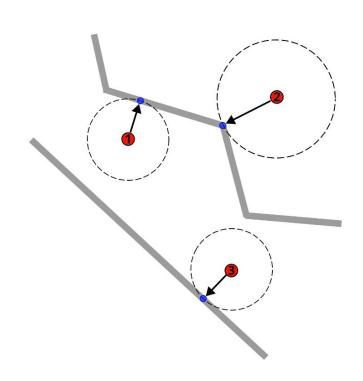
4 km

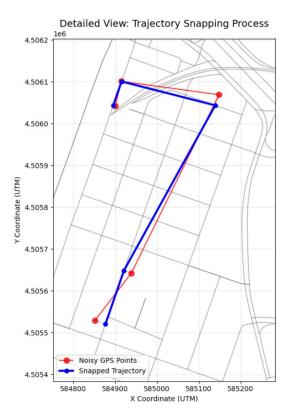
- Points can be aligned to the road network
 - Objective 1: improve accuracy of position
 - Objective 2: remove extreme cases (ref. filtering)
 - Objective 3: translate trajectories to sequences of road IDs
- Idea: project the point to the close location in the network
 - Usually there is a maximum threshold
 - Points farther than the threshold from any road are removed as noise



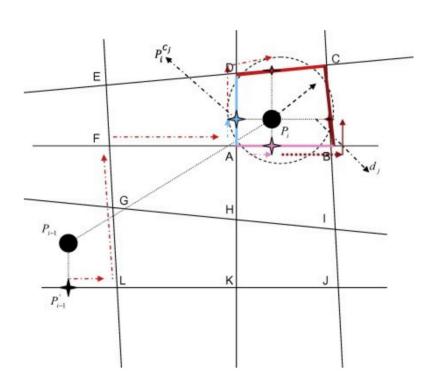
- Point projection
 - Requires to compare each point to each road segment
 - Assign closest point over closest segment

Anything wrong with ③?

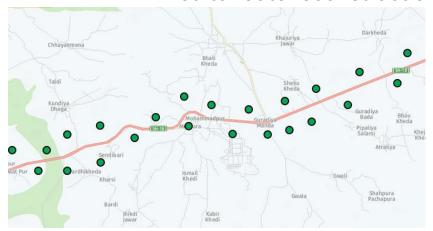


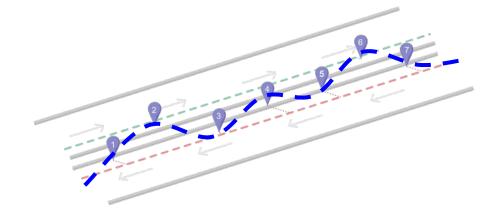


- In some contexts there can be multiple choices at reasonable distance
 - Very common in city centers
- Simply taking the closest one is "risky"

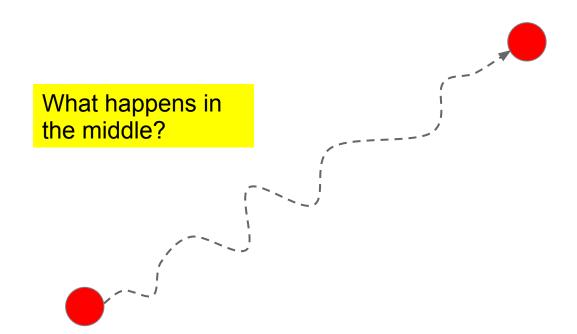


- Matching points separately can lead to inconsistent results
 - Mainly road-dense areas with position accuracy comparable to road separation
- Need a trajectory-level matching
 - Linked to route reconstruction





Sometimes the space/time gap between consecutive points is significant



- Typical solutions:
 - Free movement => straight line, uniform speed



points

- Typical solutions:
 - Constrained movement => optimal path





Optimizing what?

$$SP(o,d) = \arg\min_{p \in \mathcal{P}_{o,d}} \sum_{e \in p} cost(e)$$

- Trip length
 - cost(e) = length
- Travel time
 - cost(e) = length/speed
- Num. intersections
 - cost(e) = 1

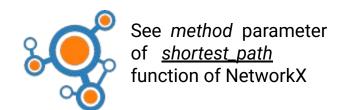


Generally referred as "Shortest paths", yet alternative "optimal paths" can be used

- Usually based on a notion of single-segment costs
- Most typical ones: path length, path duration (requires to know typical traversal times of roads)
- Alternative ones: fuel consumption, EV battery consumption, CO2 emissions, mixed costs

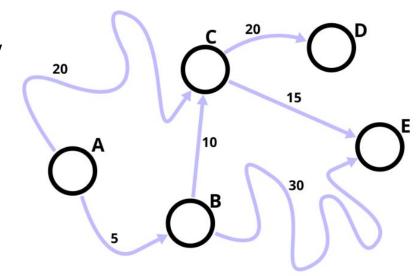
Algorithms applied are standard graph path optimization methods:

- ullet Dijkstra's algorithm \to efficient, requires that costs are non-negative
- \bullet Bellman-Ford algorithm \rightarrow less efficient, can work with negative weights (but no cycles)



Refresher: Dijkstra's minimum cost algorithm

 Simple and efficient: O(M + N log N) time complexity (M = |edges|, N = |nodes|)

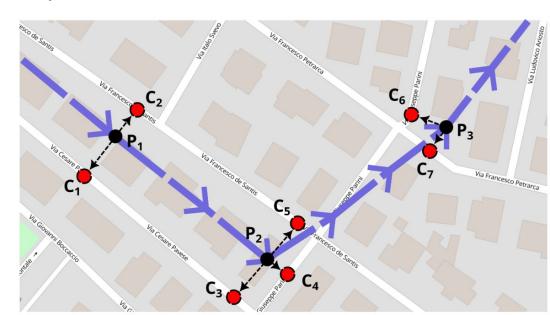


	A	В	C	D	E
Init	0	∞	∞	∞	∞

Trajectory Map Matching

- Assigns points to road segments
- Reconstructs the movement between consecutive points
- Ensures coherence of the overall process

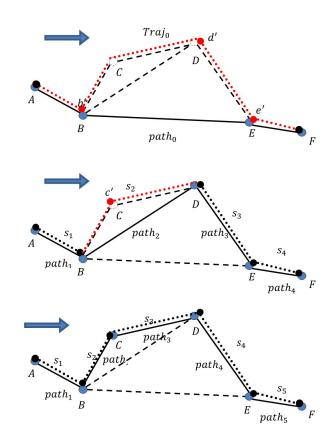
- Two sample approaches:
 - Based on shortest path
 - Based on probabilities



Shortest path-based Map Matching

Zhu-Honda-Gonder method, used by MappyMatch

- Similar ideas as trajectory simplification
 - Match first and last point
 - Compute shortest path on the network
 - Compute "fitting" of path w.r.t. Trajectory (based on distance of real points from path)
 - If fitting < threshold ⇒
 - split into two or more parts
 - run recursively the process on both



Reference: Zhu, honda & Gonder. A Trajectory Segmentation Map Matching Approach for Large-Scale, High-Resolution GPS Data. TRB 2017.

Probability-based Map Matching

Newson-Crumm method, used by <u>pyTrack</u>

Point-to-road (p->C) assignments, with probabilities

$$p(p_i \mapsto C_k) = e^{-\frac{1}{2}d(p_1, C_k)^2/\sigma^2}/\sqrt{2\pi}\sigma$$

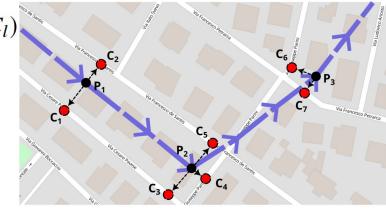
Road-to-road (C->C) assignments

$$P(C_k \to C_l) = e^{-\Delta/\beta}/\beta$$

 $\Delta = shortest_path_dist(C_k, C_l) - d(C_k, C_l)$

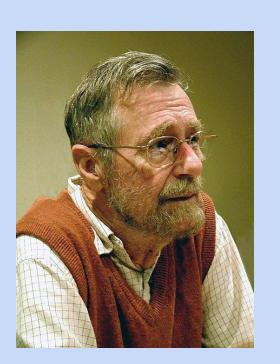
Compute best path that visits all points in the correct sequence

maximize
$$\Pi_i p(p_i \mapsto C_i) p(C_i \to C_{i+1})$$



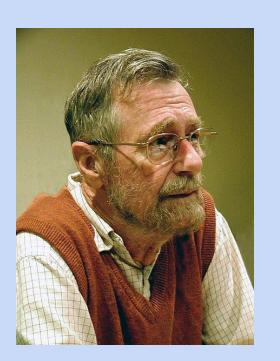
Reference: <u>Newson & Krumm. Hidden Markov Map Matching</u>
<u>Through Noise and Sparseness. ACM GIS'09.</u>

Edsger W. Dijkstra



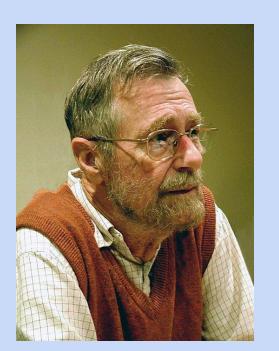
- 1930 2002
- Dutch computer scientist, programmer, software engineer, systems scientist, and science essayist
- 1972 Turing Award for "fundamental contributions to developing programming languages"

Dijkstra is famous for...



- Dijkstra's algorithm, of course
- Contributions to "self-stabilization of program computation"
 - Won him the "ACM PODC Influential Paper Award", later renamed "Dijkstra Prize"
- Hundreds of papers on computational and science philosophy issues

Dijkstra is famous for...



- His habit of writing everything with paper
 & fountain pen
- Hundreds of papers, many unpublished
 - E. W. Dijkstra Archive
- Counting should start from 0, not 1...

When dealing with a sequence of length N, the elements of which we wish to distinguish by subscript, the next vexing question is what subscript value to assign to its starting element. Adhering to convention a) yields, when starting with subscript 1, the subscript range 15 i N+1; starting with 0, however, gives the nicer range 05 i N. So let us let our ordinals start at zero: an element's ordinal (subscript) equals the number of elements preceding it in the sequence. And the moral of the story is that we had better regard -after all those centuries! - zero as a most natural number.

Dijkstra the teacher



- Chalk & blackboard, no projectors
- No textbooks
- Improvisation & long pauses
- No references in papers

"For the absence of a bibliography I offer neither explanation nor apology."

- Long exams
 - Each student was examined in Dijkstra's office or home, and an exam lasted several hours

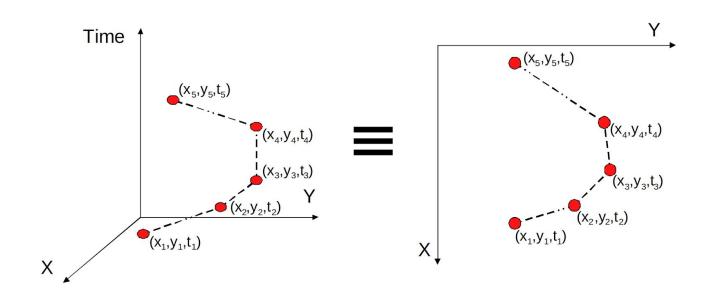
Trajectory compression / simplification

- Many algorithms for trajectories are expensive
 - Their complexity depends on the number of points
 - Sometimes trajectories have more points than needed

- Objective of compression / simplification
 - Reduce the number of points...
 - ... without affecting the quality of results

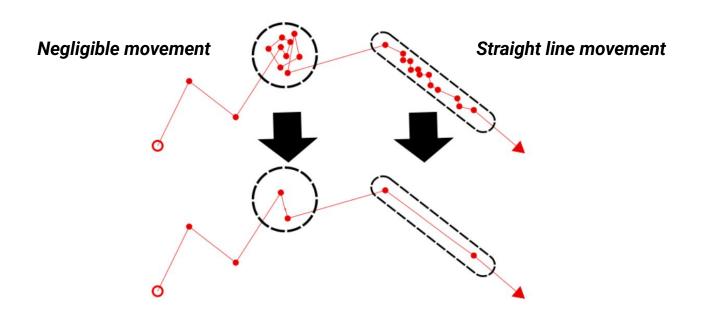
Trajectory data

- A trajectory is a temporal sequence of time-stamped locations
- Most methods focus on the spatial component



Trajectory compression / simplification

• Typical cases where points might be removed



Compression/simplification methods

Some standard methods for simplifying polygonal curves:

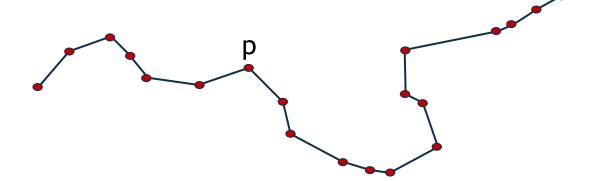
- Ramer-Douglas-Peucker, 1973
- Driemel-HarPeled-Wenk, 2010

1972 by Urs Ramer and 1973 by David Douglas and Thomas Peucker

Mostly known as Douglas-Peucker (DP) algorithm

The most successful simplification algorithm. Used in GIS, geography, computer vision, pattern recognition...

Very easy to implement and works well in practice.



```
Input: Trajectory T = \langle p_1, \dots, p_n \rangle, spatial threshold \epsilon
Output: Simplified trajectory T'

1 i \leftarrow \text{index } j \in [1, \dots, n] such that d(p_j, \overline{p_1, p_n}) is maximized;

2 if d(p_i, \overline{p_1, p_n}) > \epsilon then // p_i deviates too much, split!

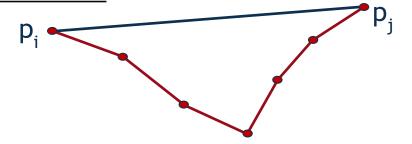
3 T'_{left} \leftarrow DP(\langle p_1, \dots, p_i \rangle);

4 T'_{right} \leftarrow DP(\langle p_i, \dots, p_n \rangle);

5 \mathbf{return} joined T'_{left} and T'_{right};

6 else

7 \mathbf{return} \mathbf{return} \mathbf{return}
```



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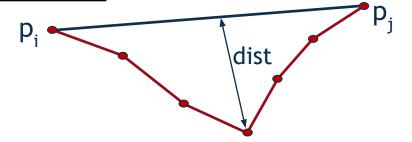
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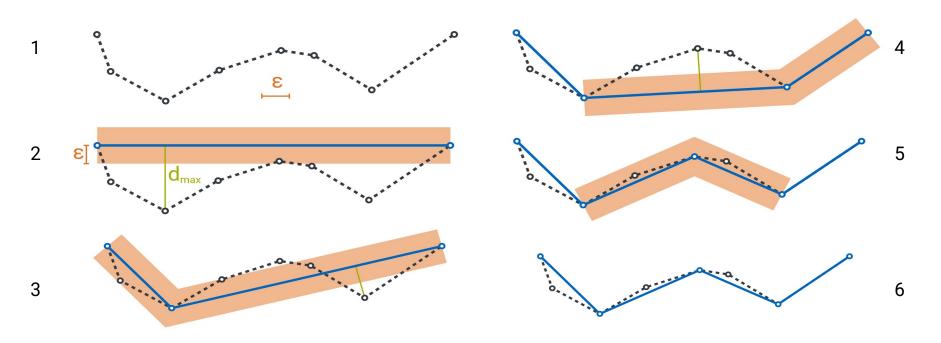
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```





Time complexity?

Testing a shortcut between p_i and p_i takes O(j-i) time.

Worst-case recursion?

 $DP(P, v_i, v_{i+1})$ $DP(P, V_{i+1}, V_i)$

Time complexity $T(n) = O(n) + T(n-1) = O(n^2)$

But rather fast in practice

```
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```

Output: Simplified trajectory T'

- 1 $i \leftarrow \text{index } j \in [1, ..., n]$ such that $d(p_i, \overline{p_1, p_n})$ is maximized; // p_i deviates too much
- 2 if $d(p_i, \overline{p_1, p_n}) > \epsilon$ then
- $3 \mid T'_{left} \leftarrow DP(\langle p_1, \ldots, p_i \rangle);$
- $T'_{right} \leftarrow DP(\langle p_i, \ldots, p_n \rangle);$
- **return** joined T'_{left} and T'_{right} ;
- 6 else
- return T;

Driemel-HarPeled-Wenk

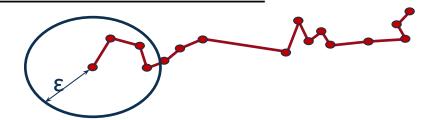
- Simplest solution
- Aims to produce well spaced points
 - \circ Basically a spatial subsampling: takes a point every ε meters
 - More local vision than The DP algorithm
- Very fast and easy to implement

Driemel-HarPeled-Wenk

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Input: Trajectory T = \langle p_1, \dots, p_n \rangle, spatial threshold \epsilon
  Output: Simplified trajectory T'
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i \leftarrow 1;
3 while i < n \operatorname{do}
       q \leftarrow p_i;
     i \leftarrow \text{smallest index } j \in [i+1,\ldots,n] \text{ such that } d(q,p_j) > \epsilon;
        if i is undefined then i \leftarrow n // We reached end of T
       Append p_i to T';
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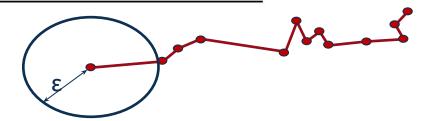
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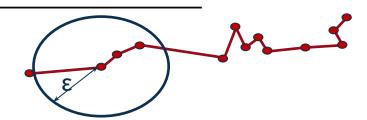
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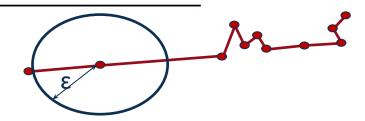
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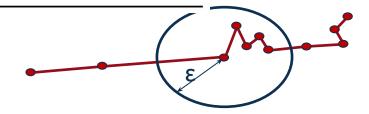
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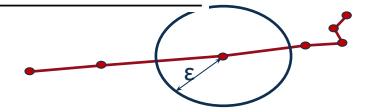
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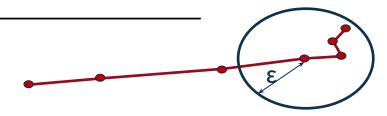
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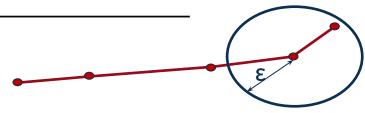
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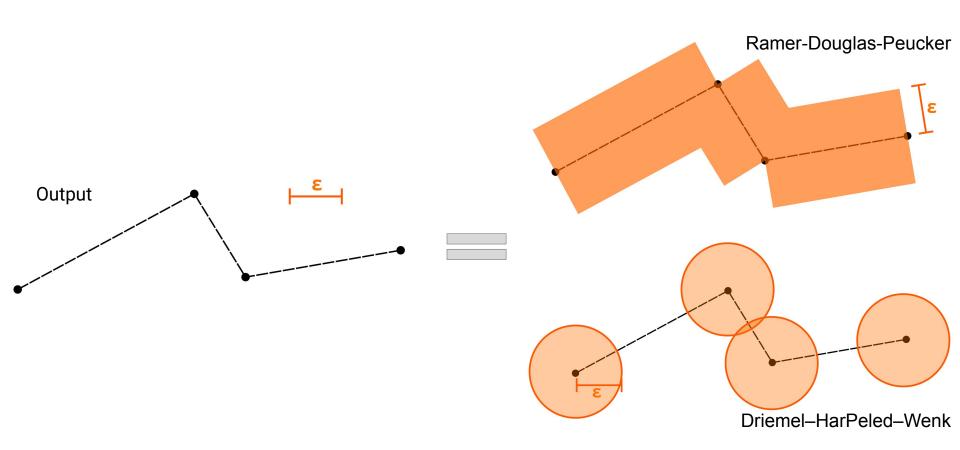
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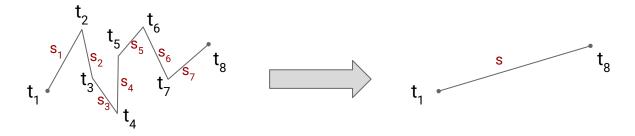
Ramer-Douglas-Peucker vs Driemel-HarPeled-Wenk



Impact on speed

- What about time and speeds?
 - Time-stamps were never considered in the algorithms
 - They considered on impact on space / geometry of trajectories
 - What impact on time-related aspects, e.g. speed?

Typically, simplification reduces average speed estimates:

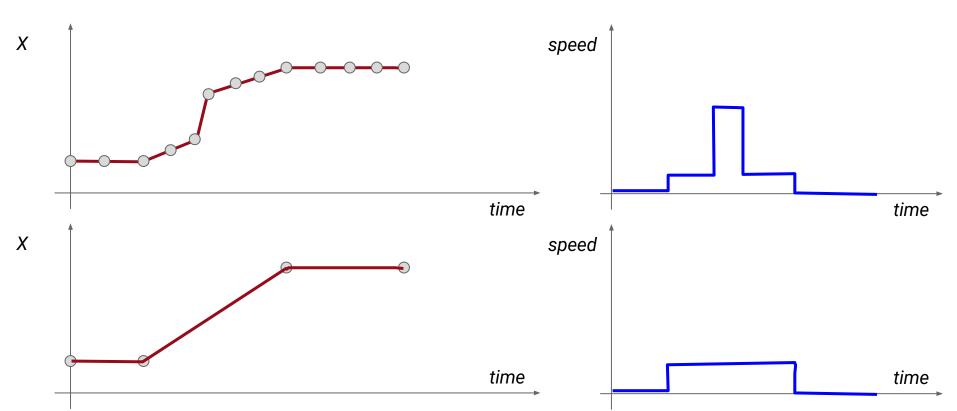


Avg speed =
$$(s_1 + ... + s_7) / (t_8 - t_1)$$

 Avg speed = $s / (t_8 - t_1)$

Impact on speed

Flattening effect



How fast is a cow?

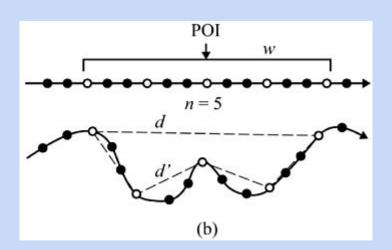


How fast is a cow?

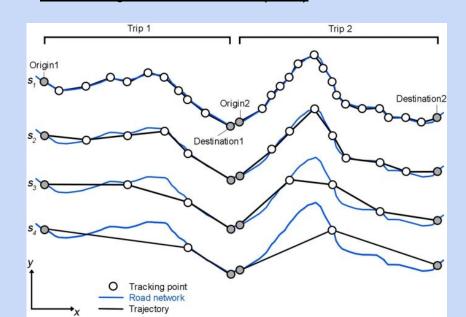
- Trajectory compression / simplification changes the scale of the analysis
 - Simplified data → macroscopic analysis
 - Detailed data → microscopic analysis
- Several movement characteristics can be affected

How fast is a cow?

How fast is a cow? Cross-Scale Analysis of Movement Data Laube P, Purves RS (2011)

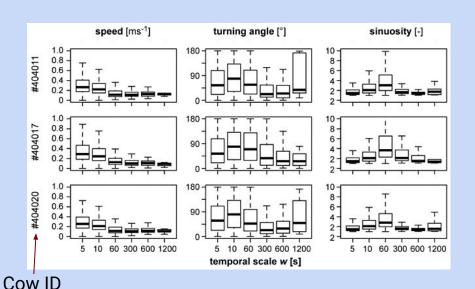


Understanding the impact of temporal scale on human movement analytics Su, R., Dodge, S. & Goulias, K.G (2022)



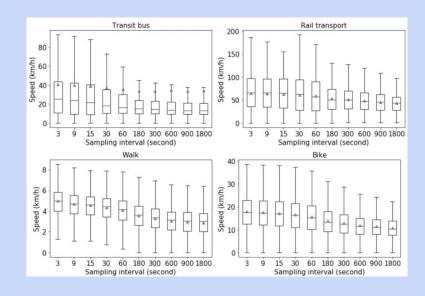
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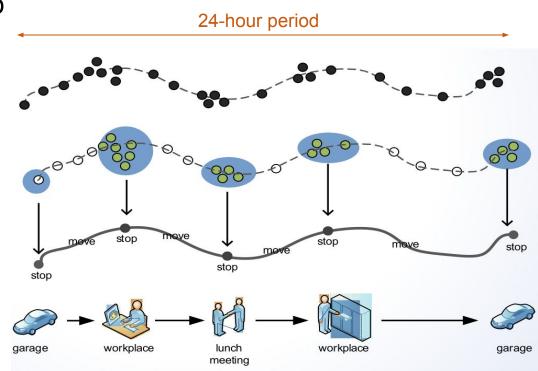


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 - route reconstruction
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 - Semantic enrichment
 - **■** stop detection / trajectory segmentation
 - home location detection (GPS & MobPhones)
 - activity recognition (POI-based)

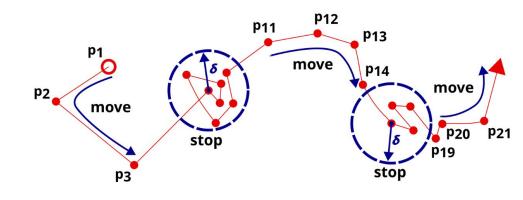
Stop detection & Trajectory segmentation

- Raw data forms a continuous stream of points
- Typical unit of analysis: the trip
- How to segment?
 - Basic idea: identify stops



Stop detection & Trajectory segmentation

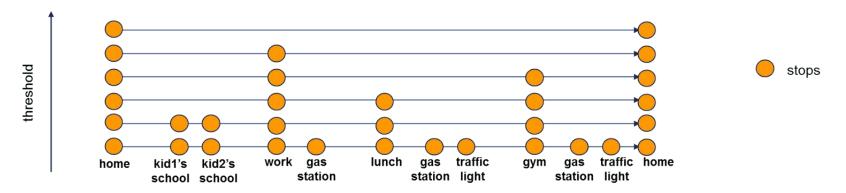
- General criteria based on speed
 - If it moves very little (threshold Th_s) over a significant time interval (threshold Th_T)
 - => it is practically a stop
 - Trajectory (trip) = contiguous sequence
 of points between two stops
 - Typical values:
 - o Th_s within [50, 250] meters
 - \circ Th_T within [1, 20] minutes



$$(Th_S = \delta, Th_T = 3', sampling rate = 1 pt/minute)$$

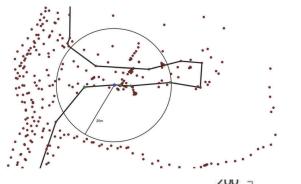
Stop detection & Trajectory segmentation

Different time thresholds yield different semantics



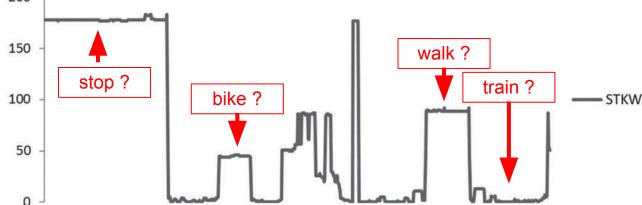
- Which one is the best for you?
 - Application dependent

Generalization: transportation means segmentation



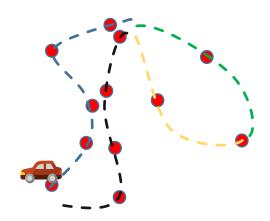
- Speed / density-based approach
- Idea: faster means less of my points around me

Number of points within radius R



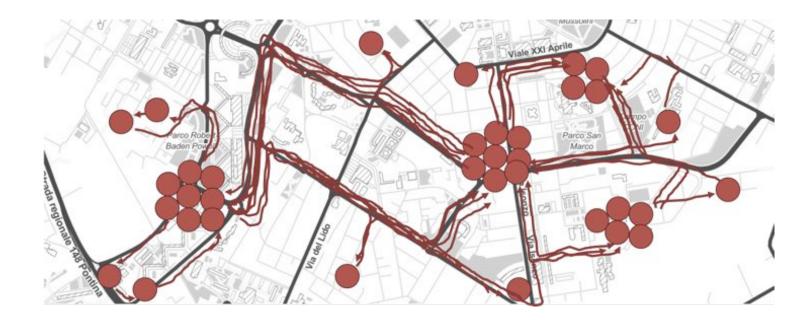
User's Mobility History

- What do we get after segmentation?
- Several trajectories associated to the same subject
- Enables individual-level analyses
 - E.g. explore user's habits, find deviations from usual, etc.



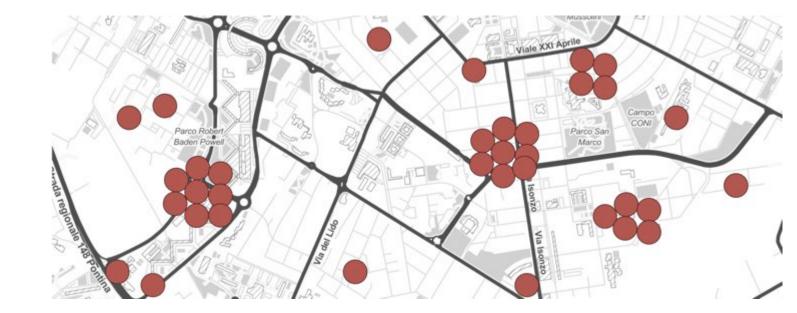
Inferring Home / Work locations

- Take all trips of a vehicle / user
- Build a "Individual Mobility Network"
 - o Graph abstraction of the overall mobility based on locations (nodes) and movements (edges).





- Focus on start and stop points
 - Dense areas represent important places





Cluster points to identify locations



Viale XXI Aprile

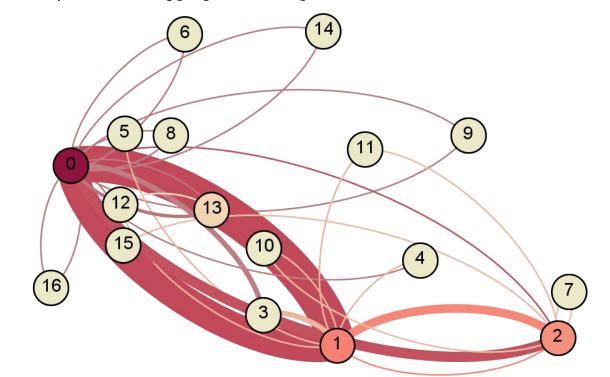


Each location is characterized by its frequency





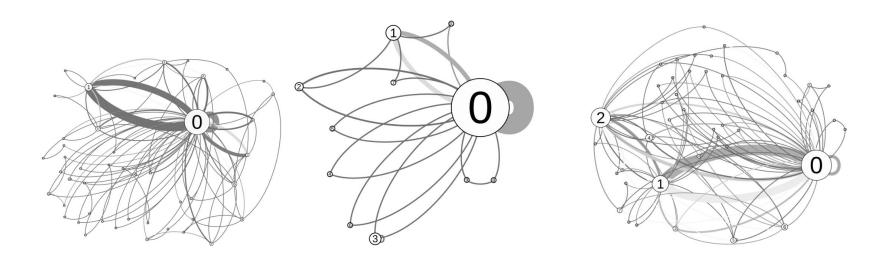
• Trips between points area aggregated as edges/flows between nodes/locations



Directions of flows:



• Some real examples

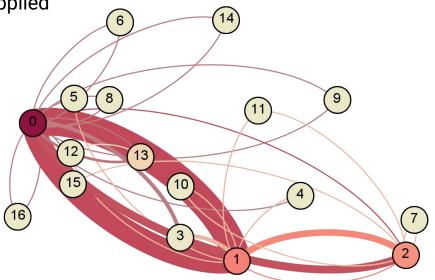


Inferring Home / Work locations

- Basic approach is based on frequency only
 - Most frequent location (L0) := Home
 - Second most frequent location (L1) := Work

A minimum frequency threshold is applied

- Various alternatives & refinement are possible
 - Check time of stop & stay duration
 - Home: stop at 20-22, stay 8-10 hrs
 - Work: stop at 7-10, stay 6-9 hrs



Objective: adding information to points / locations

- Home & work are two particular examples, we want something more general
- E.g: education activity; sport; eating; shopping; ...

Two main ways:

- Assign a single activity
- Assign a distribution of POIs / activity types

Given a dataset of GPS tracks of private vehicles, annotate trajectories with the most probable activities performed by the user.



Associates the list of possible <u>POIs</u> – Points of Interest visited by a user moving by car when he stops.

- POI collection from available sources, e.g. from OSM, Google Places.
- Association POI → Activity: Each POI is associated to an "activity".
 - E.g. "Alfredo's"→Restaurant or "Alfredo's"→ Eating/Food
 - E.g. "CS Library, Pisa" → Education
- Collect basic elements/characteristics:
 - C(POI) = {category, opening hour, location}
 - C(Trajectory) = {stop duration, stop location, time of the day}
 - C(User) = {max walking distance}
- Computation of the probability to visit a POI/ to make an activity: For each POI, the probability of ``being visited" is a function of the POI, the trajectory and the user features.
- Annotated trajectory: The list of possible activities is then associated to a Stop based on the corresponding probability of visiting POIs



POI assignment

Select POIs based on stop proximity



Walk dist = 500 m

Assign probabilities

$$f$$
 (Dentists) = $\frac{1}{4}$
 f (Churches) = $\frac{1}{4}$
 f (Banks) = $\frac{1}{2}$
 f (School) = 0

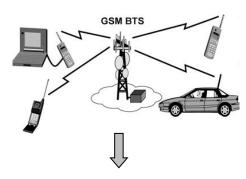


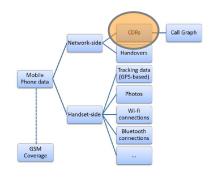
- **X** Bank: Mon Fri [8:00 15:30]
- Dentist: Mon Sat [9:00 13:00] [15:30 18:00]
- Church: Mon Sat [18:00 19:00]
 Sun [11:00 12:00]
- Primary School: Mon Sat [8:00 13:00]

Filter on time constraints (if any)

Inferring Home / Work locations with Phone Data The case of GSM traces

Data gathered from mobile phone operator for billing purpose





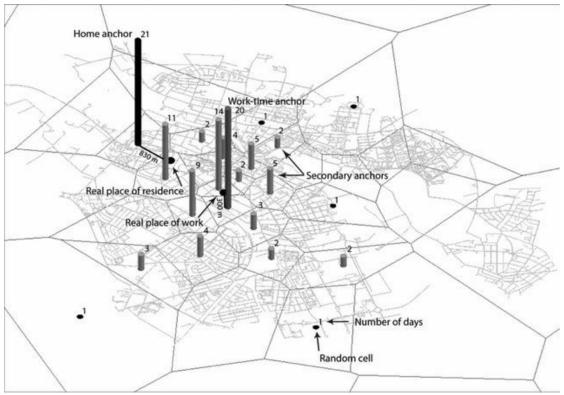
User id	Time start	Cell start	Cell end	Duration
10294595	"2014-02-20 14:24:58"	"PI010U2"	"PI010U1"	48
10294595	"2014-02-20 18:50:22"	"PI002G1"	"PI010U2"	78
10294595	"2014-02-21 09:19:51"	"PI080G1"	"PI016G1"	357

Inferring Home / Work locations with Phone Data Tha case of GSM traces - 1

- "Personal Anchor Points": high-frequency visited places of a user
 - Select top 2 cells with max number of days with calls
 - Determine home and work through time constraints:
 - Based on average start time (AST) of calls and its deviation (std)
 - IF AST<17:00 & std<0.175 ⇒ WORK
 - ELSE HOME

Inferring Home / Work locations with Phone Data Tha case of GSM traces - 1

"Personal Anchor Points"



AHAS, R., SILM, S., JARV, O., SALUVEER, E., AND TIRU, M. 2010. Using mobile positioning data to model locations meaningful to users of mobile phones. Journal of Urban Technology 17, 1, 3–27.

Inferring Home / Work locations with Phone Data Tha case of GSM traces - 2

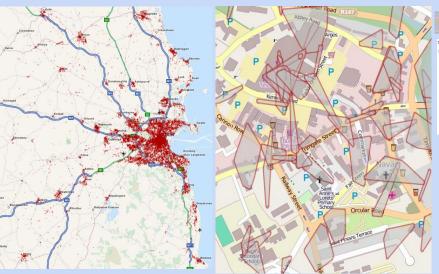
- Estimating users' residence through night activity
 - Home = region with highest frequency of calls during nighttime
 - More suitable for larger scales
 - E.g. region = municipality

Reading social media to find POIs

An Irish experiment on Twitter

The points of each trajectory taken separately were grouped into spatial clusters of maximal radius 150m. For groups with at least 5 points, convex hulls have been built and spatial buffers of small width (5m) around them.

1,461,582 points belong to the clusters (89% of 1,637,346); 24,935 personal places have been extracted.



1150 1000 250

Statistical distribution of the number of places per person

Examples of extracted places

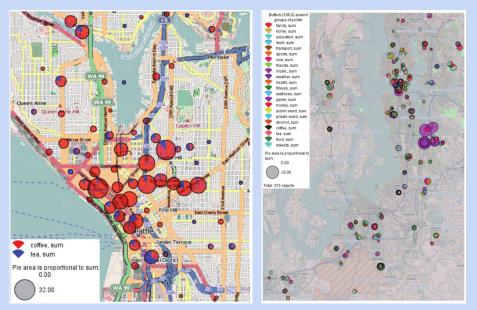
Reading social media to find POIs

Topics have been assigned to 208,391 messages (14.3% of the 1,461,582 points belonging to the personal places)

Message Feat	ures	topic=family: Occurrences of topic		topic=education: Occurrences of topic	topic=work: Occurrences of topic	
@joe_lennon I usually educ	ation	0	0	1	0	
@joe_lennon together educ	ation	0	0	1	0	
@jas_103 deadly; don work	(0	0	0	1	
Just got home and see hom	е	0	1	0	0	
So excited about my ne swe	ets	0	0	0	0	
@lamtcdizzy I haven't b shop	ping	0	0	0	0	
Get in from my night ou fami	ly;home;work	1	1	0	1	
Home again at 6pm! Nhom	R	n	1	0		
Bussing it home for tig Get in from my night out, my dad gets home from work 1 0						
Ah shite. It's been a p two m	inutes later. Gre	at timing :)	0	0	Hacks	
@ronanhutchinson be educ	ation	0	0	1		

- Some places did not get topic summaries (about 20% of the places)
- 2) In many places the topics are very much mixed
- The topics are not necessarily representative of the place type (e.g., topics near a supermarket: family, education, work, cafe, shopping, services, health care, friends, game, private event, food, sweets, coffee)

In the meanwhile, in Seattle...



G. Andrienko et al. Thematic Patterns in Georeferenced Tweets through Space-Time Visual Analytics. Computing in Science & Engineering, 2013.

Homeworks

Homework 4.1 (DIY)

How fast are users?

Choose one of the datasets seen at lesson (taxis, Geolife, etc.), select at least 10 users/vehicles and compute distributions of lengths. Remove 10% of points in each trajectory and repeat the distribution. Do the same for 20%, 30%, ... 90%. How does length distribution change?

Homework 4.2 (DIY)

Estimating GPS errors. Choose a bounding rectangle covering SF city. Download the road network/graph of that area. Select the GPS points of taxis in the same area. Assign each point P to its closest road segment R. Define pseudo-error(P) as the distance dist(P,R).

- Analyze the overall distribution of the pseudo-errors. Is it coherent with GPS.gov estimates of errors?
- Are pseudo-errors the same downtown vs. out of city?

Homework 4.3 (Can we do better?)

Implement a "speed-aware" trajectory compression method, that preserves speed, and test it on a dataset of your choice, e.g. a subset of taxis or Geolife users.

- Show the effects of simplification on some sample trajectories
- Study how the lengths of trajectories are affected

Homework 4.4 (Can we do better?)

The typical filtering algorithm relies on the individual. Define this new algorithm: label the points which have less than $\bf N$ neighbors in the set of points of all trips within a $\bf \delta$ radius as outliers, and discard them.

Test it on SF taxi data

Material

- [paper] Review and classification of trajectory summarisation algorithms: From compression to segmentation
- [paper] Algorithms for the reduction of the number of points required to represent a digitized line or its caricature (Douglas-Peucker)
- [paper] A Trajectory Segmentation Map-Matching Approach for Large-Scale, High-Resolution GPS Data
- [paper] Hidden Markov Map Matching Through Noise and Sparseness
- [chapter] An Introduction to Human Mobility, chapter 4.
 Available soon !!