## **Models of computation (MOD)** Appello straordinario – April 3, 2019

[Ex. 1] Suppose we add to IMP the command repeat c until b, whose denotational semantics is defined recursively as:

 $\mathcal{C}[\![\mathbf{repeat}\ c\ \mathbf{until}\ b]\!]\sigma = (\lambda \sigma' \mathcal{B}[\![b]\!]\sigma' \ \rightarrow \ \sigma' \ , \ \mathcal{C}[\![\mathbf{repeat}\ c\ \mathbf{until}\ b]\!]\sigma')^* \ (\mathcal{C}[\![c]\!]\sigma)$ 

- 1. Define the operational semantics of the new construct.
- 2. Extend the proof of determinacy of the operational semantics taking into account the new construct.
- 3. Define the function  $\Gamma_{c,b}$  such that  $\mathcal{C}[\![\text{repeat } c \text{ until } b]\!] = fix \Gamma_{c,b}$ .
- 4. Compute the denotational semantics of **repeat** x := x + 1 **until true**.

**[Ex. 2]** Consider the CPO<sub> $\perp$ </sub>  $\mathcal{D} \stackrel{\text{def}}{=} (\wp(\mathbb{N}), \subseteq)$  and the function  $f : \wp(\mathbb{N}) \rightarrow \wp(\mathbb{N})$  such that  $f(X) \stackrel{\text{def}}{=} \{y \mid \exists x \in X. \ y \leq x\}$ , where  $\leq$  is the usual total order on  $\mathbb{N}$ .

- 1. Is f monotone?
- 2. Is f continuous?
- 3. What is the least fixpoint of f? Does f have other fixpoints?

[Ex. 3] Let us consider the CCS processes

$$p \stackrel{\text{def}}{=} \mathbf{rec} \ x.(\tau.x + \beta.\mathbf{nil} + \alpha.(\mathbf{rec} \ y. \ \alpha.x + \tau.y)) \qquad r \stackrel{\text{def}}{=} \mathbf{rec} \ u.(\tau.u)$$
$$q \stackrel{\text{def}}{=} \mathbf{rec} \ z.(\alpha.\alpha.z + \beta.\mathbf{nil})$$

- 1. Draw the LTSs of the processes p and  $s \stackrel{\text{def}}{=} q | r$ .
- 2. Show that p and s are not strong bisimilar.
- 3. Prove that p and s are weak bisimilar.

[Ex. 4] Let  $snd_v$  and  $rec_v$  two predicates representing the fact that the value v is sent and received, respectively.

- 1. Write a LTL formula expressing the property that everytime the value 1 is sent then it is eventually received.
- 2. Write a CTL formula expressing the property that there is a future where the value 2 is sent until it is received.
- 3. Write a  $\mu$ -calculus formula expressing the property that the value 3 is never received.