

ESERCIZIO 1

Si può utilizzare il Counting Sort, opportunamente modificato.

Sort(S)

$B = \text{nuovo array di dim } n;$
 $C = \text{nuovo array di dim } 50;$

// $50 = 382 - 333 + 1$, numero di possibili valori diversi da ordinare.

for ($i=0; i < 50; i++$) $C[i] = 0;$

for ($j=0; j < n; j++$) {

$C[S[j]-333] ++;$

}

for ($i=1; i < 50; i++$) {
 $C[i] = C[i] + C[i-1];$
 }

for ($j=n-1; j > 0; j--$) {

$B[C[S[j]-333]] = S[j];$

$C[S[j]-333] --;$

}

return B;

$$T(n) = O(n)$$

$$S(n) = \cancel{O(1)}$$

(si utilizza un array di dim 50)

ESERCIZIO 2

1) La soluzione del problema si può rappresentare con un array binario B tale che

$$B[i] = 0 \iff i \in S_1$$

$$B[i] = 1 \iff i \in S_2,$$

con $0 \leq i < n$,

2)

Verifica (G, B, k)

$\text{ctr} = 0;$

for ($u=0; u < n; u++$) {

 for ($x = \text{Adj}(u)$. init(); $x \neq \text{null}; x = x.\text{next}$) {

$v = x.\text{dato};$

 if ($B(u) \neq B(v)$) $\text{ctr}++;$

 if ($\text{ctr}/2 > k$) return $\langle \text{FALSE}, \text{ctr}/2 \rangle;$

}

}

return $\langle \text{TRUE}, \text{ctr}/2 \rangle$

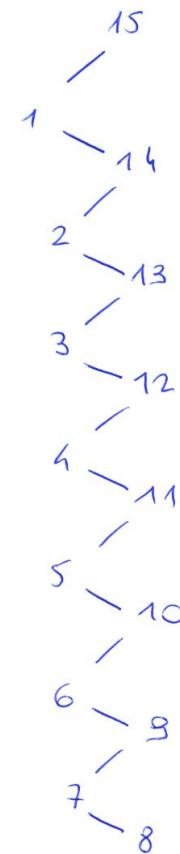
L'algoritmo scorre le liste di adiacenza di tutti i vertici del grafo, dunque la sua complessità è lineare nella dimensione dell'insieme di input:

$$T(V, E) = O(|V| + |E|).$$

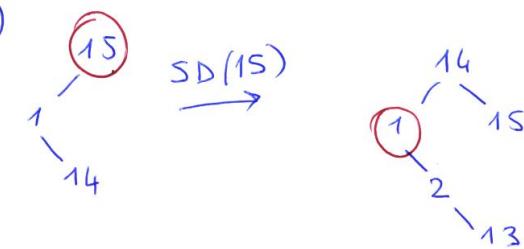
Questo dimostra che MAXCUT $\in \text{NP}$

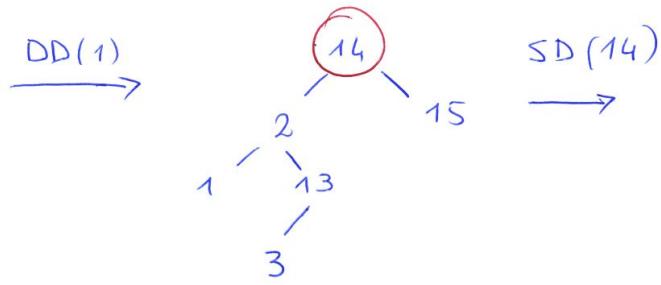
ESERCIZIO 3

1)

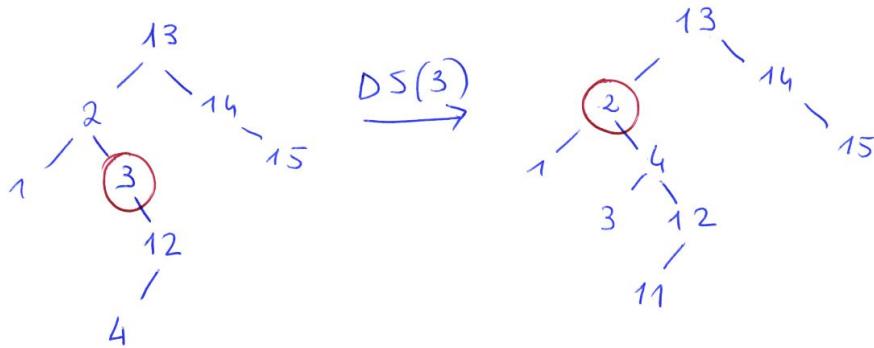


2)

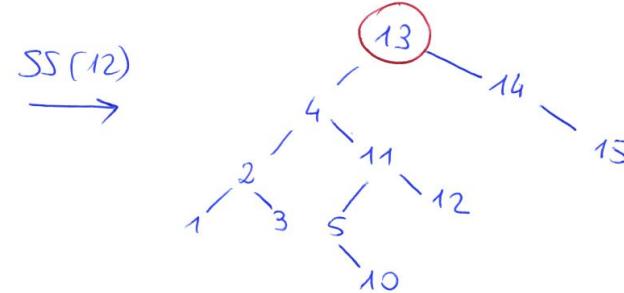
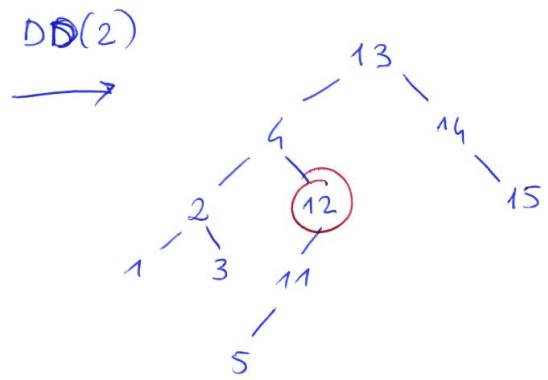
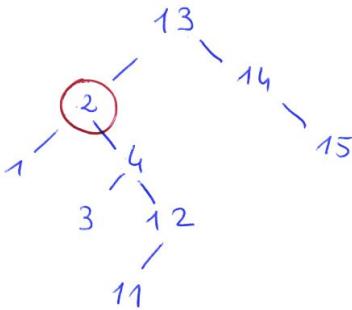




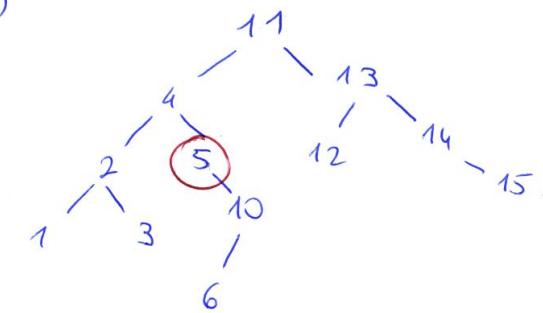
$\text{SD}(14) \rightarrow$



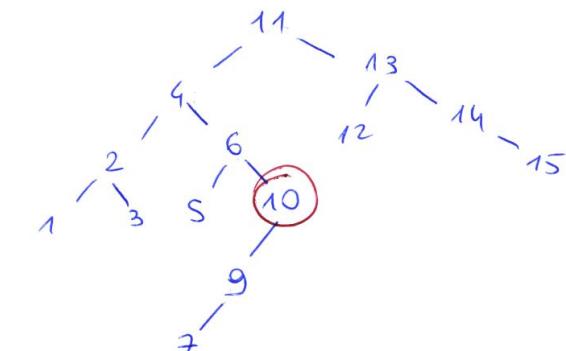
$\text{DS}(3) \rightarrow$



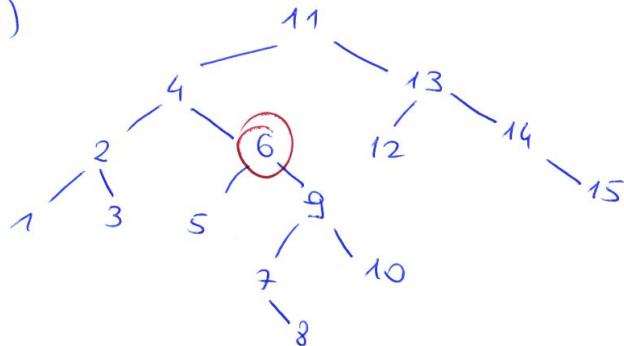
$\text{SD}(13) \rightarrow$



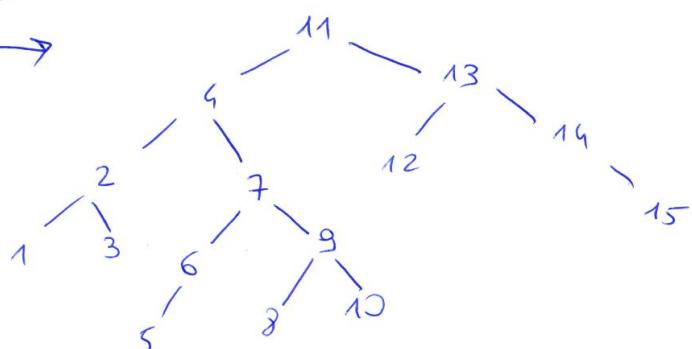
$\text{DS}(5) \rightarrow$



SS(10)



DS(6)



3)

ABR2AVL(u)

A = nuovoArrey();

A = VisitaSinistraCca(u);

return Arrey2AVL(A, 0, n-1);

$$T(n) = \Theta(n)$$

Arrey2AVL(A, sx, dx)

if ($sx > dx$) return null;

$$cx = \frac{sx+dx}{2};$$

u = nuovoNodo();

u.chiave = A(cx);

u.sx = Arrey2AVL(A, sx, cx-1);

u.dx = Arrey2AVL(A, cx+1, dx);

return u;

Dato che l'algoritmo Arrey2AVL è eseguito su un array ordinato, l'ultimo risultante è bilanciato.

$$T(n) = 2T\left(\frac{n}{2}\right) + O(1) = O(n)$$

ESERCIZIO 4

Soluzione ottima (G, k)

```
min = |E|; // globale  
B = nuovo Arey(); // globale  
GeneraBinarié(A, n, k);  
return B;
```

GeneraBinarié (A, b, k)

```
if ( $b == 0$ ) {  
    <sol, k> = Verifica ( $G, A, k$ );  
    if (sol == true) {  
        if ( $c < min$ ) {  
            min = c;  
            copia A in B;  
        }  
    }  
}  
else {  
    A(b-1) = 0;  
    GeneraBinarié ( $A, b-1, k$ );  
    A(b-1) = 1;  
    GeneraBinarié ( $A, b-1, k$ );  
}
```