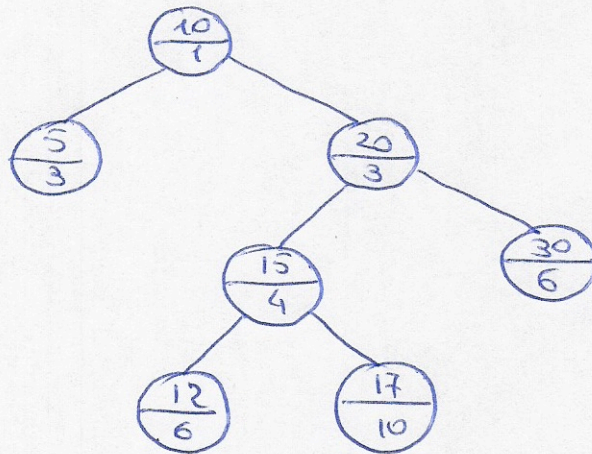
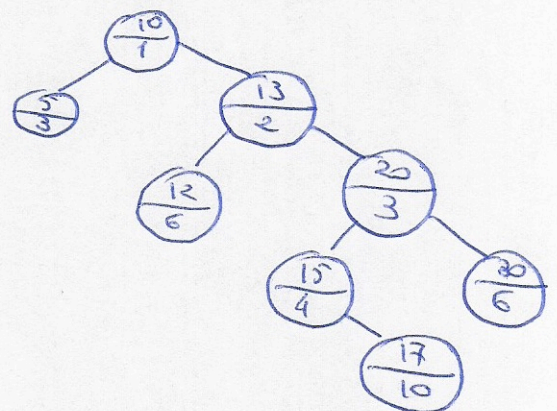
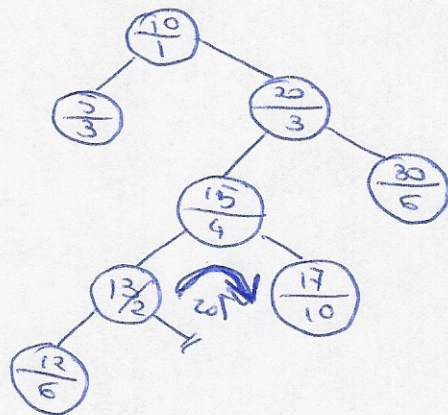
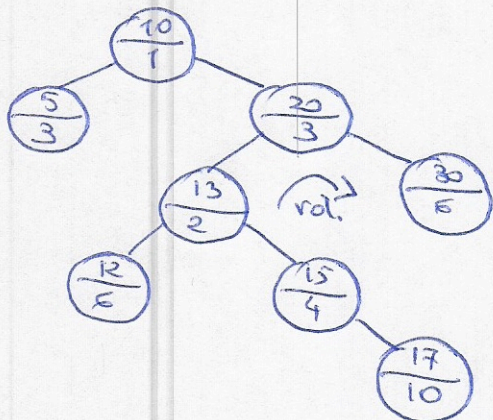
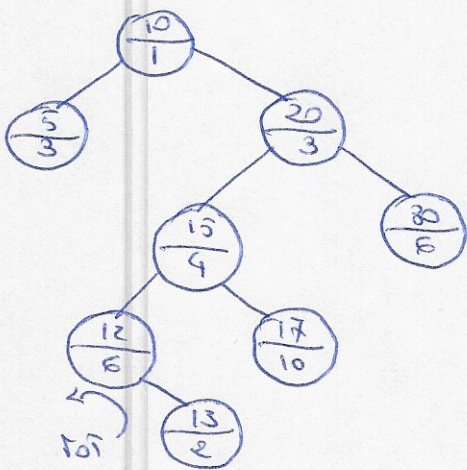


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Q1. The keys, and their priorities, are given in a way that rotations are not needed to guarantee the Treap property.



Then it is required to insert in this Treap the pair $\langle 13, 2 \rangle$. So the first step is to find the correct position to insert (13), then it starts a sequence of rotations that re-establish the underlying heap property.



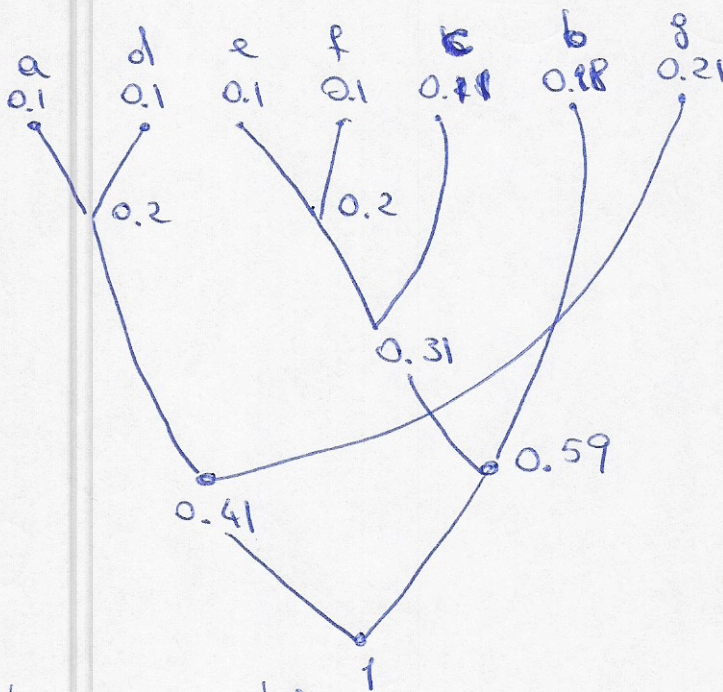
Q2.

1	2
a	b

c	d	e	f
3	1	4	2

Since $m=2$, and the values of c, e are greater than they are not inserted. Conversely, d is inserted in position 1 and f is inserted in position 2, so the final configuration is $\boxed{d|f}$.

Q3.



Σ	level
a	3
b	2
c	3
d	3
e	4
f	4
g	2

num	1	2	3	4
	0	2	3	2

sym b	fc
1	2
2	2
3	1
4	0

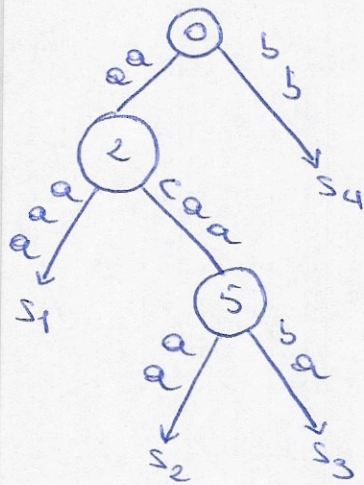
- $v=1 < fc[1]=2 \rightarrow v=(11)_2 = 3 > fc[2]=2$
 output $\text{sym b}[2, v - fc[2]] = \text{sym b}[2, 1] = g$
- $v=0 < fc[1]=2 \rightarrow v=(00)_2 < fc[2]=2 \rightarrow v=(001)_2 \geq fc[3]=1$
 $\rightarrow \text{sym b}[3, v - fc[3]] = \text{sym b}[3, 0] = a$

Q4. We know that the length in bit compressed via Arithmetic is given by the formula $d = \lceil \log_2 \frac{2}{s} \rceil$, where s is $\prod_i p_i(T_i)$

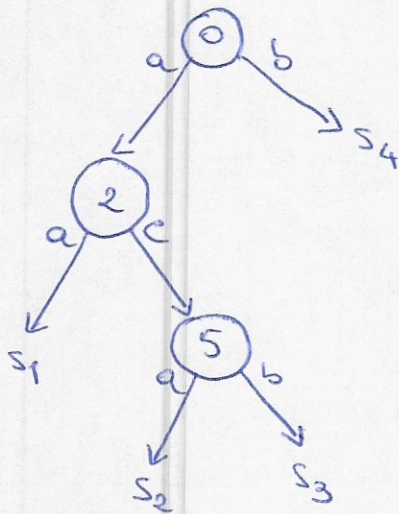
$$= [p(a)]^4 \cdot [p(b)]^2 = \left[\frac{1}{8}\right]^4 \cdot \left[\frac{1}{4}\right]^2 = \left(\frac{1}{2^3}\right)^4 \cdot \left(\frac{1}{2^2}\right)^2 = \frac{1}{2^{16}}$$

$$d = \lceil \log_2 2 \cdot 2^{16} \rceil = \lceil \log_2 2^{17} \rceil = 17 \text{ bit}$$

Q5.



Patricia tree is obtained from the compressed tree to the left by keeping only the first characters of every edge.



① $p = \frac{a}{0} a \frac{c}{2} b b \frac{a}{5} \rightarrow s_2$ is selected

② compute $\text{lcp}(p, s_2) = 3$ $\begin{matrix} \boxed{b} & \boxed{a} \\ p & s_2 \end{matrix}$

③ We percolate upward the Patricia tree until we reach the edge $\begin{matrix} \textcircled{2} \\ \downarrow c \\ \textcircled{5} \end{matrix}$ and since

$\begin{matrix} \boxed{b} & & \boxed{a} \\ p & & s_2 \end{matrix} >$ we have that p is to

the right of the subtree descending from that edge, and thus p is lexicographically between s_3 and s_4 .